

University of California  
Los Angeles  
Computer Science Department

CSM51A/EEM16 Midterm Exam  
Winter Quarter 2016  
February 8<sup>th</sup> 2016

Stats  
highest - 100  
lowest - 33  
avg -  $\approx 80$   
media -  $\approx 80$   
Final more difficult.

This is a closed book exam. Absolutely nothing is permitted except pen, pencil and eraser to write your solutions. Any academic dishonesty will be prosecuted to the full extent permissible by university regulations.

**Time allowed 100 minutes.**

Problem (possible points)	Points
1 (20)	20
2 (20)	20
3 (20)	20
4 (20)	20
5 (20)	20
Total (100)	100

20

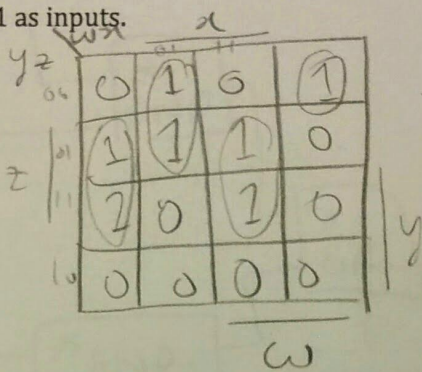
**Problem 1 (20 points)**

Use only the "E" gate defined below to implement Boolean function:

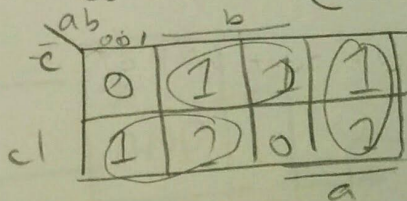
$$F = w'xy' + wxz + w'x'z + wx'y'z'$$

You may also use constants 0 and 1 as inputs.

a	b	c	E(a,b,c)
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	0

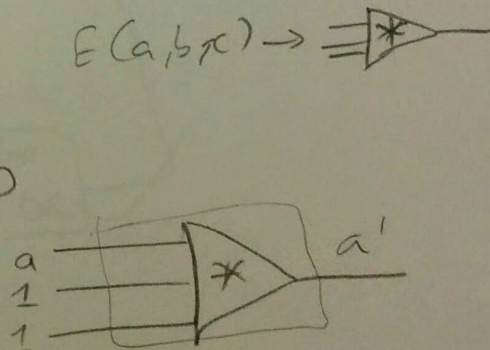


$$F(a,b,c) = (a+b+c) \cdot (a'+b'+c')$$

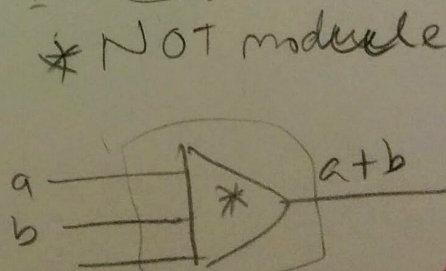


$$F(a,b,c) = a'c + bc' + ab'$$

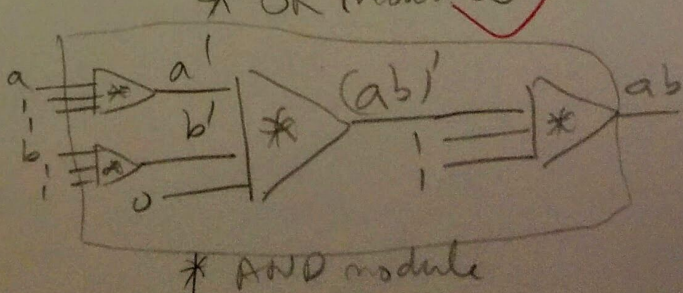
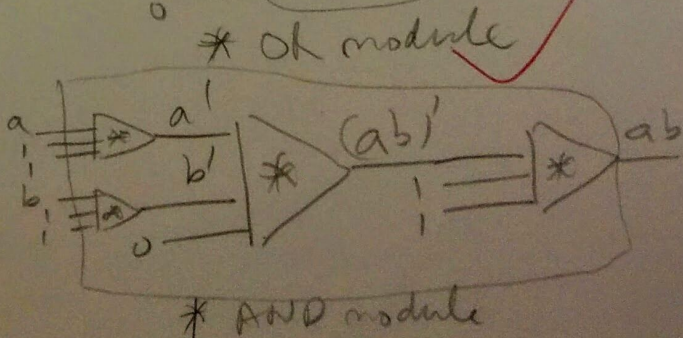
$$F(a,1,1) = a' \cdot 1 + 1 \cdot 0 + a \cdot 0 = a' \rightarrow \text{NOT}$$



