

$60/69^{13} = 63/69 = 91\%$ ~~XXXXXXXXXX~~

10/30/17

9AM DIS

Midterm exam

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LING 120C: Semantics I

Note: You may use the textbook (Heim & Kratzer 1998) and your class notes on this midterm exam.

1

4 points

-2

Please complete the following:

(a) The ^{ex} ~~intension~~ of a sentence is truth value the thing the expression refers to at a given world & time

~~60/69~~ ²/₄

(b) The ⁱⁿ ~~extension~~ of a sentence is truth condition the general concept behind the expression that determines what the expression's extension is at a given world & time

Donald Trump
(currently)

2

5 points

Please briefly describe how our semantic system based on **truth values** is able to derive **truth conditions**.

-2

To know the meaning of a sentence is to know its truth conditions. And, as sentences are composed of discrete parts — each of which denote truth values, the conditions of the sentence as a whole depends on the functional application of the individual truth values satisfying one another in the sentence to compose its meaning.

assume [S] = T and work backwards

3

10 points

Please answer the following comprehension questions about functions.

- (a) What is the domain and the range of the following function? What is $f(LA)$?

$$f = \{ \langle NYC, de\ Blasio \rangle, \langle LA, Garcetti \rangle, \langle Chicago, Emanuel \rangle \}$$

Domain $\{ NYC, LA, Chicago \}$

Range $\{ de\ Blasio, Garcetti, Emanuel \}$

$$f(LA) = \langle LA, Garcetti \rangle$$

- (b) Please define the following function using λ -notation:

$$g = \{ \langle 1, 1 \rangle, \langle 2, 4 \rangle, \langle 3, 9 \rangle, \langle 4, 16 \rangle, \langle 5, 25 \rangle \}$$

~~$x \in \mathbb{N}$~~

$$[\lambda x: x \in \mathbb{N}. x^2]$$

1-5

- (c) Please use λ -conversion to calculate the value of the following expression (show each step):

$$[\lambda x: x \in \mathbb{N}. [\lambda z: z \in \mathbb{N}. x + z]](4)(5)$$

(remember that \mathbb{N} is the set of natural numbers)

$$[\lambda z: z \in \mathbb{N}. 4 + z](5)$$

$$= 4 + 5 = 9$$

- (d) What is the semantic type of the following function?

$$\lambda x: x \in D_e . \lambda y: y \in D_e . \lambda z: z \in D_e . z \text{ introduces } x \text{ to } y$$

$$\langle e, \langle e, \langle e, + \rangle \rangle \rangle$$

✓ Please define the following function using λ -notation:

$g : D_e \rightarrow D_{(e,t)}$
 for every $x \in D_e$, $g(x) = i_x : D_e \rightarrow D_t$
 for every $y \in D_e$,
 $i_x(y) = T$ iff y hits x

