

# Math 61-1 Final exam

BENJAMIN HE

TOTAL POINTS

**55 / 90**

QUESTION 1

Multiple choice 10 pts

1.1 2 / 2

- ✓ - 0 pts Correct (c)
- 2 pts Incorrect

1.2 0 / 2

- 0 pts Correct (a)
- ✓ - 2 pts Incorrect

1.3 2 / 2

- ✓ - 0 pts Correct (c)
- 2 pts Incorrect

1.4 2 / 2

- ✓ - 0 pts Correct (b)
- 2 pts incorrect

1.5 0 / 2

- 0 pts Correct (a)
- ✓ - 2 pts Incorrect

QUESTION 2

Short answer 10 pts

2.1 2 / 2

- ✓ - 0 pts Correct  $((-2)^{100} + 3^{100})$
- 1 pts Almost correct (small arithmetic error in answer)
- 2 pts Incorrect

2.2 2 / 2

- ✓ - 0 pts Correct  $(C(7,4)6!4!)$
- 1 pts Close
- 2 pts Incorrect

2.3 2 / 2

- ✓ - 0 pts Correct (24C4)
- 1 pts Close
- 2 pts incorrect

2.4 0 / 2

- 0 pts Correct
- 1 pts Close (Three of four)
- ✓ - 2 pts Incorrect

2.5 1 / 2

- 0 pts Correct  $(2^{(n^2 - n)} + 2^{(n^2 + n / 2)} - 2^{(n^2 - n / 2)})$
- ✓ - 1 pts Close
- 2 pts Incorrect

QUESTION 3

Equivalence relation 10 pts

3.1 it is an equivalence relation 4 / 4

- ✓ - 0 pts Correct
- 1 pts issue in transitivity
- 3 pts misunderstanding of what relation is saying
- 4 pts blank
- 2 pts misunderstanding of symmetry
- 1 pts the decimal thing isn't exactly right, e.g.  $-3$  is related to  $.7$
- 0 pts Click here to replace this description.
- 1 pts issue with symmetry

3.2 defining a function 3 / 4

- 0 pts Correct
- 4 pts blank
- 2 pts need to prove uniqueness part of function
- 2 pts missing existence part of function
- 1 pts issue with uniqueness part of function

✓ - 1 pts need to consider different elements in the same equivalence class

- 1 pts thing with decimals isn't quite right, for example -3 and .7 are related

- 3 pts big misunderstanding of the equivalence relation or function

### 3.3 a function that doesn't descend 1 / 2

- 0 pts Correct

- 2 pts your g is not a function

✓ - 1 pts issue with justification

- 1 pts your g does not work

- 2 pts blank

☞ why?

### QUESTION 4

#### m-ary tree 10 pts

##### 4.1 number of internal vertices 0 / 5

- 0 pts Correct

✓ - 1 pts No/incorrect answer

✓ - 4 pts No/incorrect justification

- 2 pts Didn't justify number of total vertices

- 3 pts "Proof by example"

- 2 pts Assumed every terminal vertex had the same height as the tree

- 5 pts Nothing

- 1 pts Forgot to account for root

- 2 pts Didn't subtract off internal vertices

##### 4.2 height 1.5 / 5

- 0 pts Correct

- 1 pts No base case

- 1 pts Didn't set up/invoke induction

- 1 pts Backwards inductive step (didn't show inductive construction is exhaustive)

- 2 pts Compared to complete tree without showing this case is extremal

✓ - 3 pts Assumed tree is complete / inductive construction forms complete trees from complete trees

- 1 pts Assumed all immediate subtrees have height h-1

- 4 pts "Proof by example"

- 5 pts Nothing shown / Incorrect reasoning

- 0.5 Point adjustment

☞ Your t and h mean two different things each! This is bad.

### QUESTION 5

#### spanning trees 10 pts

##### 5.1 unique mst 3 / 6

- 0 pts Correct

- 3 pts Appeal to Prim's or Kruskal's Algorithm (without proving it can generate any MST)

- 6 pts No / Invalid reasoning

- 3 Point adjustment

☞ What if there are MST's which differ by more than a single pair of edges?