$$F(a,b,0) = a' \cdot 0 + b \cdot 1 + ab' = ab' + b = a + b$$

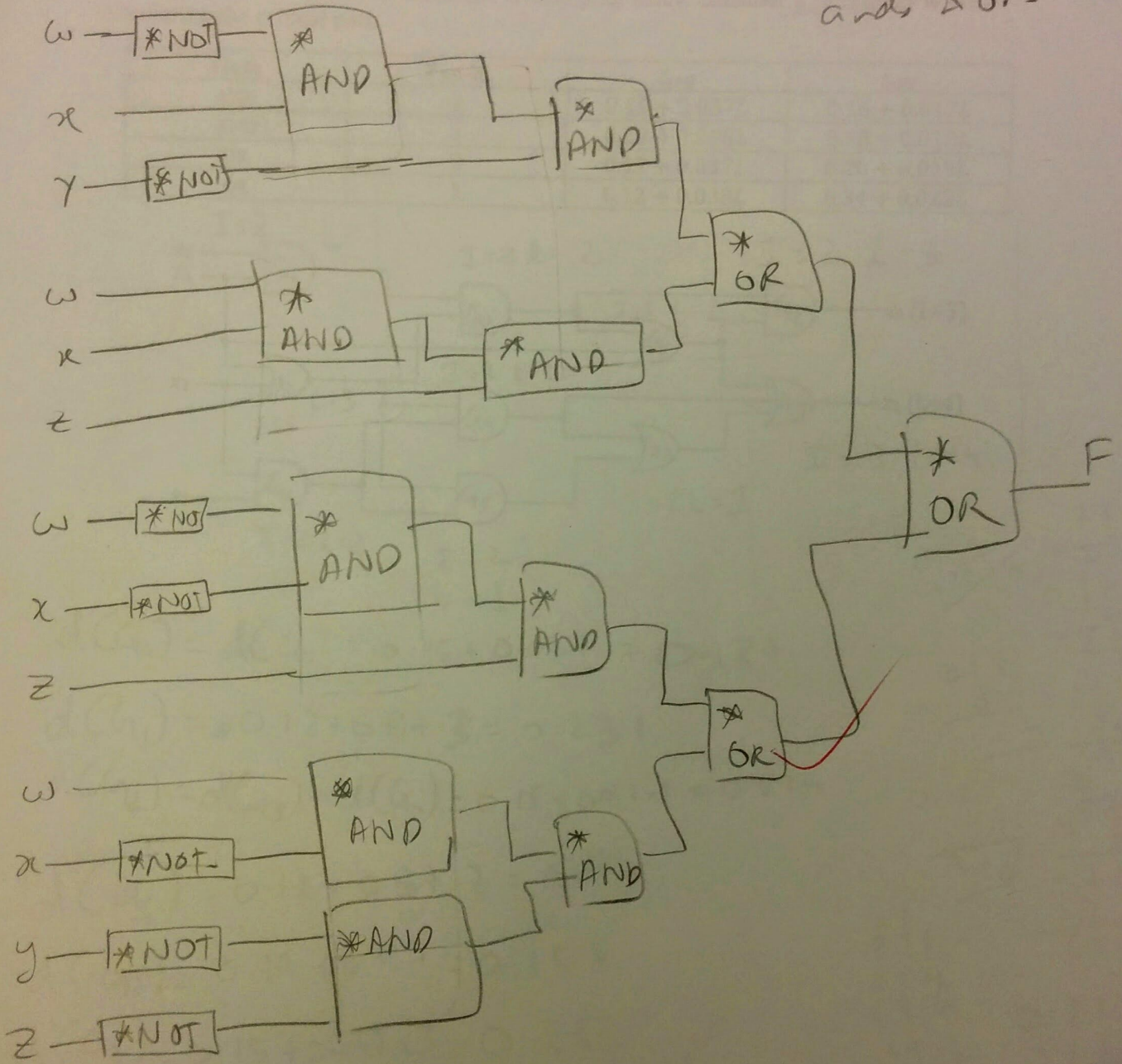


$$F(a',b',0) = a \cdot 0 + b' \cdot 1 + a'b = a' + b' = (ab)'$$



Problem 1) Extra Page

because I have 2 input  
and 2 or's

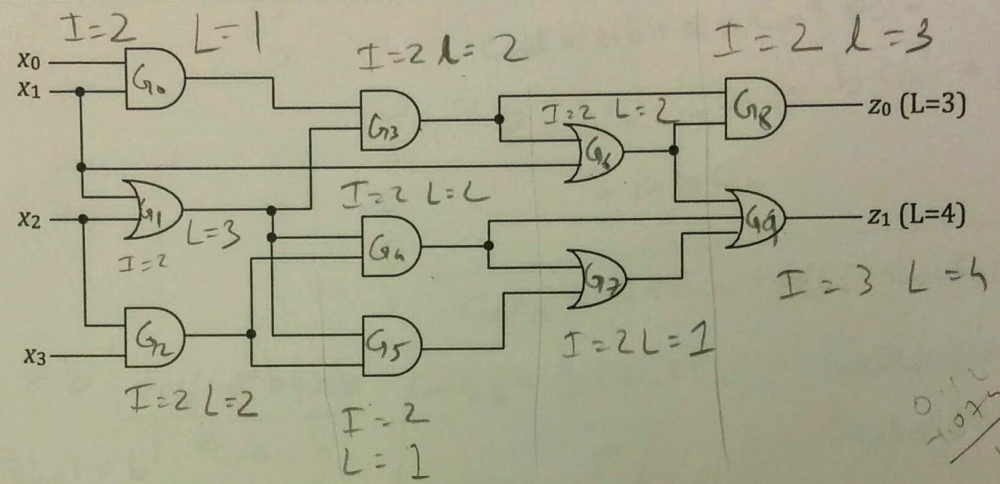


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**Problem 2 (20 points)**

Given the network below, calculate the critical path delay. Consider L → H delay when calculating the critical path.

Gate	Fan-in	$t_{pLH}$	$t_{pHL}$
AND	2	$0.15 + 0.037L$	$0.16 + 0.017L$
AND	3	$0.20 + 0.038L$	$0.18 + 0.018L$
OR	2	$0.12 + 0.037L$	$0.20 + 0.019L$
OR	3	$0.12 + 0.038L$	$0.34 + 0.022L$



$d(G_0) = d(G_5) = 0.15 + 0.037 = 0.187$

$d(G_1) = 0.12 + 0.037 \cdot 3 = 0.231$

$d(G_2) = d(G_3) = d(G_4) = 0.15 + 0.037 \times 2 = 0.224$

$d(G_6) = 0.12 + 0.037 \cdot 2 = 0.194$

$d(G_7) = 0.12 + 0.037 = 0.157$

$d(G_8) = 0.15 + 0.037 \times 3 = 0.261$

$d(G_9) = 0.12 + 0.038 \times 4 = 0.272$

Critical path

$G_1 \rightarrow G_3 \rightarrow G_6 \rightarrow G_9$

$T(CP) = 0.231 + 0.224 + 0.194 + 0.272 = 0.921$

we can see that looking gate with max delay in every layer

Handwritten calculations for gate delays:

- $0.15 + 0.037 = 0.187$
- $0.12 + 0.037 \cdot 3 = 0.231$
- $0.15 + 0.037 \cdot 2 = 0.224$
- $0.12 + 0.037 \cdot 2 = 0.194$
- $0.12 + 0.037 = 0.157$
- $0.15 + 0.037 \cdot 3 = 0.261$
- $0.12 + 0.038 \cdot 4 = 0.272$
- Sum:  $0.231 + 0.224 + 0.194 + 0.272 = 0.921$

