

Math 61-1 Final exam

TOTAL POINTS

79 / 90

QUESTION 1

Multiple choice 10 pts

1.1 0 / 2

- 0 pts Correct (c)

✓ - 2 pts Incorrect

1.2 2 / 2

✓ - 0 pts Correct (a)

- 2 pts Incorrect

1.3 2 / 2

✓ - 0 pts Correct (c)

- 2 pts Incorrect

1.4 2 / 2

✓ - 0 pts Correct (b)

- 2 pts incorrect

1.5 2 / 2

✓ - 0 pts Correct (a)

- 2 pts Incorrect

QUESTION 2

Short answer 10 pts

2.1 2 / 2

✓ - 0 pts Correct $(-2)^{100} + 3^{100}$

- 1 pts Almost correct (small arithmetic error in answer)

- 2 pts Incorrect

2.2 0 / 2

- 0 pts Correct $(C(7,4)6!4!)$

- 1 pts Close

✓ - 2 pts Incorrect

2.3 2 / 2

✓ - 0 pts Correct (24C4)

- 1 pts Close

- 2 pts incorrect

2.4 2 / 2

✓ - 0 pts Correct

- 1 pts Close (Three of four)

- 2 pts Incorrect

2.5 1 / 2

- 0 pts Correct $(2^{(n^2 - n)} + 2^{(n^2 + n / 2)} - 2^{(n^2 - n / 2)})$

✓ - 1 pts Close

- 2 pts Incorrect

QUESTION 3

Equivalence relation 10 pts

3.1 it is an equivalence relation 4 / 4

✓ - 0 pts Correct

- 1 pts issue in transitivity

- 3 pts misunderstanding of what relation is saying

- 4 pts blank

- 2 pts misunderstanding of symmetry

- 1 pts the decimal thing isn't exactly right, e.g. -3 is related to $.7$

- 0 pts Click here to replace this description.

- 1 pts issue with symmetry

3.2 defining a function 4 / 4

✓ - 0 pts Correct

- 4 pts blank

- 2 pts need to prove uniqueness part of function

- 2 pts missing existence part of function

- 1 pts issue with uniqueness part of function

- **1 pts** need to consider different elements in the same equivalence class

- **1 pts** thing with decimals isn't quite right, for example $-.3$ and $.7$ are related

- **3 pts** big misunderstanding of the equivalence relation or function

3.3 a function that doesn't descend 0 / 2

- **0 pts** Correct

✓ - **2 pts** your g is not a function

- **1 pts** issue with justification

- **1 pts** your g does not work

- **2 pts** blank

QUESTION 4

m-ary tree 10 pts

4.1 number of internal vertices 3 / 5

- **0 pts** Correct

- **1 pts** No/incorrect answer

- **4 pts** No/incorrect justification

- **2 pts** Didn't justify number of total vertices

- **3 pts** "Proof by example"

- **2 pts** Assumed every terminal vertex had the same height as the tree

- **5 pts** Nothing

- **1 pts** Forgot to account for root

- **2 pts** Didn't subtract off internal vertices

- **2 Point adjustment**

☞ Correct final answer but there are i internal vertices; you forgot to count the root vertex which is not a child, and your two mistakes cancelled out.

4.2 height 4 / 5

- **0 pts** Correct

- **1 pts** No base case

- **1 pts** Didn't set up/invoke induction

- **1 pts** Backwards inductive step (didn't show inductive construction is exhaustive)

- **2 pts** Compared to complete tree without showing

this case is extremal

- **3 pts** Assumed tree is complete / inductive construction forms complete trees from complete trees

✓ - **1 pts** Assumed all immediate subtrees have height $h-1$

- **4 pts** "Proof by example"

- **5 pts** Nothing shown / Incorrect reasoning

QUESTION 5

spanning trees 10 pts

5.1 unique mst 5 / 6

- **0 pts** Correct

- **3 pts** Appeal to Prim's or Kruskal's Algorithm (without proving it can generate any MST)

- **6 pts** No / Invalid reasoning

- **1 Point adjustment**

☞ You need to choose e to be the edge of smallest weight occurring in any cycle in G . Otherwise very well done!

