

Math 61

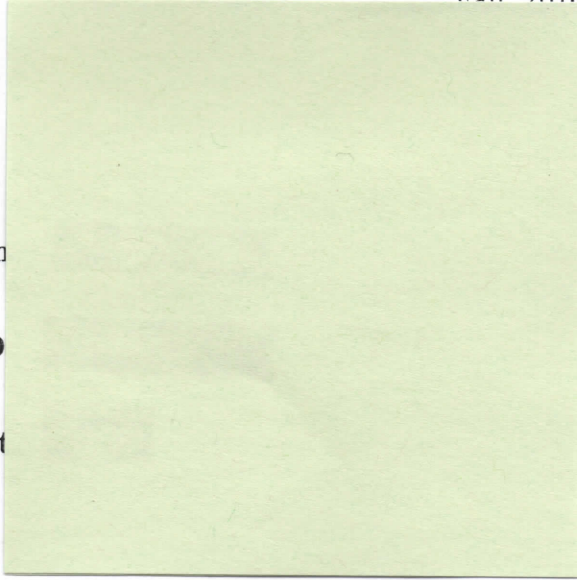
Midterm 2

Fall 2016

Name

SID

Section



There are 5 questions. Write clearly, show all of your work, and justify all of your answers. No calculators are allowed.

1	5
2	10
3	11
4	15
5	14
Total	55

1. (a) (5 pts) Show that $C(i, k) = C(i+1, k+1) - C(i, k+1)$ for $i > k$.

$$\begin{aligned}
 & \binom{i+1}{k+1} - \binom{i}{k+1} \\
 &= \frac{(i+1)!}{(k+1)!((i+1)-(k+1))!} - \frac{i!}{(k+1)!(i-(k+1))!} \\
 &= \frac{i!}{k!(i-k)!} \cdot \frac{i+1}{k+1} - \frac{i!}{k!(i-k)!} \cdot \frac{i-k}{k+1} \\
 &= \frac{i!}{k!(i-k)!} \left(\frac{i+1}{k+1} - \frac{i-k}{k+1} \right) \\
 &= \binom{i}{k} \left(\frac{i+1}{k+1} - \frac{i-k}{k+1} \right) \\
 &\stackrel{<}{=} \binom{i}{k}
 \end{aligned}$$

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- (b) (5 pts) Show that

$$C(n+1, k+1) = \sum_{i=k}^n C(i, k).$$

$$\begin{aligned}
 & \binom{n+1}{k+1} \\
 &= \frac{(n+1)!}{(k+1)!((n+1)-(k+1))!} \\
 &= \frac{n!}{k!(n-k)!} \cdot \frac{n+1}{k+1} \\
 &= \binom{n}{k} \frac{n+1}{k+1}
 \end{aligned}$$

By the binomial theorem, the right side is equal to 2^n , which is the number of possible subsets of a set of size k . But the left side also counts the number of possible subsets of a set of size k , so the two sides are equal.

$$\sum_{i=k}^n \binom{i}{k} = \sum_{i=k}^n \left[\binom{i+1}{k+1} - \binom{i}{k+1} \right] \quad \text{O}$$

$$\begin{aligned}
 &= \left[\binom{k+1}{k+1} - \binom{k}{k+1} \right] + \left[\binom{k+2}{k+1} - \binom{k+1}{k+1} \right] + \left[\binom{k+3}{k+1} - \binom{k+2}{k+1} \right] + \dots + \left[\binom{n+1}{k+1} - \binom{n}{k+1} \right] \\
 &\stackrel{<}{=} \binom{n+1}{k+1}
 \end{aligned}$$

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2. (a) (5 pts) Find the general solution for the recurrence $a_n = -4a_{n-1} - 4a_{n-2}$.

guess: $a_n = \lambda^n \rightarrow \lambda^n = -4\lambda^{n-1} - 4\lambda^{n-2}$

char poly: $0 = \lambda^2 + 4\lambda + 4 = (\lambda + 2)^2 \rightarrow \lambda = -2$

gen soln: $a_n = c_1(-2)^n + c_2 n(-2)^n$



(b) (5 pts) Find the solution to the recurrence $a_n = -4a_{n-1} - 4a_{n-2}$ for $n \geq 2$ with initial conditions $a_0 = 3$ and $a_1 = 0$.

init cond: $\begin{cases} 3 = c_1(-2)^0 + c_2 \cdot 0 \cdot (-2)^0 \\ 0 = c_1(-2)^1 + c_2 \cdot 1 \cdot (-2)^1 \end{cases} \rightarrow \begin{cases} 3 = c_1 \\ 0 = -2c_1 - 2c_2 \end{cases}$

$0 = c_1 + c_2$

$0 = 3 + c_2$

$c_2 = -3$

$\therefore a_n = 3(-2)^n - 3n(-2)^n$



3. (a) (10 pts) Suppose $G = (V, E)$ is a simple graph with n vertices. Prove that if $\deg(v) \geq \frac{n-1}{2}$ for every vertex $v \in V$, then G is connected.

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Assume $n > 1$ (for $n \leq 1$, the graph is trivially "connected"). Let u, v be any two vertices in V . There are two cases:

- u and v are directly connected
- u and v are not ^{directly} connected, but are each connected to at least $\frac{n-1}{2}$ of the $n-2$ remaining vertices that are not u or v . But since $\deg(u) + \deg(v) \geq \frac{n-1}{2} + \frac{n-1}{2} = n-1 > n-2$, there must be at least one vertex connected to both u and v . Then u and v are connected.