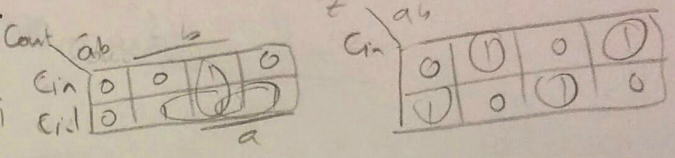
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**Problem 3 (20 points)**

Four 4-bit numbers A,B,C, and D are given as inputs.  $E=A+B$ ,  $F=C+D$ . Design a system that outputs the larger number between E and F. If  $E=F$ , output either E or F. You can use any type of gates to implement your design.

First 1 bit adder

a	b	Cin	cout	z <sub>i</sub>
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1



$$cout = ab + aCin + bCin$$

$$z = a'b'Cin + a'bCin' + ab'Cin' + abcin$$

$$cout_a = b + Cin + bCin \quad cout_a' = bCin$$

$$z_{cin} = a'b' + ab \quad z_{cin'} = a'b + ab'$$

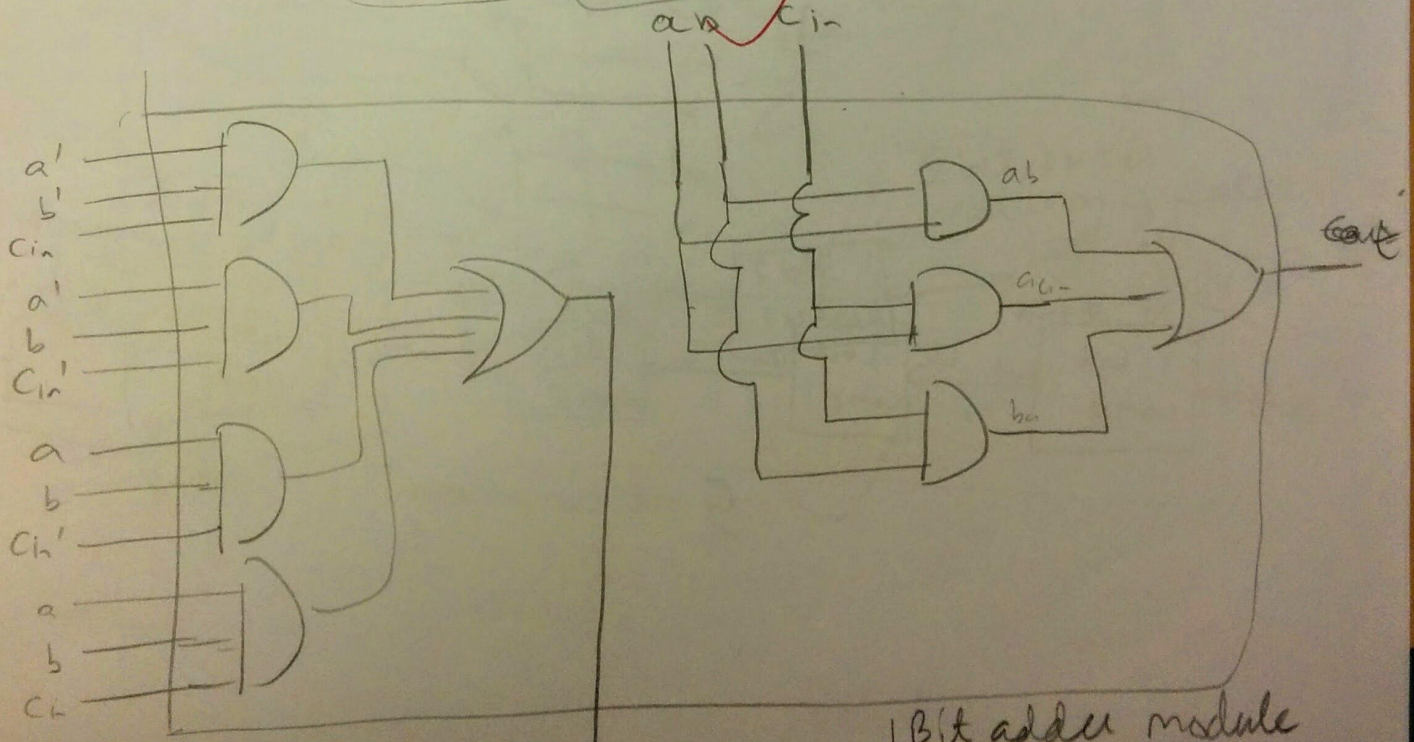
$$z_{cina} = b \quad z_{cinal} = b' \quad z_{cin'a} = b' \quad z_{cin'a'} = b$$