5.2 non unique spanning tree 4 / 4

✓ - **0 pts** Correct

- **4 pts** Not an example

- **4 pts** Claimed no such graph exists

- **4 pts** Nothing

QUESTION 6

planar graphs 10 pts

6.1 $2e > 3f$ 3 / 3

✓ + **3 pts** Correct

+ **2 pts** ≥ 3 edges for each face

+ **1 pts** ≥ 3 edges for each face (w/ mistake)

+ **1 pts** ≤ 2 faces for each edge

+ **0 pts** Incorrect

6.2 $e < 3v - 6$ 3 / 3

✓ + **3 pts** Correct

+ **2 pts** Euler's formula

+ 1 pts Correct application with (a)

+ 0 pts Incorrect

6.3 nonplanar graph 4 / 4

✓ + 4 pts Correct

+ 3 pts Isomorphic to $K_{3,3}$

+ 2 pts Mistaken/missing isomorphism to $K_{3,3}$

+ 1 pts $E \leq 2v-4$ or $2E \geq 4F$

+ 1 pts Other partial credit

+ 0 pts Incorrect

QUESTION 7

10 pts

7.1 7^{n-1} divisible by 6 5 / 5

✓ + 5 pts Correct

+ 1 pts Base case

+ 1 pts Inductive hypothesis

+ 2 pts factoring out a 7 in inductive step as $(6+1)$ or adding/subtracting 7

+ 1 pts Conclusion

+ 0 pts Incorrect

7.2 number with only 1s divisible by 7 5 / 5

✓ + 5 pts Correct

+ 0 pts Click here to replace this description.

+ 1 pts Look at 8 consecutive terms

+ 1 pts Pigeonhole remainder

+ 1 pts 7 divides a number of the form $111\dots000\dots$

+ 2 pts This implies that 7 divides $10^k - 1$

+ 1 pts Unsuccessful attempt with substantial work

QUESTION 8

balanced binary trees 10 pts

8.1 4 / 4

✓ - 0 pts Correct

- 2 pts incomplete, need to describe how a height n minimal balanced binary tree is made out of ones of smaller height

- 3 pts can't just do examples

- 4 pts blank

- 1 pts how are you adding in these trees/ vertices?

- 3 pts can't do induction without using some properties of minimal balanced binary trees

- 4 pts incorrect numbers/ equation

8.2 relationship to fibonacci numbers 3 / 3

✓ - 0 pts Correct

- 1.5 pts that is not the recurrence/ equation for the fibonacci numbers/ minimal balanced binary trees

- 1 pts you are assuming the desired conclusion

- 3 pts blank

- 1.5 pts need to use recurrence for fibonacci numbers

- 1.5 pts missing inductive step

- 1 pts the two recurrences aren't exactly the same, you need to account for this difference

- 0.5 pts error in equations

- 1 pts need to check initial conditions

8.3 Theta 3 / 3

✓ - 0 pts Correct

- 0.5 pts need to account for other term in equation for fibonacci numbers (sometimes it is contributing something positive, something something negative)

- 2 pts wrong formula for fibonacci numbers/ v_n

- 1 pts issue with big O

- 1 pts issue with omega

- 3 pts blank/ no gradable work

- 1 pts wrong equations/ issues with constants

- 2 pts need to use equation for v_n / Fibonacci numbers

QUESTION 9

binomial coefficients 10 pts

9.1 3^n 4 / 4

✓ + 4 pts Correct

+ 3 pts Minor error

+ 2 pts Binomial theorem

+ 1 pts Attempted induction or counting argument

+ 0 pts Incorrect

9.2 vandermonde identity 6 / 6

✓ + 6 pts Correct

+ 5 pts Minor error

+ 3 pts One part of counting argument or $(x+y)^{n+m}$

+ 1 pts Attempted to use induction/binomial

thrm/Pascal's identity

+ 0 pts Incorrect

Final

Name: _____

Student ID: _____

Section:

Tuesday:

Thursday:

1A

1B

TA: Albert Zheng

1C

1D

TA: Benjamin Spitz

1E

1F

TA: Eilon Reisin-Tzur

Instructions: Do not open this exam until instructed to do so. Please print your name and student ID number above, and circle the number of your discussion section. **You may not use calculators, books, notes, or any other material to help you.** Please make sure **your phone is silenced** and stowed where you cannot see it. Remember that you are bound by a conduct code.

Please get out your id and be ready to show it during the exam.

Please do not write below this line.