##### 5.2 non unique spanning tree 4 / 4

✓ - 0 pts Correct

- 4 pts Not an example

- 4 pts Claimed no such graph exists

- 4 pts Nothing

### QUESTION 6

#### planar graphs 10 pts

##### 6.1 $2e > 3f$ 3 / 3

✓ + 3 pts Correct

+ 2 pts  $\geq 3$  edges for each face

+ 1 pts  $\geq 3$  edges for each face (w/ mistake)

+ 1 pts  $\leq 2$  faces for each edge

+ 0 pts Incorrect

##### 6.2 $e < 3v - 6$ 3 / 3

✓ + 3 pts Correct

+ 2 pts Euler's formula

+ 1 pts Correct application with (a)

+ 0 pts Incorrect

### 6.3 nonplanar graph 1 / 4

- + 4 pts Correct
- + 3 pts Isomorphic to  $K_{3,3}$
- + 2 pts Mistaken/missing isomorphism to  $K_{3,3}$
- + 1 pts  $E \leq 2v-4$  or  $2E \geq 4F$
- ✓ + 1 pts Other partial credit
- + 0 pts Incorrect

#### QUESTION 7

10 pts

### 7.1 $7^n - 1$ divisible by 6 5 / 5

- ✓ + 5 pts Correct
- + 1 pts Base case
- + 1 pts Inductive hypothesis
- + 2 pts factoring out a 7 in inductive step as  $(6+1)$  or adding/subtracting 7
- + 1 pts Conclusion
- + 0 pts Incorrect

### 7.2 number with only 1s divisible by 7 0 / 5

- + 5 pts Correct
- ✓ + 0 pts [Click here to replace this description.](#)
- + 1 pts Look at 8 consecutive terms
- + 1 pts Pigeonhole remainder
- + 1 pts 7 divides a number of the form 111...000...
- + 2 pts This implies that 7 divides  $10^k - 1$
- + 1 pts Unsuccessful attempt with substantial work

#### QUESTION 8

### balanced binary trees 10 pts

#### 8.1 4 / 4

- ✓ - 0 pts Correct
- 2 pts incomplete, need to describe how a height  $n$  minimal balanced binary tree is made out of ones of smaller height
- 3 pts can't just do examples
- 4 pts blank
- 1 pts how are you adding in these trees/ vertices?
- 3 pts can't do induction without using some properties of minimal balanced binary trees

- 4 pts incorrect numbers/ equation

### 8.2 relationship to fibonacci numbers 3 / 3

- ✓ - 0 pts Correct
- 1.5 pts that is not the recurrence/ equation for the fibonacci numbers/ minimal balanced binary trees
- 1 pts you are assuming the desired conclusion
- 3 pts blank
- 1.5 pts need to use recurrence for fibonacci numbers
- 1.5 pts missing inductive step
- 1 pts the two recurrences aren't exactly the same, you need to account for this difference
- 0.5 pts error in equations
- 1 pts need to check initial conditions

### 8.3 Theta 2.5 / 3

- 0 pts Correct
- ✓ - 0.5 pts need to account for other term in equation for fibonacci numbers (sometimes it is contributing something positive, something something negative)
- 2 pts wrong formula for fibonacci numbers/  $v_n$
- 1 pts issue with big O
- 1 pts issue with omega
- 3 pts blank/ no gradable work
- 1 pts wrong equations/ issues with constants
- 2 pts need to use equation for  $v_n$ / Fibonacci numbers

#### QUESTION 9

### binomial coefficients 10 pts

#### 9.1 $3^n$ 4 / 4

- ✓ + 4 pts Correct
- + 3 pts Minor error
- + 2 pts Binomial theorem
- + 1 pts Attempted induction or counting argument
- + 0 pts Incorrect

