

# EEM16 Final

VAIBHAV AGGARWAL

TOTAL POINTS

**114.5 / 125**

QUESTION 1

Problem #1 30 pts

1.1 (a), (b) 6 / 6

✓ - 0 pts a is correct

✓ - 0 pts b is correct

1.2 (c) - (i) 22.5 / 24

✓ - 0 pts c - correct

✓ - 0 pts d - correct

✓ - 0.5 pts e - 3 product terms

✓ - 0.5 pts f - columns or rows

✓ - 0 pts g - correct

✓ - 0 pts h - correct

✓ - 0.5 pts i - 12 states

✓ - 0 pts Correct

QUESTION 7

7 Problem #7 12 / 15

✓ - 1 pts (e)incorrect setup

✓ - 1 pts (e)incorrect hold

✓ - 1 pts (e)incorrect c2q

QUESTION 2

2 Problem #2 15 / 15

✓ - 0 pts Correct

QUESTION 3

3 Problem #3 13 / 15

✓ - 2 pts missing one thing or has an extra state

QUESTION 4

4 Problem #4 20 / 20

✓ - 0 pts Correct

QUESTION 5

Problem #5 18 pts

5.1 (a)-(c) 7 / 11

✓ - 4 pts (c) incorrect

5.2 (d)-(f) 7 / 7

✓ - 0 pts Correct

QUESTION 6

6 Problem #6 12 / 12

# Final Exam

Name (Last, First): Aggarwal, Vaibhav

Student Id #:

Do not start working until instructed to do so.

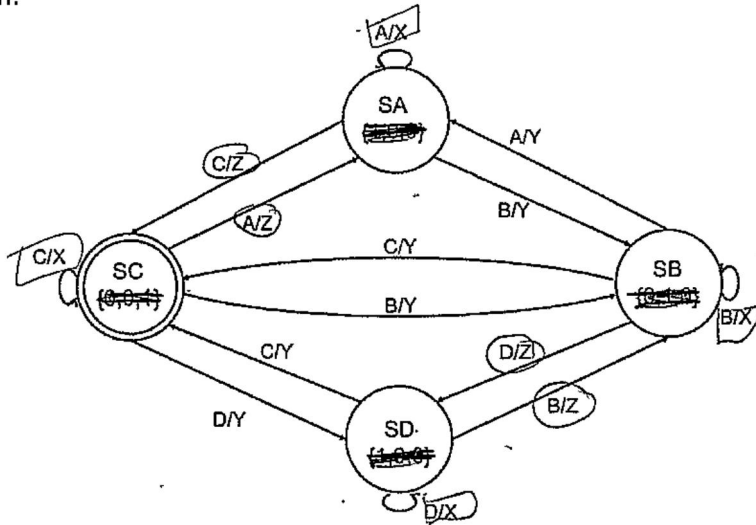
1. You must answer in the space provided for answers after every question. We will ignore answers written anywhere else in the booklet. All pages in this booklet must be accounted for otherwise it will not be graded.
2. You are permitted 2 page of notes 8.5x11 (front and back).
3. You may not use any electronic device.

Following table to be filled by course staff only

	Maximum Score	Your Score
Question 1	30	
Question 2	15	
Question 3	15	
Question 4	20	
Question 5	18	
Question 6	12	
Question 7	15	
TOTAL	125	

Question #1 State Machine to Logic (30)

A Mealy FSM state diagram is shown below. This is a decoder for a 3-level to 4-level encoding. A 4-level signal is communicated between two endpoints (A, B, C, and D). This signal is the input to the FSM. Transitions between the levels map to 3 symbols (X, Y, and Z); these symbols are the outputs. Note that some transitions are eliminated to enhance the quality of the communication.



(a) (2) Explain the difference between a Mealy and a Moore FSM.

*In a Mealy FSM, the output depends on the current state and the input, while in a Moore FSM it only depends on the current state.*

(b) (4) Fill in the blanks in this partial state transition table.

state	in	out	nx_state
SA	A	X	SA
SB	B	X	SB
SD	B	Z	SB
SC	C	X	SC
SDVSB	C	Y	SC
SC	A	Z	SA
SB	D	Z	SD

Assume for the following parts that inputs, outputs and states are all one-hot encoded,

(c) (3) How many bits are needed for the input, output, and states?