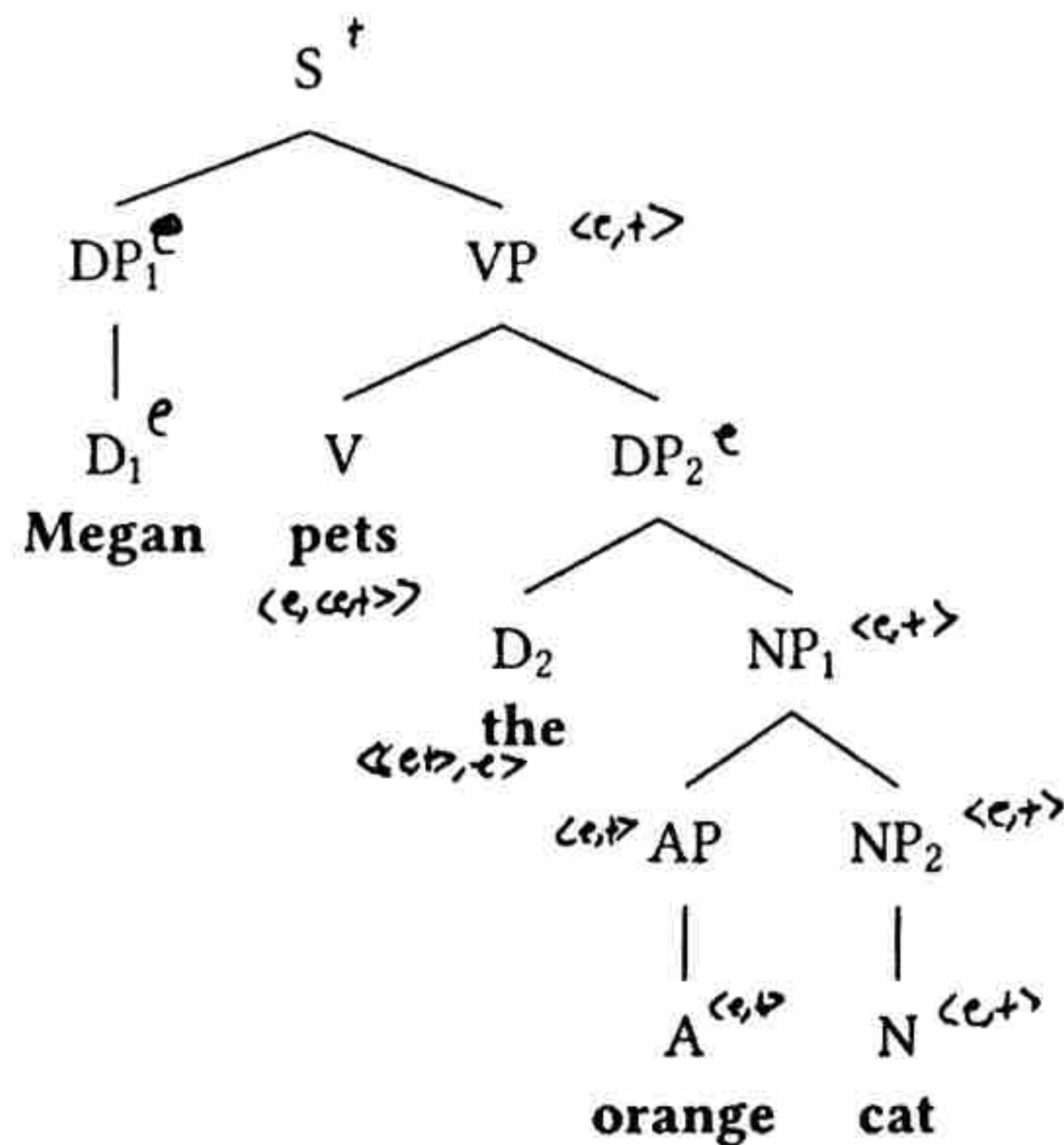
$[\lambda x_e . [\lambda y_e . y \text{ hits } x]]$

4

30 points

On the next sheet, please use our semantic-composition rules (listed on the last page) to compute the truth conditions of the following sentence:

-1



Be sure to show each step of the derivation and to indicate the semantic-composition rule or other deduction (e.g. λ -conversion or a subproof) used at each step.

SUBPROOF 1
 $\llbracket DP_1 \rrbracket = NN$
 $\llbracket D_1 \rrbracket = TN$

★ Megan

SUBPROOF 2
 $\llbracket AP \rrbracket = NN$
 $\llbracket A \rrbracket = TN$

★ $\llbracket \lambda x_e. x \text{ is orange} \rrbracket$

SUBPROOF 3
 $\llbracket NP_2 \rrbracket = NN$
 $\llbracket N_2 \rrbracket = TN$

★ $\llbracket \lambda x_e. x \text{ is a cat} \rrbracket$

SUBPROOF 4
 $\llbracket NP_1 \rrbracket = PM$

$\llbracket \lambda x_e. \llbracket AP \rrbracket(x) = T \text{ and } \llbracket NP_2 \rrbracket(x) = T \rrbracket$
 = by λ comm

$\llbracket \lambda x_e. \llbracket \lambda x_e. x \text{ is orange} \rrbracket(x) = T \text{ and } \llbracket \lambda x_e. x \text{ is a cat} \rrbracket(x) = T \rrbracket$
 = by λ comm

★ $\llbracket \lambda x_e. x \text{ is orange and } x \text{ is a cat} \rrbracket$

SUBPROOF 5

$\llbracket D_2 \rrbracket = TN$

★ $\llbracket \lambda f_{cat} \lambda t. \exists! x [f(x) = T] \cdot \forall y [f(y) = T] \rrbracket$