### 9.2 vandermonde identity 0 / 6

- + 6 pts Correct

+ **5 pts** Minor error

+ **3 pts** One part of counting argument or  $(x+y)^{n+m}$

+ **1 pts** Attempted to use induction/binomial  
thrm/Pascal's identity

✓ + **0 pts** **Incorrect**

# Final

Name: Benjamin He

Student ID: 804962948

Section:

Tuesday:

Thursday:

1A

1B

TA: Albert Zheng

1C

1D

TA: Benjamin Spitz

1E

1F

TA: Eilon Reisin-Tzur

**Instructions:** Do not open this exam until instructed to do so. Please print your name and student ID number above, and circle the number of your discussion section. **You may not use calculators, books, notes, or any other material to help you.** Please make sure **your phone is silenced** and stowed where you cannot see it. Remember that you are bound by a conduct code.  
**Please get out your id and be ready to show it during the exam.**

Please do not write below this line.

Question	Points	Score
1	10	
2	10	
3	10	
4	10	
5	10	
6	10	
7	10	
8	10	
9	10	
Total:	90	

1. (10 points) Circle the correct answer (only one answer is correct for each question)

$$1. \frac{n!}{k!(n-k)!} + \frac{n!}{(k+1)!(n-k-1)!} =$$

(a)  $\frac{(n+k)!}{k!n!}$

(b)  $\frac{(n+1)!}{k!(n+1-k)!}$

(c)  $\frac{(n+1)!}{(k+1)!(n-k)!}$

(d) none of the above

$\frac{n!}{k!}$

$$\binom{n}{k} + \binom{n}{k+1} = \binom{n+1}{k+1}$$

$$n+1 - (k+1)$$

2) The decision tree of a sorting algorithm for sorting  $n$  items (where at each step we can only decide whether or not one item is less than other) necessarily has:

(a) a height of  $\geq \lg(n!)$

(b) a height of  $\Omega \lg(n!)$  (but not necessarily a height of  $\geq \lg(n!)$ )

(c) a height of  $O(\lg(n!))$

(d) a height of  $O(n \lg n)$

3. If  $G$  is a graph with  $n$  vertices and  $n - 2$  edges, then:

(a)  $G$  is a tree

(b)  $G$  is connected

(c)  $G$  is disconnected

(d)  $G$  is simple

Question 1 continues on the next page...

Question 1 continued...

4. Which of these graphs has an Euler cycle?

- (a)  $K_4$
- (b)  $K_5$
- (c)  $K_{3,3}$
- (d)  $K_{2,3}$

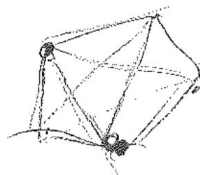
*concluded  
all degrees even*



5. What is the *fewest* number of edges (i.e. in the best case) that could be examined by Dijkstra's algorithm on a graph with  $n$  vertices? (We examine edges in the part of the algorithm where we update labels.)  
Your answer should be true for all  $n$ .

- (a) Less than or equal to  $n$
- (b) More than  $n$  but less than or equal to  $n^2/2$
- (c) More than  $n^2/2$  but less than or equal to  $n^2$
- (d) More than  $n^2$

$n(n-1)$



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8

2. In this question write down your answer, no need for any justification. Leave your answers in a form involving factorials,  $P(n, m)$ ,  $\binom{n}{m}$ , exponents, etc.

(a) (2 points) If  $s_n = s_{n-1} + 6s_{n-2}$  and  $s_0 = 2, s_1 = 1$ , what is  $s_{100}$ ?