$$cout_{ab} = 1 + Cin + Cin = 1$$

$$cout_{ab'} = Cin$$

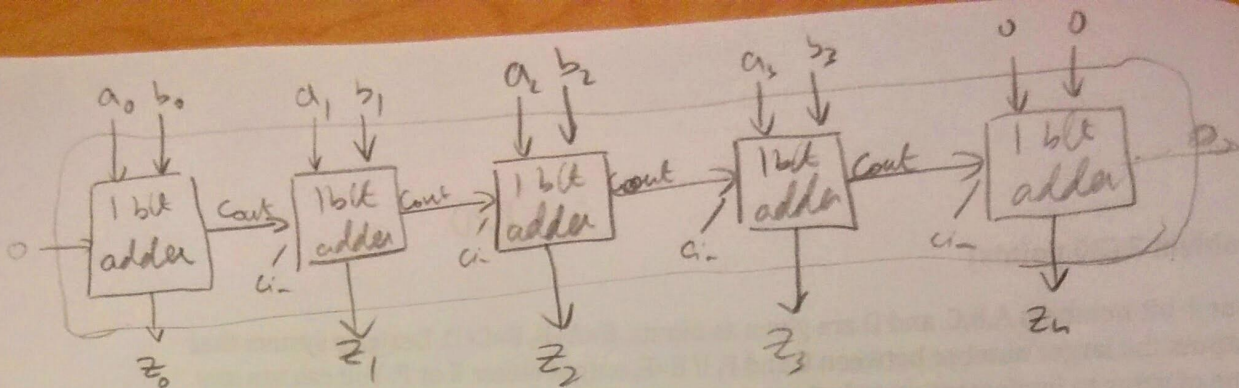
$$cout_{a'b} = Cin$$

$$cout_{a'b'} = 0$$



1 Bit adder module  
cout from each 1 bit cin to next

Z<sub>p</sub>

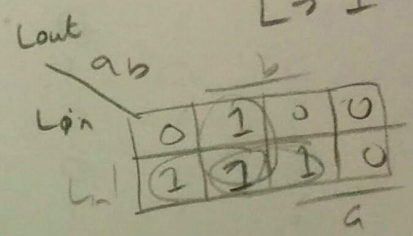


4-bit adder module

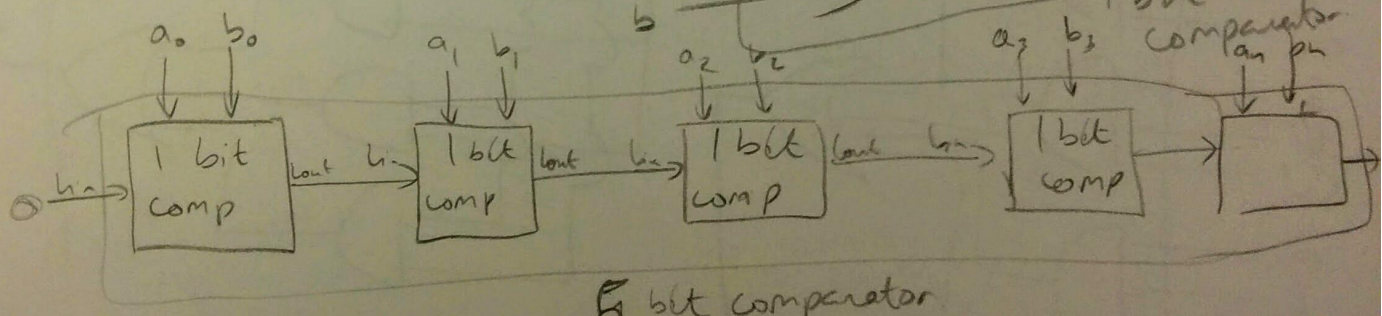
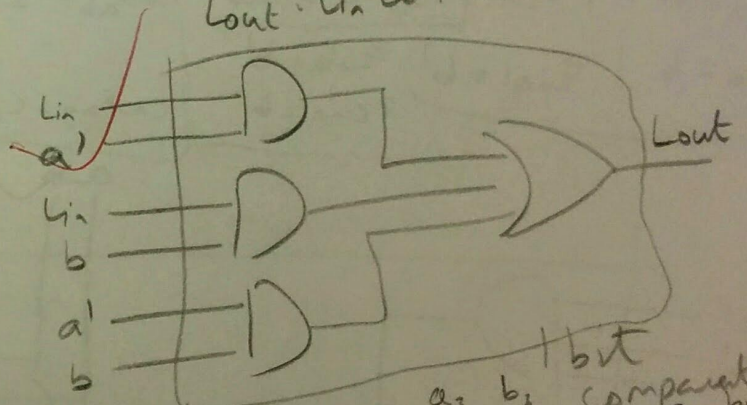
Comparator

a	b	$L_{in}$	$L_{out}$