SUBPROOF 6

$\llbracket DP_2 \rrbracket = FA$

$\llbracket P_2 \rrbracket(\llbracket NP_1 \rrbracket) = 5, 4$

$\llbracket \lambda f_{cat} \lambda t. \exists! x [f(x) = T] \cdot \forall y [f(y) = T] \rrbracket (\lambda x_e. x \text{ is orange} \ \& \ x \text{ is a cat}) = \lambda \text{ comm}$

$\forall y [\llbracket \lambda x_e. x \text{ is orange} \ \& \ x \text{ is a cat} \rrbracket(y) = T]$, where $\exists! x [f(x) = T]$ = $\lambda \text{ comm}$
 $x \text{ is orange and } x \text{ is a cat} - 1$

★ $\forall y [y \text{ is orange} \ \& \ y \text{ is a cat}]$, where $\exists! y [f(y) = T]$

SUBPROOF 7

$\llbracket V \rrbracket = TN$

$\llbracket \lambda x_e. (\lambda y_e. y \text{ pets } x) \rrbracket$

SUBPROOF 8

$\llbracket VP \rrbracket = FA$

$\llbracket V \rrbracket(\llbracket DP_2 \rrbracket) = 7, 6$

$\llbracket \lambda x_e. (\lambda y_e. y \text{ pets } x) \rrbracket (\llbracket \lambda z [z \text{ is orange} \ \& \ z \text{ is a cat}] \text{, where } \exists! z [z \text{ is orange} \ \& \ z \text{ is a cat}] \rrbracket)$

★ $\llbracket \lambda y_e. y \text{ pets } [\llbracket \lambda z [z \text{ is orange} \ \& \ z \text{ is a cat}] \text{, where } \exists! z [z \text{ is orange} \ \& \ z \text{ is a cat}] \rrbracket] \rrbracket$

$\llbracket S \rrbracket = T$ iff by FA

$\llbracket VP \rrbracket(\llbracket DP_1 \rrbracket) = T$ iff by 1, 8

$\llbracket \lambda y_e. y \text{ pets } [\llbracket \lambda z [z \text{ is orange} \ \& \ z \text{ is a cat}] \text{, where } \exists! z [z \text{ is orange} \ \& \ z \text{ is a cat}] \rrbracket] \rrbracket$ (Megan) = T iff by λ -comm

★ Megan pets $\llbracket \lambda z [z \text{ is orange} \ \& \ z \text{ is a cat}] \text{, where } \exists! z [z \text{ is orange} \ \& \ z \text{ is a cat}] \rrbracket$