$$s_n = t^n$$

$$t^n = t^{n-1} + 6t^{n-2} \rightarrow t^2 - t - 6 = 0$$

$$(t-3)(t+2) = 0 \rightarrow t=3, t=-2$$


$$s_n = a3^n + b(-2)^n \quad a=1, b=1$$

$$a+b=2$$

$$3a-2b=1$$

$$s_{100} = 3^{100} + (-2)^{100}$$

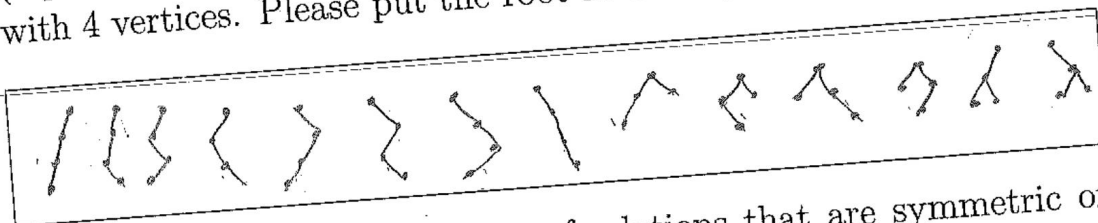
(b) (2 points) How many ways can 7 distinct math majors and 4 distinct CS majors sit in a circle, if the CS majors won't sit by each other and we say that two seatings are the same if they are related by a rotation?

$$6! P(7, 4)$$


(c) (2 points) A squirrel has 20 identical acorns that she is going to hide among 5 distinct holes. In how many ways can the squirrel hide the acorns?

$$\binom{24}{4} = \binom{24}{20}$$

(d) (2 points) Draw all the distinct (up to isomorphism) rooted trees with 4 vertices. Please put the root at the top.



(e) (2 points) What is the number of relations that are symmetric or reflexive on a set with  $n$ -elements?

$$\frac{1}{2}(n^2 - n) + 2^n - 2^{\frac{1}{2}(n^2 - n)} = 2^{\frac{1}{2}(n^2 - n)} (2^n - 1)$$



3. Consider the relation on the real numbers defined by  $C = \{(x, y) \in \mathbb{R} \times \mathbb{R} : x - y \in \mathbb{Z}\}$  <sup>real</sup>  $\downarrow$  integers

(a) (4 points) Show that  $C$  is an equivalence relation.

Reflexive:

Any number minus itself is 0, which is an integer, thus for any ~~pair~~ pair  $(x, x)$  w/  $x \in \mathbb{R}$ ,  $(x, x) \in C$  making it reflexive

Symmetric:

If  $(x, y) \in C$ , then  $x - y = \text{integer} = -(y - x)$ , therefore  $y - x$  must also be an integer, so  $(y, x) \in C$  and this relation is symmetric:

Transitive:

If  $(x, y), (y, z) \in C$ , then  $x - y = \text{integer}$  and  $y - z = \text{integer}$  <sup>that</sup> ~~is~~ <sup>another</sup> ~~integer~~ <sup>integer</sup>  
 $x - y + (y - z) = \text{integer} + \text{integer} = x - z = \text{integer}$  b/c  
an integer minus another integer is always an integer. Thus, if  $(x, y), (y, z) \in C$ , then  $(x, z) \in C$  making  $C$  transitive.

Since  $C$  is Reflexive, Symmetric, and Transitive, it's an equivalence relation.

- (b) (4 points) Let  $\tilde{\mathbb{R}}$  denote the set of equivalence classes of  $\mathbb{C}$ , i.e.  $\tilde{\mathbb{R}} = \{[x] : x \in \mathbb{R}\}$ . Consider the function  $f : \mathbb{R} \rightarrow \mathbb{R}$  defined by  $f(x) = x + 1/2$ . Show that the relation  $\tilde{f}$  from  $\tilde{\mathbb{R}}$  to  $\tilde{\mathbb{R}}$  defined by  $\tilde{f} = \{([a], [b]) \in \tilde{\mathbb{R}} \times \tilde{\mathbb{R}} : f(a) = b\}$  is a function.