Question	Points	Score
1	10	
2	10	
3	10	
4	10	
5	10	
6	10	
7	10	
8	10	
9	10	
Total:	90	

1. (10 points) Circle the correct answer (only one answer is correct for each question)

$$1. \frac{n!}{k!(n-k)!} + \frac{n!}{(k+1)!(n-k-1)!} = \binom{n}{k} + \binom{n}{k+1} = \binom{n+1}{k+1}$$

(a) $\frac{(n+k)!}{k!n!}$

(b) $\frac{(n+1)!}{k!(n+1-k)!}$

(c) $\frac{(n+1)!}{(k+1)!(n-k)!}$

(d) none of the above

$$\frac{(n+1)!}{(k+1)!(n-k-1)!}$$

2. The decision tree of a sorting algorithm for sorting n items (where at each step we can only decide whether or not one item is less than other) necessarily has:

(a) a height of $\geq \lg(n!)$

(b) a height of $\Omega \lg(n!)$ (but not necessarily a height of $\geq \lg(n!)$)

(c) a height of $O(\lg(n!))$

(d) a height of $O(n \lg n)$

3. If G is a graph with n vertices and $n - 2$ edges, then:

~~(a)~~ G is a tree

(b) G is connected

(c) G is disconnected

(d) G is simple



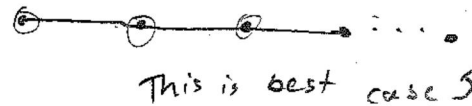
Question 1 continued...

4. Which of these graphs has an Euler cycle?

- (a) K_4
- (b) K_5
- (c) $K_{3,3}$
- (d) $K_{2,3}$

5. What is the fewest number of edges (i.e. in the best case) that could be examined by Dijkstra's algorithm on a graph with n vertices? (We examine edges in the part of the algorithm where we update labels.)
Your answer should be true for all n .

- (a) Less than or equal to n
- (b) More than n but less than or equal to $n^2/2$
- (c) More than $n^2/2$ but less than or equal to n^2
- (d) More than n^2



2. In this question write down your answer, no need for any justification. Leave your answers in a form involving factorials, $P(n, m)$, $\binom{n}{m}$, exponents, etc.

(a) (2 points) If $s_n = s_{n-1} + 6s_{n-2}$ and $s_0 = 2, s_1 = 1$, what is s_{100} ?

$$\begin{aligned} t^2 - t - 6 &= 0 & t &= 3, -2 & s_n &= A3^n + B(-2)^n & A+B &= 2 \\ (t-3)(t+2) &= 0 & & & & & 3A-2B &= 1 & B &= 1 \\ & & & & & & 5A &= 5 & & \Rightarrow A=1 \end{aligned}$$

$$s_{100} = 3^{100} + (-2)^{100}$$

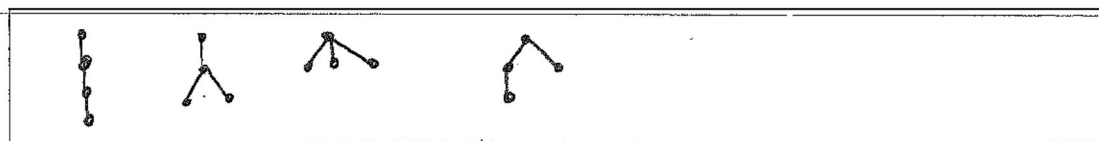
(b) (2 points) How many ways can 7 distinct math majors and 4 distinct CS majors sit in a circle, if the CS majors won't sit by each other and we say that two seatings are the same if they are related by a rotation?

$$\frac{7! P(8, 4)}{11!}$$

(c) (2 points) A squirrel has 20 identical acorns that she is going to hide among 5 distinct holes. In how many ways can the squirrel hide the acorns?

$$\binom{20+5-1}{5-1} = \binom{24}{4}$$

(d) (2 points) Draw all the distinct (up to isomorphism) rooted trees with 4 vertices. Please put the root at the top.



(e) (2 points) What is the number of relations that are symmetric or reflexive on a set with n -elements?

$$2^{\binom{n}{2}} = \# \text{ of symmetric} \quad 2^{P(n,2)} = \# \text{ of reflexive}$$

$$2^{\binom{n}{2}} + 2^{P(n,2)}$$

symmetric: choose 2 elements a, b , include $(a, b), (b, a)$ $2^{\binom{n}{2}}$
 reflexive: include the n elements $(a, a), \dots$ then choose other elements $2^{P(n,2)}$

3. Consider the relation on the real numbers defined by $C = \{(x, y) \in \mathbb{R} \times \mathbb{R} : x - y \in \mathbb{Z}\}$.