# bits for in = 4 (4 choices)

# bits for out = 3 (3 choices)

# bits for state = 4 (4 states)

(d) (4) What is the logic for nx\_state:SA? out:Z? You can define your mapping for part (c) to write this Boolean function.

SA → state[0]    SD → state[3]    A → inp[0]    D → inp[3]    X → out[0]  
 SB → state[1]    SC → state[2]    B → inp[1]    C → inp[2]    Y → out[1]  
 SC → state[2]    C → inp[2]

$$nx\_state:SA = (inp[0] \wedge (state[2] \vee state[0] \vee state[1]))$$

$$out:Z = (state[0] \wedge inp[2]) \vee (state[2] \wedge inp[0]) \vee (state[1] \wedge inp[3]) \vee (state[3] \wedge inp[1])$$

(e) (3) If nx\_state:SA is written as a fully-disjunctive normal form, how many product terms are there?

# product terms = 3

Now assume that states,  $st[1:0]$ , are assigned as gray code: SA=2'b00, SB=2'b01, SC=2'b11, SD=2'b10. The inputs,  $in[1:0]$ , are also assigned as gray code where A=2'b00, B=2'b01, C=2'b11, D=2'b10; and outputs,  $out[1:0]$ , are X=00, Y=01, Z=11.

(f) (2) How many columns (inputs+outputs) and rows are in this truth table?

# columns = 4

# rows = 14 (not including don't cares) with don't cares: 16

(g) (5) Use the Karnaugh map below to determine the logic for out[0]. How many prime implicants are there? How many are essential?

in[1:0]

	"00"	"01"	"11"	"10"
out[0]	"00"	"01"	"11"	"10"
"00"	0	1	1	X
"01"	1	0	1	1
"11"	1	1	0	1
"10"	X	1	1	0

# prime implicants = 4

# essential prime implicants = 4

(h) (3) Write the Boolean expression for out[0].

$$out[0] = (in[1] \wedge \neg st[0]) \vee (in[0] \wedge \neg st[0]) \vee (\neg in[0] \wedge st[0]) \vee (\neg in[1] \wedge st[0])$$

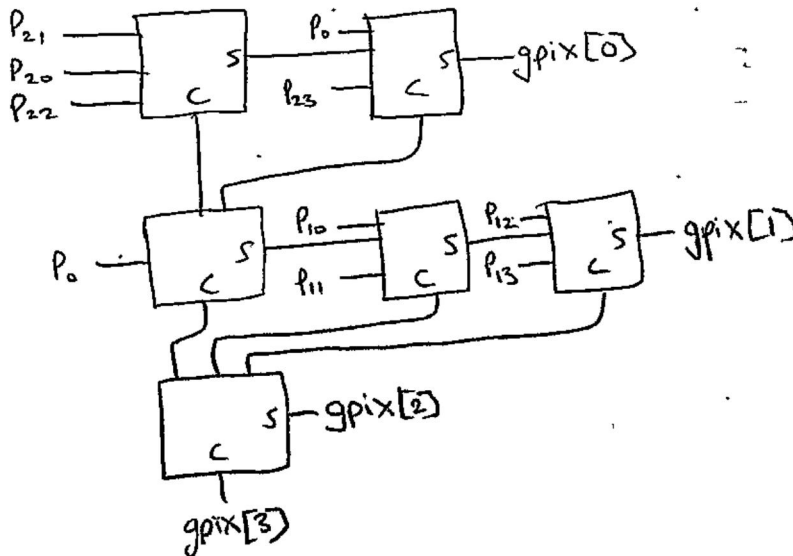
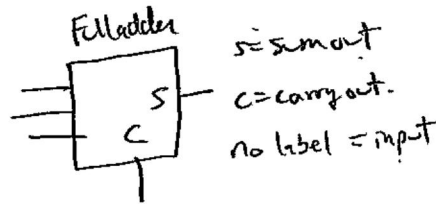
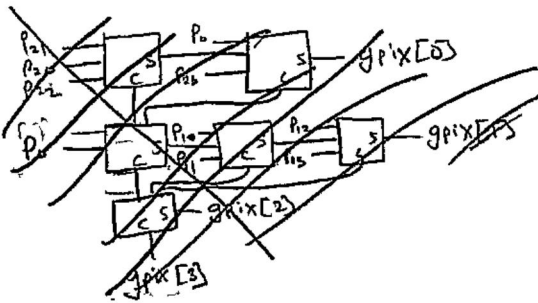
(i) (4) How many states do you need if you want to convert the FSM to a Moore Machine?