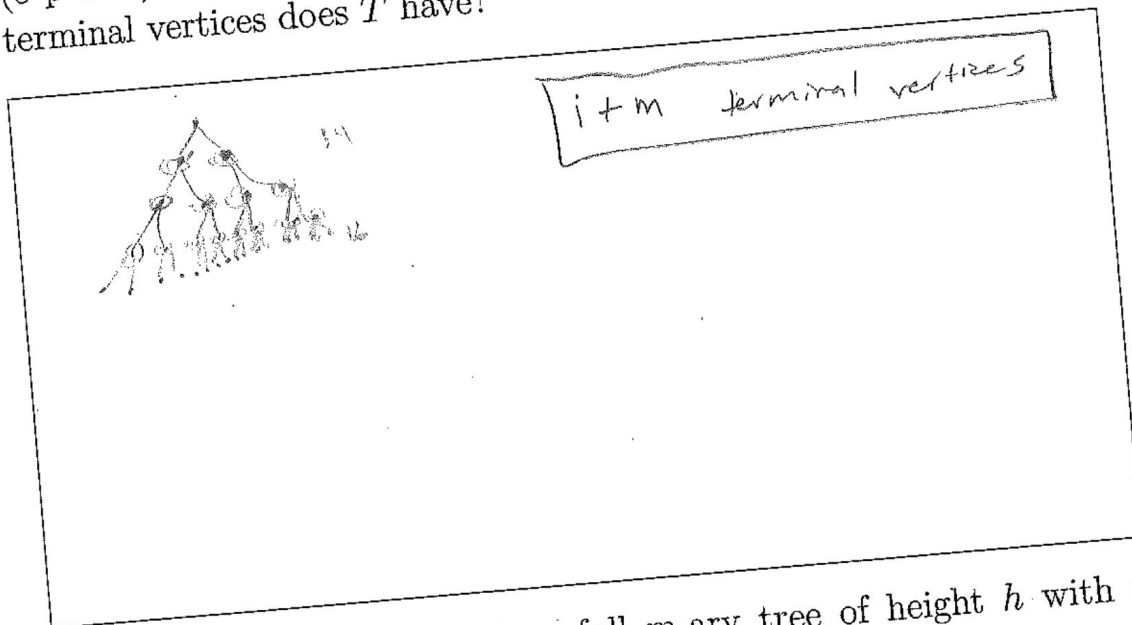
Function: for each  $x \in X$  there is a unique  $y \in Y$  where  $f(x) = y$   
 $\tilde{f}$  is ~~not~~ injective since  $f(a_1) = a_1 + \frac{1}{2}$  and  $f(a_2) = a_2 + \frac{1}{2}$   
 if  $f(a_1) = f(a_2) \rightarrow a_1 + \frac{1}{2} = a_2 + \frac{1}{2} \rightarrow a_1 = a_2$ , so for every  $a$  there's only one  $b$ . Since  $[a]$  are all numbers that are an integer distance from  $a$  and same w/  $[b]$  (all numbers are integer distance from  $b$ ) then if  $f(a) = b$  then every element in  $[a]$  will have a unique corresponding one in  $[b]$  satisfying the function. Thus for every  $[a]$ , where  $a \in \mathbb{R}$ , it will only map to one  $[b]$ , making  $\tilde{f}$  a function. This is b/c there is a uniform distance between <sup>consecutive</sup>  $f(x)$  values

- (c) (2 points) Give an example of a function  $g : \mathbb{R} \rightarrow \mathbb{R}$  so that the relation  $\tilde{g}$  from  $\tilde{\mathbb{R}}$  to  $\tilde{\mathbb{R}}$  defined by  $\tilde{g} = \{([a], [b]) \in \tilde{\mathbb{R}} \times \tilde{\mathbb{R}} : g(a) = b\}$  is not a function. (Be sure to justify your answer.)

$g(a) = a^2$  would have  $[a]$  mapped to multiple  $[b]$ 's. Since the distances between consecutive  $g(a)$  change so each value in  $[a]$  can associate with a different  $[b]$  set due to the nature of the  $\mathbb{C}$  relation.

4. For  $m$  a positive integer, a full  $m$ -ary tree is a rooted tree where every parent has exactly  $m$  children.

(a) (5 points) If  $T$  is a full  $m$ -ary tree with  $i$  internal vertices, how many terminal vertices does  $T$  have?



(b) (5 points) Show that if  $T$  is a full  $m$ -ary tree of height  $h$  with  $t$  terminal vertices, then  $t \leq m^h$ .