(a) (4 points) Show that C is an equivalence relation.

• For all $x \in \mathbb{R}$, $x - x = 0$ which is an integer so $x - x \in \mathbb{Z}$ and $x C x$. Thus, C is reflexive.

• If $(x, y) \in C$, then $(x - y) \in \mathbb{Z}$ and $-(x - y) \in \mathbb{Z}$, so $y - x \in \mathbb{Z}$ and $(y, x) \in C$. Thus, for any $(x, y) \in C$, $(y, x) \in C$, so C is symmetric.

• If $(x, y), (y, z) \in C$, then we know that $x - y \in \mathbb{Z}$, and $y - z \in \mathbb{Z}$. The sum of two integers is also always an integer, so $(x - y) + (y - z) = x - z \in \mathbb{Z}$. Therefore $x C z$ and C is transitive.

• C is a reflexive, symmetric, and transitive relation, so it is an equivalence relation.

- (b) (4 points) Let $\tilde{\mathbb{R}}$ denote the set of equivalence classes of C , i.e. $\tilde{\mathbb{R}} = \{[x] : x \in \mathbb{R}\}$. Consider the function $f : \mathbb{R} \rightarrow \mathbb{R}$ defined by $f(x) = x + 1/2$.

Show that the relation \tilde{f} from $\tilde{\mathbb{R}}$ to $\tilde{\mathbb{R}}$ defined by $\tilde{f} = \{([a], [b]) \in \tilde{\mathbb{R}} \times \tilde{\mathbb{R}} : f(a) = b\}$ is a function.

For any $x_1, x_2 \in \tilde{\mathbb{R}}$, if $[x_1] = [x_2]$, then $\tilde{f}([x_1]) = [x_1 + \frac{1}{2}]$ and $\tilde{f}([x_2]) = [x_2 + \frac{1}{2}]$. We know that $x_1 + \frac{1}{2} - (x_2 + \frac{1}{2}) = 0 \in \mathbb{Z}$, so $[x_1 + \frac{1}{2}] = [x_2 + \frac{1}{2}]$. Also every $x \in \mathbb{R}$ has an equivalence class because C is an equivalence relation, so \tilde{f} is a function. The reverse is also true: if $\tilde{f}([x_1]) \neq \tilde{f}([x_2])$ then we know that $\tilde{f}([x_1]) - \tilde{f}([x_2]) \in \mathbb{Z}$ so $[x_1] \neq [x_2]$ because otherwise \uparrow would not be true.

- (c) (2 points) Give an example of a function $g : \mathbb{R} \rightarrow \mathbb{R}$ so that the relation \tilde{g} from $\tilde{\mathbb{R}}$ to $\tilde{\mathbb{R}}$ defined by $\tilde{g} = \{([a], [b]) \in \tilde{\mathbb{R}} \times \tilde{\mathbb{R}} : g(a) = b\}$ is **not** a function. (Be sure to justify your answer.)

$$g(x) = \frac{1}{x}$$

$$[\frac{1}{2}] = \{\frac{1}{2}, \frac{3}{2}, \frac{5}{2}, \dots\}$$

$$\tilde{g}([\frac{1}{2}]) = [2]$$

$$\tilde{g}([\frac{5}{2}]) = [\frac{2}{5}]$$

$$2 - \frac{2}{5} \notin \mathbb{Z}, \text{ so } [2] \neq [\frac{2}{5}] \text{ and}$$

\tilde{g} is not a function.

4. For m , a positive integer, a *full m -ary tree* is a rooted tree where every parent has exactly m children.

(a) (5 points) If T is a full m -ary tree with i internal vertices, how many terminal vertices does T have?



4 + 5



1 + 2



27

each internal vertex has m children, so $m \cdot i$ children, where $(i-1)$ of those are internal vertices, so subtracting internal vertices $(i-1)$ from children

(m.i): $m \cdot i - (i-1) = i(m-1) + 1$

This holds for the binary tree case as $t = i + 1$

(b) (5 points) Show that if T is a full m -ary tree of height h with t terminal vertices, then $t \leq m^h$.

This can be shown by induction on h .