\therefore Any two vertices u, v in V are connected, so G is connected.

- (b) (5 pts) Suppose $G = (V, E)$ is a simple graph with n vertices where $n > 1$. Prove that there must be two different vertices of G that have the same degree.

OK

• Base case: For $n = 2$, the two vertices must have the same degree, since all edges are between the two vertices.

• Inductive step: If the proposition is true for $n = k$, then it must be true for $n = k + 1$. When n goes from even to odd \dots

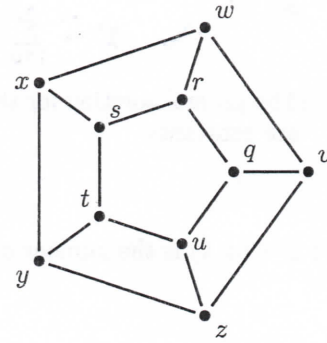
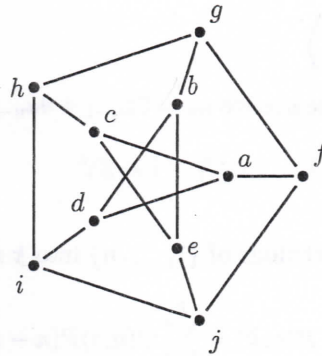
Since G has n vertices, its vertices can have degree $0, 1, \dots, n-1$ for a total of n possible values. But vertices of degree 0 and degree $n-1$ cannot coexist in a single graph, since a vertex of degree $n-1$ must be connected to every other vertex; then there can only be $n-1$ distinct values for the degree of a vertex in G . By the pigeonhole principle, there must be two different vertices of G that have the same degree.

4. (a) (5 pts) State the definition of when two simple graphs $G_1 = (V_1, E_1)$ and $G_2 = (V_2, E_2)$ are isomorphic.

$G_1 = (V_1, E_1)$ and $G_2 = (V_2, E_2)$ are isomorphic iff \exists a bijection $f: V_1 \rightarrow V_2$ st.
 $[\forall u, v \in V_1, (u, v) \in E_1 \text{ iff } (f(u), f(v)) \in E_2]$

- (b) (5 pts) Are the following two graphs isomorphic? If so, give an isorphism using the table below. If not, explain why.

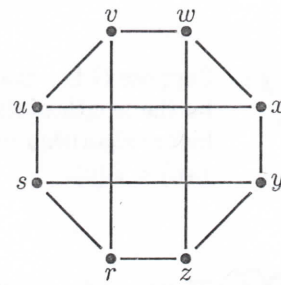
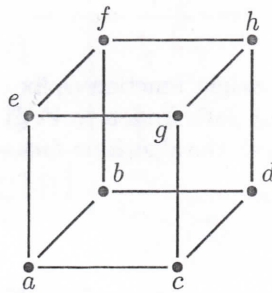
$a \mapsto$
 $b \mapsto$
 $c \mapsto$
 $d \mapsto$
 $e \mapsto$
 $f \mapsto$
 $g \mapsto$
 $h \mapsto$
 $i \mapsto$
 $j \mapsto$



No. The right graph contains 4-cycles but the left graph does not.

- (c) (5 pts) Are the following two graphs isomorphic? If so, give an isomorphism using the table below. If not, explain why.

$a \mapsto r$
 $b \mapsto v$
 $c \mapsto z$
 $d \mapsto w$
 $e \mapsto s$
 $f \mapsto u$
 $g \mapsto y$
 $h \mapsto x$



5. (20 pts) Circle whether the following are True or False. You do not need to justify these answers.

✓ T / (F): There is a simple graph with 6 vertices whose vertices have degree 0, 1, 2, 3, 4, and 4.

✓ (T) / F: The graph $K_{3,3}$ has exactly $3^2 2^4 = 144$ simple cycles.

✗ T / (F): $4^n = \sum_{i=0}^n 2^i C(n, i)$ for every integer $n \geq 0$.

$3^n \quad \rightarrow \quad 2^n = \sum_{i=0}^n \binom{n}{i}$

✓ T / (F): The general solution for the recurrence $a_n = 6a_{n-1} + 9a_{n-2} + 2^n$ is $a_n = b_1 3^n + b_2 2^n$ where b_1, b_2 are constants.

$$0 = \lambda^2 - 6\lambda - 9 = (\lambda - 3)^2$$

✗ (T) / F: If $P(n, k)$ is the number of partitions of $\{1, \dots, n\}$ into k many pieces, then for $n, k > 1$,

$$P(n, k) = \sum_{i=1}^k C(n, i) P(n - i, k - 1).$$

$$P(n+1, k+1) = \sum_{i=k}^n \binom{n}{i} P(i, k)$$

✓ (T) / F: If G is a graph with n vertices, then any path of length n in G must include some vertex at least twice.

✓ T / (F): Suppose A is the adjacency matrix of a graph G with n vertices. If G is connected, then every entry of the matrix A^n is nonzero.

✓ (T) / F: Suppose G is a graph with weight function w , fix a vertex a , and for every vertex v in G , let $L(v)$ be the length of the shortest path from a to v . If we use Dijkstra's algorithm to find $L(v)$, then before algorithm returns $L(v)$, the algorithm first correctly finds $L(u)$ for every vertex u such that $L(u) < L(v)$.

✗ (T) / (F): There are $(n - 1)!$ different isomorphisms from K_n to K_n .

$n!$

✓ (T) / F: If G is a simple graph that has an Euler cycle, and G' is a subgraph obtained from G by removing only a single edge and removing no vertices, then G' cannot have an Euler cycle.

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