Base case:  $\forall h=0, t=1$  for any  $m$  and  $1 \leq m^0 \rightarrow 1 \leq 1$  ✓  
 Base case is true.

Inductive step: Assume  $t \leq m^h$ , then adding another

layer will cause  $t = mt$  since it's a full  $m$ -ary tree.

Thus  $mt \leq m^{h+1} \Rightarrow m(t) \leq m(m^h) \rightarrow t \leq m^h$

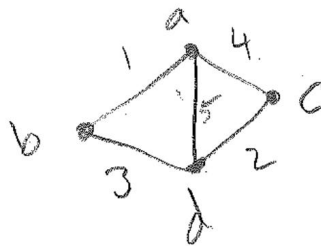
Thus if  $t \leq m^h$  is true for a given full  $m$ -ary tree of height  $h$ , then it's also true for given full  $m$ -ary tree of height  $h+1$ . Thus

$t \leq m^h$  is true for any full  $m$ -ary tree of height  $h$  and  $t$  terminal vertices.  $\square$

5. (a) (6 points) Show that if  $G$  is a connected weighted graph where all the edges of  $G$  have distinct weights then  $G$  has a unique minimal spanning tree.

Since  $G$  is connected, we know that it has a minimal spanning tree. Let's assume that there are 2 minimal spanning trees for  $G$  where all edges have distinct weights. Let's say you find one minimal spanning tree  $T$ . You can then connect any two vertices to create a cycle and then remove an original edge in  $T$  to create a new ~~edge~~ spanning tree  $T'$  however, because all the weights are distinct, ~~the~~  $\text{weight}(T') \neq \text{weight}(T)$  meaning that one tree will be "heavier" than the other, which contradicts the claim that there are 2 min spanning trees. The same logic is applicable to having 3+ min spanning trees. Thus by contradiction  $G$  must have a single unique min spanning tree if all its edges have distinct weights.

(b) (4 points) Give an example of a connected weighted graph  $G$  so that all the edges of  $G$  have distinct weights and  $G$  has at least two distinct spanning trees that have the same total weight, i.e. the sums of the weights of the edges in these two distinct trees agree, or prove that no such weighted graph exists.



6. (a) (3 points) Show that for  $G$  a connected simple planar graph containing a cycle if  $G$  has  $E$  edges and  $F$  faces, then  $2E \geq 3F$ .

Every edge can be a part of at most 2 faces.

Every face must be ~~made of~~ <sup>enclosed by</sup> at least 3 edges

$$\text{Thus } 2E \geq 3F$$

- (b) (3 points) Show that for  $G$  a connected simple planar graph containing a cycle if  $G$  has  $E$  edges and  $V$  vertices, then  $E \leq 3V - 6$ .

For any simple planar graph,

$$V - E + F = 2 \rightarrow E = V + F - 2$$

$$2E \geq 3F$$

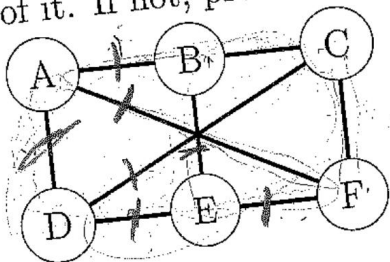
$$\rightarrow 3E \geq 3V + 3F - 6$$

since  $2E \geq 3F$  (part (a))

$$3E \leq 3V + 2E - 6$$

$$\Rightarrow \boxed{E \leq 3V - 6}$$

(c) (4 points) Is the following graph planar? If it is give a planar drawing of it. If not, prove that it is not planar.



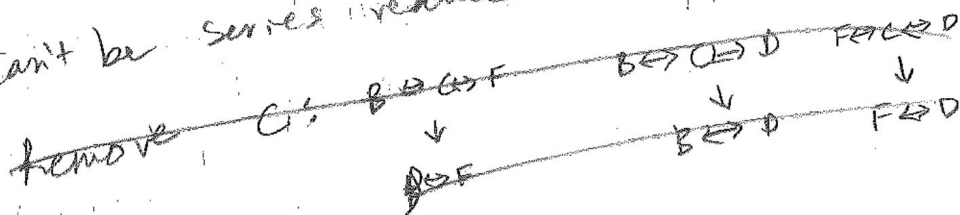
$$V = 4$$

$$E = 5$$

$$F = 3$$

Not planar if a subgraph ~~is~~ can be reduced to  $K_5$  or  $K_{3,3}$ .