# states = 12

A field of black ("0") and white ("1") pixels can be "blurred" into gray values by taking a weighted-average that includes neighboring pixels as shown in the figure. Each converted gray-valued pixel,  $gpix$ , is a 4-bit value and is computed from 9 binary inputs,  $p_{xy}$ , based on the equation  $gpix_0[3:0] = 3 \cdot p_0 + 2 \cdot \sum p_{1n} + \sum p_{2n}$ .

$p_{21}$	$p_{11}$	$p_{20}$
$p_{12}$	$p_0$	$p_{10}$
$p_{22}$	$p_{13}$	$p_{23}$

You have available 1-bit Full Adders (FA with 3-inputs and 2-outputs) as building blocks for implementing a design. The design should output  $gpix[3:0]$ . Use the fewest number of FAs to achieve this task. Show your design.



Question #3 FSM State Diagram (15)

The input to an FSM,  $Y$ , is a string of 1's and 0's. Design a **Moore** FSM that detects when a "01" sequence is followed a "10" sequence. The FSM is reset/initialized to a state for which prior inputs are all 0's. An example of the input and output is shown below and key transitions are underlined. Note that "010" does not constitute a "01" followed by a "10". The output,  $Z$ , asserts for only 1 cycle when the sequence is detected. As a design constraint, use the fewest number of states.

$Y = 0000\underline{11\underline{1000\underline{1001010101101111}}$

$Z = 00000000\underline{10000000\underline{10000000\underline{1000}}$



Question #4 System Partitioning (20)

The following algorithm calculates the combinatorics function  $C(n,k) = n! / ((n-k)!k!)$  (commonly referred to as n-choose-k or  $nCk$ ).