★ QED □

because I'm fancy

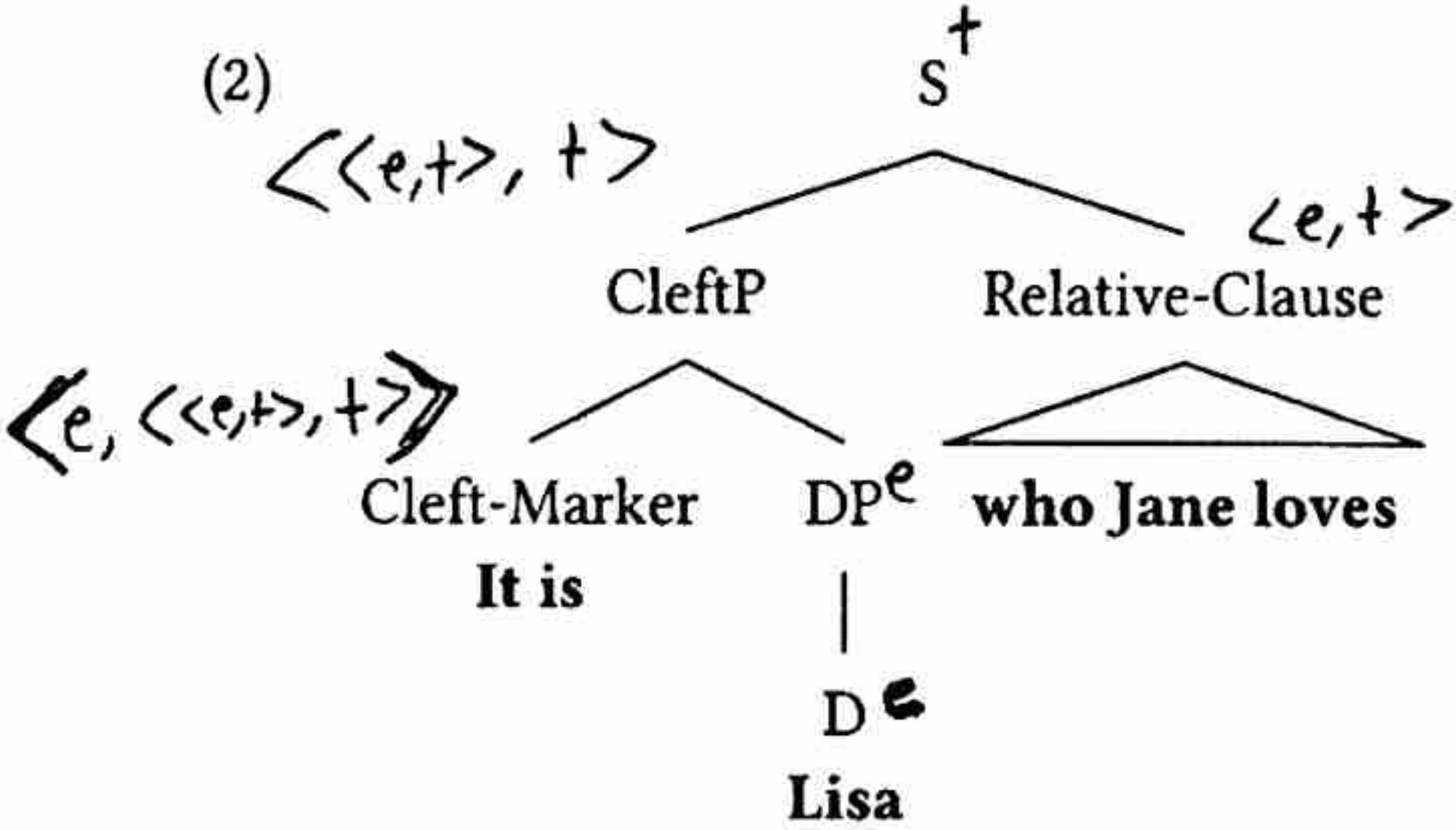
5

20 points

The English cleft construction is exemplified by sentences like those below:

- (1) a. It is Lisa who Jane loves.
- b. It is Lisa who loves Jane.
- c. It is Ellen who Jane loves.
- d. It is Ellen who loves Jane.

We will assume that this construction has the following, oversimplified syntax:



In this exercise, you will be working out a compositional semantics for the English cleft construction. To begin with, let us assume the lexical entries below:

- (3) a. $[[\text{Lisa}]] = \text{Lisa}$ e
- b. $[[\text{Ellen}]] = \text{Ellen}$ e
- c. $[[\text{who Jane loves}]] = \lambda x_e . \text{Jane loves } x$
- d. $[[\text{who loves Jane}]] = \lambda x_e . x \text{ loves Jane}$

$\langle e, t \rangle$
 $\langle e, t \rangle$

Question 1:

[5 points]

Assuming the syntax in (2) and the lexical entries in (3), what could the semantic type of the Cleft-Marker *it is* be? [Note: There are two possibilities.]

$\langle e, \langle \langle e, t \rangle, t \rangle \rangle$

Linguists have long observed that the English cleft carries a particular *presupposition*. As indicated below, a cleft of the form **It is DP Relative-Clause** presupposes that there is exactly one thing that the relative clause is true of. For example:

PRE
SUPPOSITIONS

- (4) a. **It is Lisa who Jane loves** presupposes there is exactly one x such that Jane loves x .
 $\exists! x [\text{Jane loves } x]$
- b. **It is Ellen who loves Jane** presupposes there is exactly one x such that x loves Jane.
 $\exists! x [x \text{ loves Jane}]$

Question 2:

[5 points]

Use a diagnostic to show that the presuppositions in (4) are indeed correct.

one way to prove a presupposition's correctness is to check if presupposition failure (aka undefined extension) persists despite negating the original sentence ^{True neither T nor F}

under the presupposition, **It is DP Rel-Clause** can only take as argument ~~the~~ entity x that ~~is~~ unique in this world.

Thus, one example of ~~an~~ an input that would give an undefined extension is cats.