This graph is ~~not~~ planar b/c it can't be series reduced to  $K_5$  or  $K_{3,3}$ .



$$18 \leq 4F$$

$$4 \leq \frac{18}{2} = 9$$

Also fulfills Euler's Theorem

$$V - E + F = 2$$

$$6 - 9 + 3 = 0$$

7. (a) (5 points) Show that for all  $n \geq 1$ ,  $7^n - 1$  is divisible by 6.

Base case:  $n=1$ ,  $7^1 - 1 = 6$  which is divisible by 6, thus base case satisfied.

Inductive step: Assume  $7^n - 1$  is div by 6, then we want to prove  $7^{(n+1)} - 1$  also div by 6:

$$7^{(n+1)} - 1 = 7(7^n) - 1 = 7(7^n) - 7 + 6$$

$$= 7(7^n - 1) + 6 \rightarrow \text{since } 7^n - 1 \text{ div by 6,}$$

$\rightarrow 7^n - 1 = 6 \cdot x$  for some integer  $x$

$$\rightarrow 7(7^n - 1) + 6 = 7 \cdot 6(x) + 6 = 6(7x + 1) \text{ which is div by 6, therefore } 7^n - 1 \text{ is div by 6 for all } n \geq 1 \quad \square$$

(b) (5 points) Show that there is a number of the form  $\sum_{i=0}^n 10^i$  (i.e. a number consisting only of 1s) that is divisible by 7.

$$\sum_{i=0}^n 10^i = 7m \text{ for some } m.$$



8. A *balanced binary tree* is a binary tree where for each vertex the heights of the left and right subtrees of that vertex differ by at most one. Let  $v_n$  denote the minimum number of vertices in a balanced binary tree of height  $n$ .

(a) (4 points) Show that  $v_n$  satisfies for  $n \geq 2$  the recurrence  $v_n = v_{n-1} + v_{n-2} + 1$

Since the trees for  $v_{n-1}$  and  $v_{n-2}$  must both be balanced, and they differ by a height of 1, they can simply be chained together with a parent root. Thus, you add the vertices of the two older trees plus 1 for the new root node.

Ex:  $v_0 = 1$      $v_1 = 2$      $v_2 = 4 = v_0 + v_1 + 1$      $v_3 = 7 = v_1 + v_2 + 1$      $v_n = v_{n-1} + v_{n-2} + 1$

(b) (3 points) Show that for  $n \geq 0$ ,  $v_n = F_{n+2}$  where  $F_k$  is the  $k^{\text{th}}$  Fibonacci number.  
 $v_n = F_{n+3} - 1$

$v_n = v_{n-1} + v_{n-2} + 1 \rightarrow$  define  $z_n = v_n + 1$   
 $v_0 = 1, v_1 = 2$      $z_{n-1} = v_n$   
 then

$z_{n+1} = z_{n-1} + z_{n-2} + 1 + 1 = z_{n-1} + z_{n-2} + 2$

$\rightarrow z_n = z_{n-1} + z_{n-2}$  with  $z_0 = 2, z_1 = 3$

The Fibonacci sequence's recurrence relation is also  $F_n = F_{n-1} + F_{n-2}$  but  $F_0 = 0$  and  $F_1 = 1, F_2 = 1, F_3 = 2, F_4 = 3$

Here,  $z_0 = F_3$  and  $z_1 = F_4$ , so  $z_n = F_{n+3}$  but since  $z_n = v_n + 1$

$\rightarrow v_n + 1 = F_{n+3} \rightarrow v_n = F_{n+3} - 1$

(c) (3 points) Show that  $v_n = \Theta(\phi^{n+2})$ , where  $\phi = \frac{1+\sqrt{5}}{2}$ .