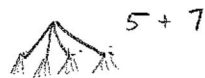
Base Case: For a height 0 m -ary tree, there will be exactly 1 vertex, so $t \leq m^0 = 1$ is true.

Inductive Step: Assume the number of terminal vertices of any height h m -ary tree is $t \leq m^h$. We must show for a height $h+1$ m -ary tree, $t \leq m^{h+1}$.

• The number of terminal vertices is just the sum of all the terminal vertices of the individual height h child subtrees of the root. Thus, by the inductive hypothesis: $t \leq m \cdot m^h = m^{h+1}$ so $t \leq m^{h+1}$ for all $h \geq 0$ by mathematical induction. \square



4 + 5



5

16

5. (a) (6 points) Show that if G is a connected weighted graph where all the edges of G have distinct weights then G has a unique minimal spanning tree.

Base case: If G is a graph with 1 edge, then the graph is either O^w , \xrightarrow{w} and each have a unique minimal spanning tree, so the base case holds.

Inductive step: Assume all connected weighted graphs G with n edges with unique weights have a unique minimal spanning tree, we must show a graph with $n+1$ edges also has a unique minimal spanning tree. The graph G must have an edge incident on a vertex of degree 1 like so:



or have a cycle:



In the first case, remove edge e and then use the inductive hypothesis to construct a minimal spanning tree of the n edge graph. Then add edge e to the tree and there is now a unique minimal spanning tree.

For the second case, simply remove e and use the inductive hypothesis to construct the unique minimal spanning tree.

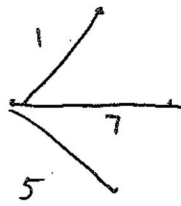
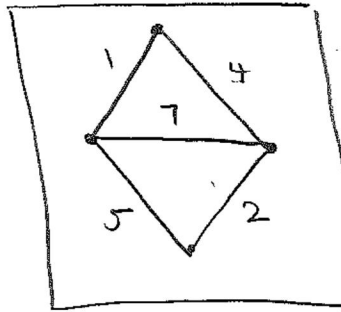
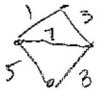
We have proved that a connected weighted graph G has a unique minimal spanning tree by the principle of mathematical induction. \square

- (b) (4 points) Give an example of a connected weighted graph G so that all the edges of G have distinct weights and G has at least two distinct spanning trees that have the same total weight, i.e. the sums of the weights of the edges in these two distinct trees agree, or prove that no such weighted graph exists.

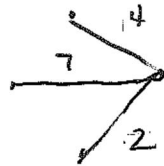
1 5

3

3



$$1 + 5 + 7 = 13$$



$$4 + 2 + 7 = 13$$

6. (a) (3 points) Show that for G a connected simple planar graph containing a cycle if G has E edges and F faces, then $2E \geq 3F$.

Count number of boundary edges:

• each edge is a boundary to at most two faces: $\# \text{ boundaries} \leq 2E$

• Each face has at least 3 boundary edges
 $\# \text{ boundaries} \geq 3F$

$$3F \leq \# \text{ boundaries} \leq 2E$$

$$3F \leq 2E$$

- (b) (3 points) Show that for G a connected simple planar graph containing a cycle if G has E edges and V vertices, then $E \leq 3V - 6$.

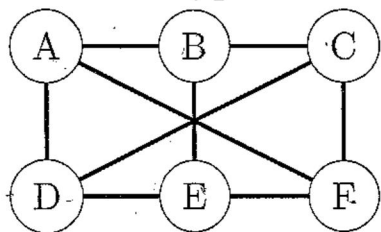
From part (a), $3F \leq 2E$. Euler equation
is $F - E + V = 2$, so $F = E - V + 2$

$$3(E - V + 2) \leq 2E$$

$$3E - 3V + 6 \leq 2E$$

$$E \leq 3V - 6$$

(c) (4 points) Is the following graph planar? If it is give a planar drawing of it. If not, prove that it is not planar.



$$E = 9 \quad V = 6$$

$$F - E + V = 2$$

$$F = 9 - 6 + 2 = 5$$

• Every cycle in this graph has length at least 4, so this implies that we should have

$$2E \geq 4F$$

$$2(9) \geq 4(5)$$

$18 \not\geq 20$, so it is not planar.

7. (a) (5 points) Show that for all $n \geq 1$, $7^n - 1$ is divisible by 6.