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	1

$A, B$   
 $L \rightarrow 0 \rightarrow A > B$   
 $L \rightarrow 1 \rightarrow A \leq B$



$$L_{out} = L_{in} a' + L_{in} b + a' b$$



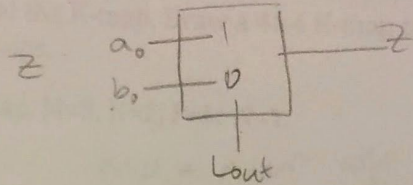
5 bit comparator

Problem 3) Extra Page

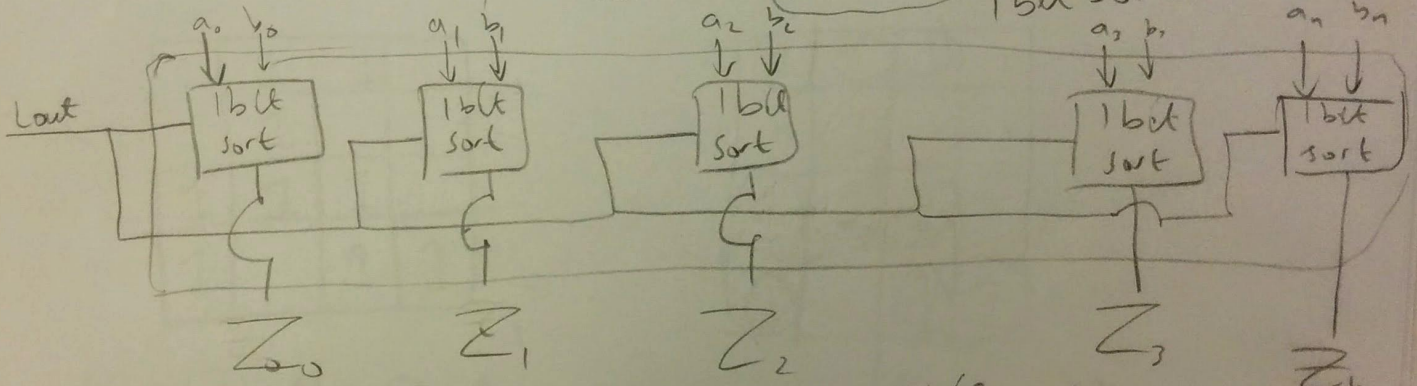
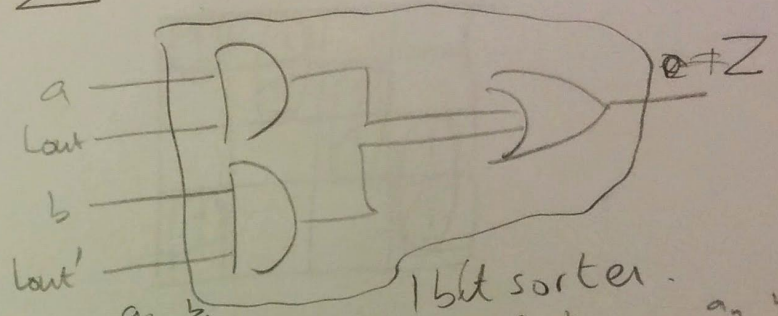
Ex 9)  $a > b$