It is ~~cats~~ that Jane loves
 presupposes: there is exactly one x such that DP Jane loves x

It is not the case that it is ~~cats~~ that Jane loves.
 presupposes: there is exactly one x such that Rel Clause Jane loves x .

It's not the case that it's Lisa who Jane loves

It's Lisa who Jane loves
 = T iff Jane loves Lisa + presuppos

Indeed, the ambiguity of what "cat" being referred to persists in the keeping the truth value of the sentence as a whole unclear even after negation

Linguists have also long observed that, aside from its presupposition, the English cleft construction does not seem to differ much in meaning from the 'non-clefted' counterpart. That is, the following truth-conditional statements seem accurate:

- clefted (5) a. **It is Lisa who Jane loves** is T iff Lisa is the unique y such that Jane loves y .
- non-clefted b. **It is Ellen who loves Jane** is T iff Ellen is the unique y such that y loves Jane.

TRUTH
CONDITIONS

$\langle e, \langle \langle e, t \rangle, t \rangle \rangle$

~~1~~ - ~~4~~ 4

Question 3:

[10 points]

Provide a lexical entry for the Cleft-Marker **it is** that predicts the truth-conditional statements in (5) and the presuppositions in (4). (You only need to give the lexical entry. You do *not* have to provide a truth-conditional derivation.)

$\lambda f_{\langle e, t \rangle} : \exists! x [f(x) = T]$
 $\lambda x_e : \exists! x [\lambda f_{\langle e, t \rangle} . f(x) = T] . \lambda g_{\langle \langle e, t \rangle, t \rangle} . g(f)$
 $\lambda g_{\langle e, t \rangle} . [\lambda f_{\langle e, t \rangle} . f(x) = T] [\lambda g_{\langle e, t \rangle} . g(x)]$
 $\lambda x_e : \text{there is exactly one } x \text{ such that } [\lambda f_{\langle e, t \rangle} . f(x) = T] [\lambda g_{\langle e, t \rangle} . g(x)]$

presupp here w/ λf , otherwise f doesn't get bound by the correct meaning requirement
 $\langle \langle e, t \rangle, e \rangle$
 $\llbracket \text{the} \rrbracket = \lambda f_{\langle e, t \rangle} : \exists! x [f(x) = T]$
 • by $\llbracket f(y) = T \rrbracket$

$\lambda x_e \lambda f_{\langle e, t \rangle} : \exists! y [f(y) = T] \cdot f(x) = T$
 not type $\langle e, \langle \langle e, t \rangle, t \rangle \rangle$

Big Hint: Consider the way that our lexical entry for **the** captures the presupposition with which it is associated.

Bigger Hint: Your answer should analyze CleftP (i.e. **It is Lisa**) as being of semantic type $\langle \langle e, t \rangle, t \rangle$.

(6) **Semantic-composition rules**

a. **Nonbranching Nodes (NN)**

If α is a nonbranching node, and β is its daughter node, then $\llbracket \alpha \rrbracket = \llbracket \beta \rrbracket$.

b. **Functional Application (FA)**

If α is a branching node, $\{\beta, \gamma\}$ is the set of α 's daughters, and $\llbracket \beta \rrbracket$ is a function whose domain contains $\llbracket \gamma \rrbracket$, then $\llbracket \alpha \rrbracket = \llbracket \beta \rrbracket (\llbracket \gamma \rrbracket)$.

c. **Terminal Nodes (TN)**

If α is a terminal node, $\llbracket \alpha \rrbracket$ is specified in the lexicon.

d. **Predicate Modification (PM)**

If α is a branching node, $\{\beta, \gamma\}$ is the set of α 's daughters, and $\llbracket \beta \rrbracket$ and $\llbracket \gamma \rrbracket$ are both in $D_{\langle e, t \rangle}$, then $\llbracket \alpha \rrbracket = \lambda x \in D_e . \llbracket \beta \rrbracket (x) = T$ and $\llbracket \gamma \rrbracket (x) = T$.

"It is Lisa who Jane loves" = T iff
 Jane loves Lisa
 $\llbracket \text{it is} \rrbracket (\llbracket \text{DP} \rrbracket) (\llbracket \text{DP} \rrbracket)$
 = DP is the unique x such that Jane loves x