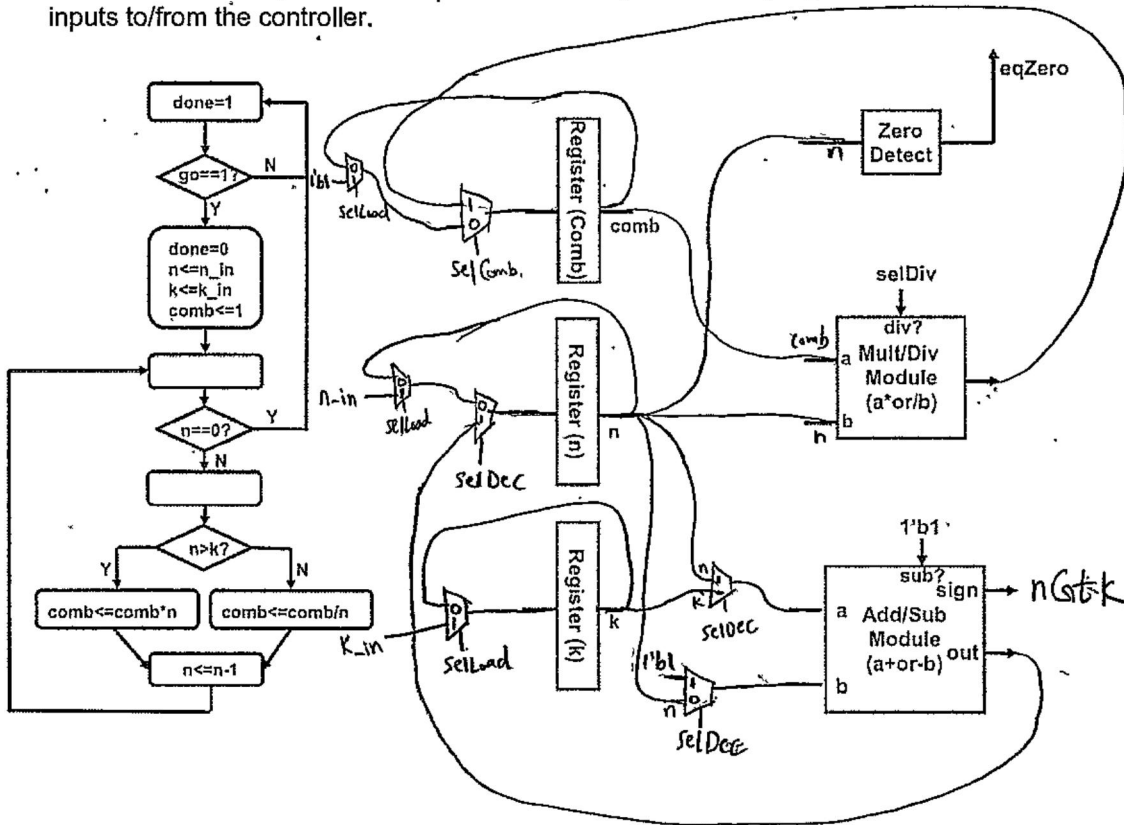
```

n = n_in;
k = k_in;
comb = 1;
while (n > 0) {
  if (n > k)
    comb = comb * n;
  else
    comb = comb / n;
  n--;
}
done = 1
    
```

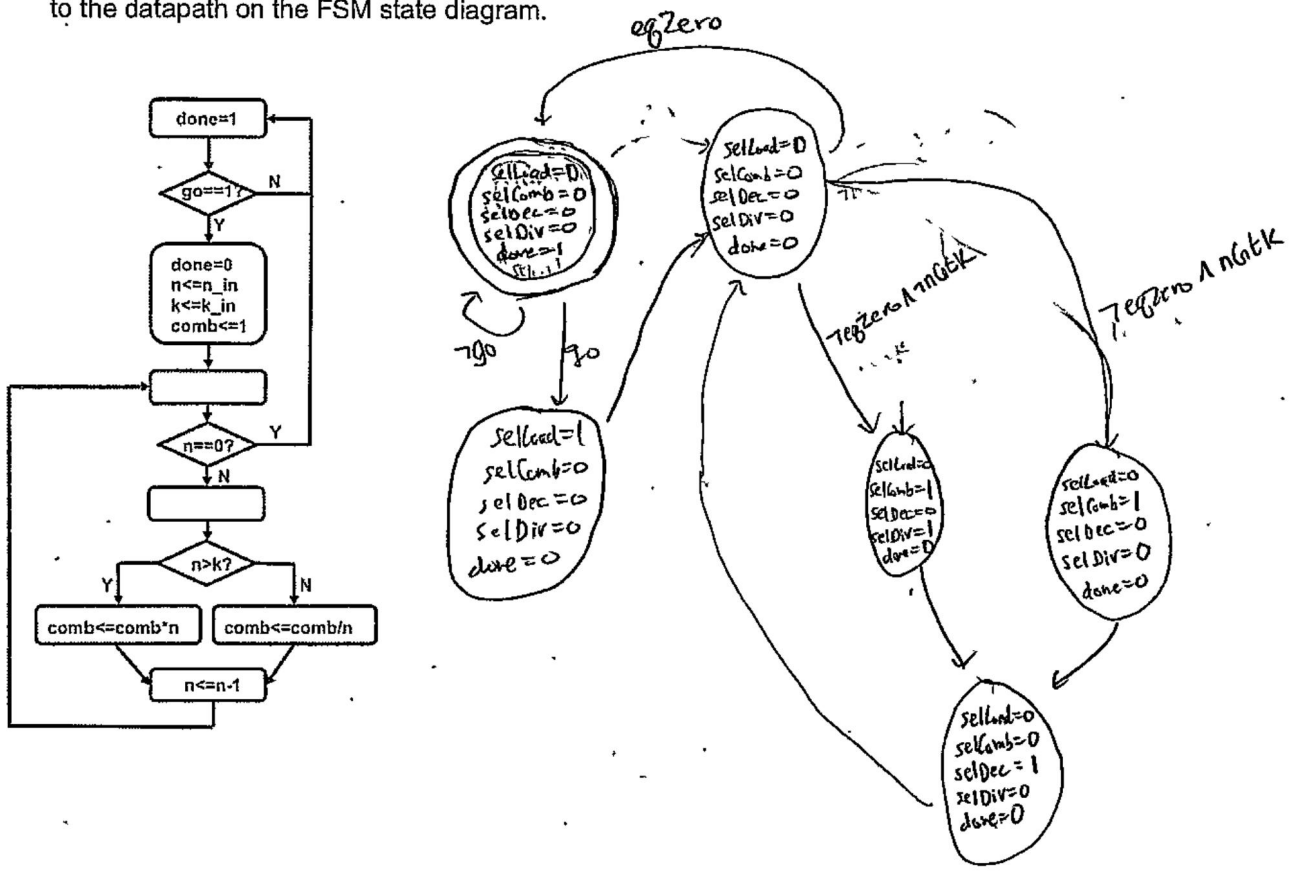
The flow diagram is already designed as shown below. A signal, *go*, is an input to the controller that triggers this algorithm. Output, *done*, is asserted by the controller when the algorithm completes and is waiting for the *go* signal. The previous computation is held in the register, *comb*. You are to complete the controller and datapath design.