sorter

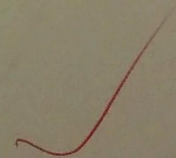
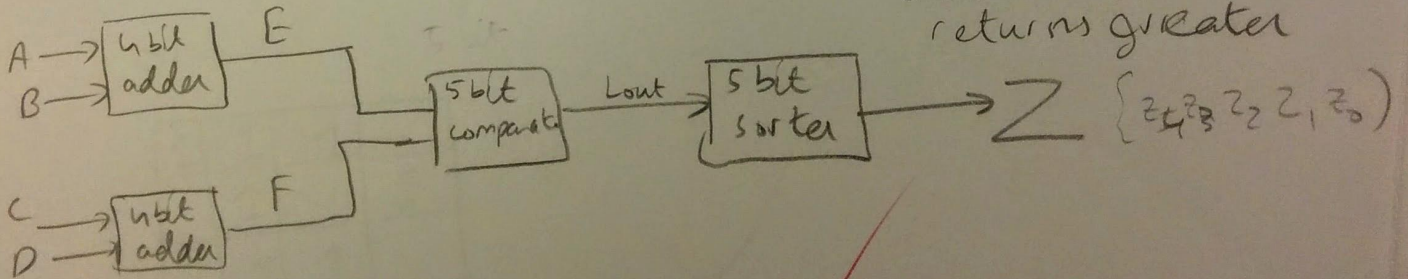
a	b	lout
0	0	0
0	0	1
0	1	0
0	1	1
1	0	0
1	0	1
1	1	0
1	1	1



$$Z = \text{lout} a + \text{lout}' b$$



5 bit sorter.  
returns greater



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**Problem 4 (20 points)**

For a K-map,  $M$  denotes the number of prime implicants of the K-map, and  $N$  denotes the number of essential prime implicants of the K-map. Draw a  $4 \times 4$  K-map that has the largest value of  $P=M-N$  among all the  $4 \times 4$  K-maps.

For example, in the following  $4 \times 4$  K-map,  $M=3$ ,  $N=2$ ,  $P=M-N=1$ .

	$x_0$				
	0	0	0	0	
	1	1	0	0	
	1	1	1	0	
$x_3$	0	0	1	0	$x_2$
					$x_1$

max num of prime  
min number of ep

1	1		
1	1	1	1
		1	1
1			1

1	1	1	
	1		
1	1	1	
1		1	1

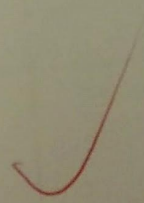
1	1	1	
	1		
1	1	1	
1		1	1

$M=12$   $N=0$   $P=12$

or

1	1	1	
	1		1
1	1	1	
1		1	1

$M=13$   $N=1$   
 $P=12$





Problem 5 (20 points)

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Use only multiplexers to design a system with input  $x \in \{0,1,2, \dots, 8\}$ , outputs  $y$  and  $z$  that implements the following equation

$$(x)_{10} = (yz)_3$$

In the system,  $x$  is encoded as  $x_3x_2x_1x_0$  in binary.  $y$  is encoded as  $y_1y_0$  in binary, and  $z$  is encoded as  $z_1z_0$  in binary.

Note that the outputs  $y$  and  $z$  represent the two digits of a base-3 number.

For example, if  $x=7$  ( $x_3x_2x_1x_0=0111$ ), then the system will solve:  $(7)_{10} = (21)_3$ . Thus  $y = 2$  ( $y_1y_0=10$ ) and  $z = 1$  ( $z_1z_0=01$ ).