$$F_n = \frac{1}{\sqrt{5}} \left( \frac{1+\sqrt{5}}{2} \right)^n - \frac{1}{\sqrt{5}} \left( \frac{1-\sqrt{5}}{2} \right)^n$$

$$v_n = F_{n+3} - 1 = \frac{1}{\sqrt{5}} \left( \frac{1+\sqrt{5}}{2} \right)^{n+3} - \frac{1}{\sqrt{5}} \left( \frac{1-\sqrt{5}}{2} \right)^{n+3} - 1$$

$$\leq \frac{1}{\sqrt{5}} \left( \frac{1+\sqrt{5}}{2} \right)^{n+3} \leq \frac{\sqrt{5}+5}{10} \left( \frac{1+\sqrt{5}}{2} \right)^{n+2}$$

then  $v_n \leq \frac{\sqrt{5}+5}{10} (\phi^{n+2})$

so  $v_n = O(\phi^{n+2})$

$$v_n = \frac{1}{\sqrt{5}} \left( \frac{1+\sqrt{5}}{2} \right)^{n+3} - \frac{1}{5} \left( \frac{1-\sqrt{5}}{2} \right)^{n+3} - 1$$

since  $\frac{1+\sqrt{5}}{2} > \frac{1-\sqrt{5}}{2}$  and  $\frac{1+\sqrt{5}}{2} \geq 1$ ,

$$v_n = \frac{1}{\sqrt{5}} \left( \frac{1+\sqrt{5}}{2} \right)^{n+3} - \frac{1}{5} \left( \frac{1-\sqrt{5}}{2} \right)^{n+3} - 1 \geq \frac{1}{\sqrt{5}} \left( \frac{1+\sqrt{5}}{2} \right)^{n+3} - \frac{1}{\sqrt{5}} \left( \frac{1+\sqrt{5}}{2} \right)^{n+3} - \left( \frac{1+\sqrt{5}}{2} \right)^{n+3}$$

$$v_n \geq \frac{1}{\sqrt{5}} (\phi)^{n+3} - \frac{1}{\sqrt{5}} (\phi)^{n+3} - \phi^{n+3}$$

$$v_n \geq -\left( \frac{1+\sqrt{5}}{2} \right) \phi^{n+2} \text{ therefore } v_n = \Omega(\phi^{n+2})$$

Since  $v_n = O(\phi^{n+2})$  and  $v_n = \Omega(\phi^{n+2})$ ,  $v_n = \Theta(\phi^{n+2})$

9. (a) (4 points) Show that  $\sum_{i=0}^n 2^i \binom{n}{i} = 3^n$ .

$(x+z)^n = \sum_{i=0}^n \binom{n}{i} x^{n-i} z^i$  v/c Binomial thm.  
 let  $x=1$ :  $(1+z)^n = 3^n = \sum_{i=0}^n \binom{n}{i} z^i$

(b) (6 points) Show that  $\binom{n+m}{r} = \sum_{i=0}^r \binom{n}{i} \binom{m}{r-i}$ .

$\binom{n+m}{r} = \binom{n}{r} + \binom{n}{r-1} + \dots$   
 $\frac{n! m!}{\sum_{i=0}^r (i! (r-i)! (n-i)! (m-r+i)!)} = n! m! \sum_{i=0}^r \frac{1}{i! (r-i)! (n-i)! (m-r+i)!}$   
 $= n! m! \left( \frac{1}{r! n! (m-r)!} + \frac{1}{(r-1)! (n-1)! (m-r+1)!} + \dots \right)$   
 $= \frac{n! m!}{r! n! (m-r)!} (1 + (r)(n)(m-r) + (r)(r-1)(n)(n-1)(m-r)(m-r-1) + \dots)$   
 $= \binom{m}{r} \binom{n}{r-i}$

This page has been left intentionally blank. You may use it as scratch paper. It will not be graded unless indicated very clearly here and next to the relevant question.