Handwritten notes:  $n-k$ ,  $0$ ,  $1$ ,  $k > n$ ,  $k-n$ ,  $1$ ,  $n > k$

(a) (10) The datapath blocks available to you are also shown below (a combined multiply/divide module, an add/subtract module, and a zero detect module). You may also use as many 2:1 MUX as you choose (Note that you can only use 2:1 MUX so the select signals for each MUX is a single-bit signal from the controller). You can ignore the bit-width of any of the signals. Show the necessary connections within the datapath and any signals that need to pass as inputs to/from the controller.



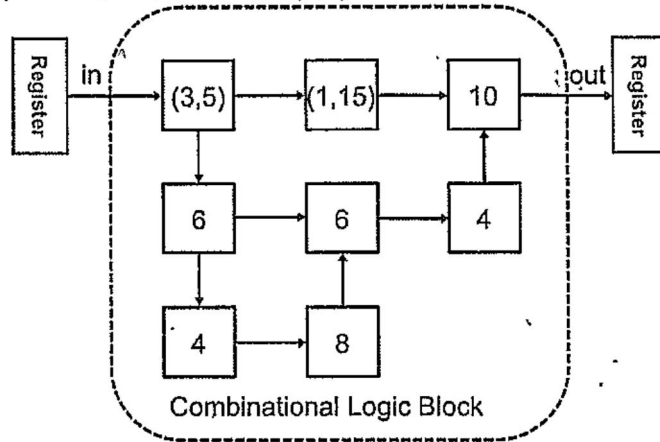
(b) (10) Design a Moore FSM for the controller. Indicate the desired control signals from controller to the datapath on the FSM state diagram.





Question #5 Timing and Pipelining (18)

The following combinational logic block can be broken down into modules. Each module have their delay as shown. For each module, the propagation and contamination delay are the same ( $t_c = t_d$ ) with the except of two blocks where the ( $t_c, t_d$ ) is shown in the block. The registers comprise of DFF with the properties  $t_s = 3$ ,  $t_H = 1$ ,  $t_{cq} = (1, 2)$ .



- (a) (4) Determine the contamination and propagation delay of the combinational logic block.