Base case: If $n=1$, $7^1 - 1 = 6$ which is clearly divisible by 6, so the base case holds.

Inductive Step: Assume $7^n - 1$ is divisible by 6 for all n . We must show $7^{n+1} - 1$ is also divisible by 6.

$$7^{n+1} - 1 = 7 \cdot 7^n - 1 = 6 \cdot 7^n + (7^n - 1)$$

From the inductive hypothesis we know that $7^n - 1$ is divisible by 6. $6 \cdot 7^n$ is also clearly divisible by 6 because it has a factor of 6. The sum of two numbers divisible by 6 is also divisible by 6, so we have shown $7^{n+1} - 1$ is divisible by 6 for all $n \geq 1$ by mathematical induction. \square

- (b) (5 points) Show that there is a number of the form $\sum_{i=0}^n 10^i$ (i.e. a number consisting only of 1s) that is divisible by 7.

There are 7 possible remainders when dividing by 7, so considering the 8 numbers
1, 11, 111, 1111, 11111, 111111, 1111111, 11111111

by the pigeonhole principle, 2 distinct numbers from above have the same remainder when divided by 7.

By subtracting these two numbers, $x_1 - x_2$ is divisible by 7. $x_1 - x_2$ is of the form

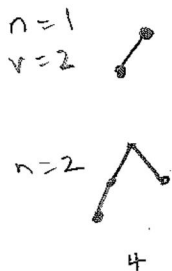
$$\underbrace{111 \dots 1}_{n-m} \cdot 10^m$$

if we ensure $x_1 > x_2$ and m is the number of digits in the smaller number x_2 .

Thus we have a number $\underbrace{11 \dots 1}_{n-m}$ divisible by 7. which is of the form $\sum_{i=0}^{n-m} 10^i$

8. A *balanced binary tree* is a binary tree where for each vertex the heights of the left and right subtrees of that vertex differ by at most one. Let v_n denote the minimum number of vertices in a balanced binary tree of height n .

(a) (4 points) Show that v_n satisfies for $n \geq 2$ the recurrence $v_n = v_{n-1} + v_{n-2} + 1$



$\xleftrightarrow{\text{swap}}$

height $n-1$ height $n-2$

$$v_n = \text{root vertex} + \binom{\# \text{ of vertices of right subtree}}{\# \text{ of vertices of left subtree}}$$

$$= 1 + \binom{\text{min vertices for height } n-1}{\text{min vertices for height } n-2}$$

$$v_n = 1 + v_{n-1} + v_{n-2}$$

(b) (3 points) Show that for $n \geq 0$, $v_n = \frac{F_{n+2}}{F_{n+3}-1}$, where F_k is the k^{th} Fibonacci number. 1, 1, 2, 3, 5

$n=0 : v_0 = 1, F_3 - 1 = 2 - 1 = 1 \quad \checkmark$
 $n=1 : v_1 = 2, F_4 - 1 = 3 - 1 = 2 \quad \checkmark$
 The base cases hold.

Inductive Step: Assume $v_k = F_{k+3} - 1$ for all $k \leq n$. We must show $v_{n+1} = F_{n+4} - 1$.

$$v_{n+1} = 1 + v_n + v_{n-1} \quad \leftarrow \text{inductive hypothesis.}$$

$$= 1 + F_{n+3} - 1 + F_{n+2} - 1 = F_{n+4} - 1$$

So $v_n = F_{n+3} - 1$ for all $n \geq 0$ by the principle of mathematical induction. \square

(c) (3 points) Show that $v_n = \Theta(\phi^{n+2})$, where $\phi = \frac{1+\sqrt{5}}{2}$.

$$v_n = \left(\frac{1}{\sqrt{5}}\right) \left(\frac{1+\sqrt{5}}{2}\right)^{n+3} - \left(\frac{1}{\sqrt{5}}\right) \left(\frac{1-\sqrt{5}}{2}\right)^{n+3} - 1$$

$$\geq \left(\frac{1}{\sqrt{5}}\right) \left(\frac{1+\sqrt{5}}{2}\right)^{n+3} - \frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2}\right)^{n+3} - \left(\frac{1+\sqrt{5}}{2}\right)^{n+3}$$

$$= \left(-\frac{2}{\sqrt{5}} - 1\right) \left(\frac{1+\sqrt{5}}{2}\right)^{n+3}$$

$$\geq$$

because $\left(\frac{1+\sqrt{5}}{2}\right) > 1$
 because $\left(\frac{1+\sqrt{5}}{2}\right)^n$ is an increasing function