$t_{cCL} = 3 + 1 + 10 = 14$

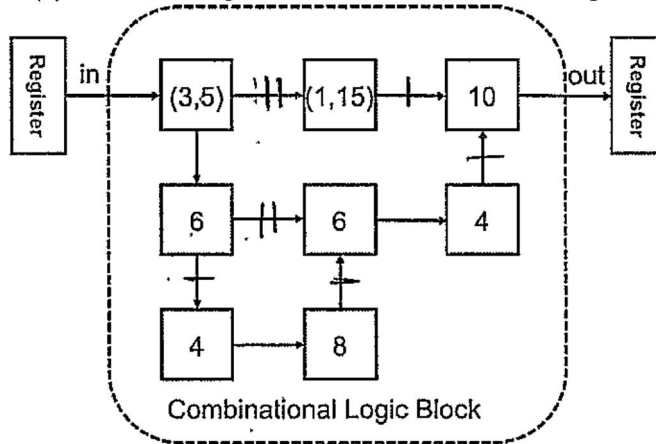
$t_{dCL} = 5 + 6 + 4 + 8 + 6 + 4 + 10 = 43$

- (b) (3) What is the minimum cycle time of the combinational logic block?

$\min(T_{cycle}) = 2 + 43 + 3 = 48$

$t_{cq} + t_{max} + t_s$

- (c) (4) We can minimize the cycle time by inserting registers. Show on the diagram below where to insert the register(s). Indicate a register with a line. Use as few registers as possible.



(d) (2) Based on the answer in (c) determine the new minimum cycle time.

$$\min(T_{\text{cycle}}) = 2 + 3 + 15 = \boxed{20}$$

$$t_{\text{dmax}} = 15$$

(e) (3) During verification of the design in (d), an engineer found that the DFF hold time is actually longer,  $t_H=3$ . Does this pose a problem? Explain your answer.

Yes or No

Explain:

$$t_H \leq t_{\text{clkQ}} + t_{\text{comin}}$$

$$3 \leq 1 + 1$$

$3 \leq 2$  (X) hold time violated.

(f) (2) Name as many ways as you can to fix this problem?

It is hard to fix as hold time isn't dependent on the clock.

We could increase  $t_{\text{clkQ}}$  in order to satisfy the constraint, or modify the pipeline to increase  $t_{\text{comin}}$ , which would fix the problem at the cost of cycle time probably.

## Question #6 (12)

An incomplete Verilog code for a module is shown below:

```

module final (
    input [3:0] st,
    output [3:0] nx_st,
    input [1:0] in,
    output [1:0] out,
    input go, reset, done, clock
);

<(a) missing TYPE> [1:0] out;
<(a) missing TYPE> [2:0] nx_st;

always @(<(b)-missing activation list>) begin
    case (st)
        3'b000: nx_st=3'b001;
        3'b001: nx_st=3'b010;
        3'b010:
            if (in[0] != go)
                nx_st=3'b100;
            else
                nx_st=3'b010;
        3'b100:
            if (in[0] != go)
                nx_st=3'b100;
            else
                nx_st=3'b001;
        default:
            nx_st = {in[0],1'b0, reset};
    endcase
end

assign out[0] = nx_st[1] | in[1];
//(c) out[1] is the output of a mux that selects 1'b0 when reset else nx_st[0]
endmodule

```

(a) (2) What should be the declared type for the following signals:

wire [1:0] out;  
reg [2:0] nx\_st;

(b) (2) What should go in the activation list of the always @()? Choose only the signals that needs to be there. You may not use \*.

Activation list = {st or in or reset or go}

(c) (5) The signal, *out[1]*, is the output of a 2:1 MUX that uses input, *reset*, to choose between input of 1'b0 (when *reset*==1), and *nx\_st[0]* (when *reset*==0). Write the Verilog code for this signal in three different ways (continue next page):

```

// Library module provided
module mux21(muxout, muxselA, inputA, inputB);
    // muxselA ==1 chooses inputA
    // module details not shown
endmodule

```

Declarative Verilog:

```
assign out[1] = reset ? 1'b0 : nx-st[0];
```

Structural Verilog (using the library module above)

```
mux2 muxout1 (out[1], reset, 1'b0, nx-st[0]);
```