Thus

$$v_n = A \left(\frac{1+\sqrt{5}}{2}\right)^{n+3} + B \left(\frac{1-\sqrt{5}}{2}\right)^{n+3} - 1$$

$$\leq$$

Let $\psi = \frac{1-\sqrt{5}}{2}$

$$\rightarrow v_n = F_{n+3} - 1 = F_{n+2} + F_{n+1} - 1$$

$$= \left(\frac{1}{\sqrt{5}}\right) \phi^{n+2} - \left(\frac{1}{\sqrt{5}}\right) \psi^{n+2} + \left(\frac{1}{\sqrt{5}}\right) \phi^{n+1} - \left(\frac{1}{\sqrt{5}}\right) \psi^{n+1} - 1$$

This is an exponential function where $\phi > \psi$ and ϕ^n is increasing, so the greatest power is ϕ^{n+2} , so $v_n = \Theta(\phi^{n+2})$

$$A \left(\frac{1+\sqrt{5}}{2}\right) + B \left(\frac{1+\sqrt{5}}{2}\right) = 1 \quad A = \frac{1 - B \left(\frac{1-\sqrt{5}}{2}\right)}{\left(\frac{1+\sqrt{5}}{2}\right)}$$

$$A \left(\frac{6+2\sqrt{5}}{4}\right) + B \left(\frac{6-2\sqrt{5}}{4}\right) = 1$$

$$\left(\frac{1 - B \left(\frac{1-\sqrt{5}}{2}\right)}{\left(\frac{1+\sqrt{5}}{2}\right)}\right) \left(\frac{6+2\sqrt{5}}{4}\right) + B \left(\frac{3-\sqrt{5}}{2}\right) = 1$$

$$\left(1 - B \left(\frac{1-\sqrt{5}}{2}\right)\right) \left(\frac{3+\sqrt{5}}{1+\sqrt{5}}\right) + B \left(\frac{3-\sqrt{5}}{2}\right) = 1$$

9. (a) (4 points) Show that $\sum_{i=0}^n 2^i \binom{n}{i} = 3^n$.

Binomial theorem:

$$\sum_{i=0}^n 2^i \binom{n}{i} = \sum_{i=0}^n \binom{n}{i} 2^i 1^{n-i} = (2+1)^n = \boxed{3^n}$$

(b) (6 points) Show that $\binom{n+m}{r} = \sum_{i=0}^r \binom{n}{i} \binom{m}{r-i}$.

Choosing r items out of a group of n items and a group of m items is $\binom{n+m}{r}$. We can also choose i items from group of n , then the remaining $(r-i)$ items from the group of m , so total for i items from group of n is $\binom{n}{i} \binom{m}{r-i}$. We want all possible numbers from group of n , so $\binom{n+m}{r} = \sum_{i=0}^r \binom{n}{i} \binom{m}{r-i}$.

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$$\begin{aligned}
 &= \frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2} \right) - \frac{1}{\sqrt{5}} \left(\frac{1-\sqrt{5}}{2} \right) \\
 &= \frac{1+\sqrt{5}-1+\sqrt{5}}{2\sqrt{5}} = 1 \checkmark
 \end{aligned}$$

$$\frac{3+\sqrt{5}}{1+\sqrt{5}} - B \left(\frac{3-5-2\sqrt{5}}{2+2\sqrt{5}} \right) + B \left(\frac{3-\sqrt{5}}{2} \right) = 1$$

$$\frac{3+\sqrt{5}-3\sqrt{5}-5}{1-5} - B(-1) + B \left(\frac{3-\sqrt{5}}{2} \right) = 1$$

$$\frac{-2-2\sqrt{5}}{-4} + B \left(1 + \frac{3-\sqrt{5}}{2} \right) = 1$$

$$\begin{aligned}
 B \left(\frac{5-\sqrt{5}}{2} \right) &= -\frac{1+\sqrt{5}}{2} & B &= \frac{-1-\sqrt{5}}{5-\sqrt{5}} \\
 & & &= \frac{-5-5\sqrt{5}-\sqrt{5}-5}{25-5} \\
 & & &= \frac{-10-6\sqrt{5}}{20} =
 \end{aligned}$$