Procedural Verilog (note that out variable will need to be declared differently):

```
reg[1:0] out;
```

```
always @(nx-st or reset) begin
    out[1] = reset ? 1'b0 : nx-st[0];
end.
```

(d) (3) Four different ways of implementing a function is shown below. Which of them are the same? Circle all that are the same.

```
(1)
always (@posedge clock) begin
    y <= z;
    x <= y;
end
```

```
(2)
always (@posedge clock) begin
    y = z;
    x = y;
end
```

```
(3)
always (@posedge clock) begin
    x = y;
    y = z;
end
```

```
(4)
always (@posedge clock) begin
    z = y;
    y = x;
end
```

Question #7 Short Answers (15)

(a) (4) For the following Karnaugh map, the Boolean expression for the function

$$Z = (\neg A \wedge \neg B \wedge \neg C) \vee (A \wedge \neg B \wedge D) \vee (A \wedge \neg B \wedge C)$$

		AB			
		"00"	"01"	"11"	"10"
CD	"00"	1	0	0	0
	"01"	1	0	0	1
	"11"	0	0	0	1
	"10"	0	0	0	1

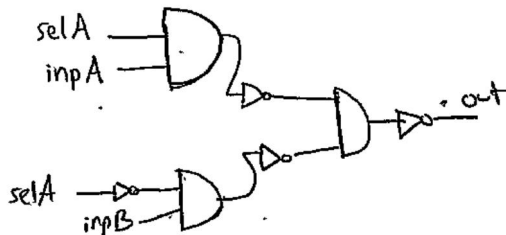
What input conditions and transition has a potential for causing a glitch (static hazard) at the output?

a hazard could happen if we transition from  $\overset{BCD}{1001}$  to  $\overset{BCD}{0001}$ .

(b) (2) How would you resolve the issue in (a) by adjusting the Boolean expression?

$$Z = (\neg A \wedge \neg B \wedge \neg C) \vee (A \wedge \neg B \wedge D) \vee (A \wedge \neg B \wedge C) \vee (\neg B \wedge \neg C \wedge D)$$

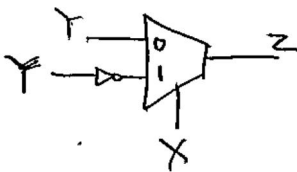
(c) (3) If you only have 2-input AND gates and inverters available, how would you build a 2:1 multiplexer? (out selects between *inpA* and *inpB* with the select signal, *selA*)



		selA					
		0	1				
AB	00	0	0	0	0	0	0
	01	1	0	1	0	1	0
	11	1	1	1	1	1	1
	10	0	1	0	1	0	1

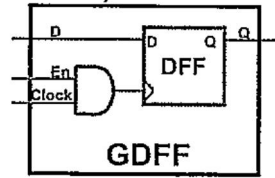
$$out = (B \wedge \neg selA) \vee (A \wedge selA)$$

(d) (3) If you only have 2:1 MUX and inverters available, how would you implement  $Z = X \text{ xor } Y$ ?



X	Y	Z
0	0	0
0	1	1
1	0	1
1	1	0

(e) (3) A designer modified the basic DFF as shown below to make a GDFF where the clock signal is ANDed with an Enable signal. This approach is known as "clock gating". How does the GDFF's characteristics compare to that of the DFF? Select the answer.



Setup:

$t_{s\_GDFF} > t_{s\_DFF}$

$t_{s\_GDFF} = t_{s\_DFF}$

$t_{s\_GDFF} < t_{s\_DFF}$

Hold:

$t_{H\_GDFF} > t_{H\_DFF}$

$t_{H\_GDFF} = t_{H\_DFF}$

$t_{H\_GDFF} < t_{H\_DFF}$

Clock-Q Delay:

$t_{c2Q\_GDFF} > t_{c2Q\_DFF}$

$t_{c2Q\_GDFF} = t_{c2Q\_DFF}$

$t_{c2Q\_GDFF} < t_{c2Q\_DFF}$