$x_3$	$x_2$	$x_1$	$x_0$	$y_1$	$y_0$	$z_1$	$z_0$
0	0	0	0	0	0	0	0
0	0	0	1	0	0	0	1
0	0	1	0	0	0	1	0
0	0	1	1	0	1	0	0
0	1	0	0	0	1	0	1
0	1	0	1	0	1	1	0
0	1	1	0	1	0	0	0
0	1	1	1	1	0	0	1
1	0	0	0	1	0	1	0
1	0	0	1	-	-	-	-
1	0	1	0	-	-	-	-
1	0	1	1	-	-	-	-
1	1	0	0	-	-	-	-
1	1	0	1	-	-	-	-
1	1	1	0	-	-	-	-
1	1	1	1	-	-	-	-

ref 8  $3 \times 2 + 2$   
 $6 + 2 = 8$

00
01
02
10
11
12
20
21
22
$3 \times 2 + 2$
8

$y_1$	$x_3$	$x_2$	$x_2$	
$x_1, x_0$	0	0	-	
	0	0	-	
$x_0$	0	0	-	
	0	0	-	

$y_0$	$x_2$			
	0	1	-	0
	0	1	-	-
	0	1	-	-
	0	1	-	-

$z_1$	$x_2$			
	0	0	-	0
	0	0	-	-
$x_0$	1	1	-	-
	1	1	-	-

$z_0$	$x_2$			
	0	0	-	0
$x_0$	1	1	-	-
	1	1	-	-
	0	0	-	-

$y_1 = x_3$   
 $y_0 = x_2$   
 $z_1 = x_1$   
 $z_0 = x_0$

$y_1$	$x_2$			
	0	0	-	1
$x_0$	0	0	-	-
	0	1	-	-
	0	1	-	-

$y_0$	$x_2$			
	0	1	-	0
$x_1$	0	1	-	-
	1	0	-	-
	0	0	-	-

$z_1$	$x_2$			
	0	0	-	1
$x_0$	0	1	-	-
	0	0	-	-
	1	0	-	-

$z_0$	$x_2$			
	0	1	-	0
$x_0$	1	0	-	-
	0	1	-	-
	0	0	-	-

Problem 5) Extra Page

$$y_1 = x_3 + x_2 x_1 \quad y_0 = x_2 x_1' + x_2' x_1 x_0$$

$$z_1 = x_3 + x_2 x_1' x_0 \quad z_0 = x_2 x_1' x_0' + x_2' x_1 x_0 + x_2 x_1 x_0$$

$$y_1 x_3 = 1 + x_2 x_1 = 1$$

$$y_1 x_3' = x_2 x_1$$

$$y_1 x_3' x_2 = x_1$$

$$y_1 x_3' x_2' = 0$$

$$y_0 x_2 = x_1'$$

$$y_0 x_2' = x_1 x_0$$

$$y_0 x_2' x_1 = x_0$$

$$y_0 x_2' x_1' = 0$$

$$z_1 x_3 = 1$$

$$z_1 x_3' = x_2 x_1' x_0$$

$$z_1 x_3' x_2 = x_1' x_0$$

$$z_1 x_3' x_2' = 0$$

$$z_1 x_3' x_2 x_0 = x_1'$$

$$z_1 x_3' x_2 x_0' = 0$$

$$z_0 x_1 = x_2' x_0$$

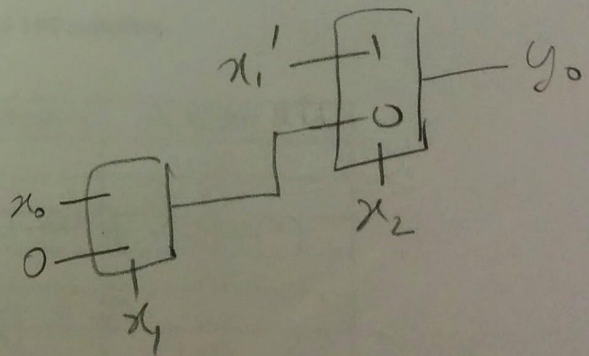
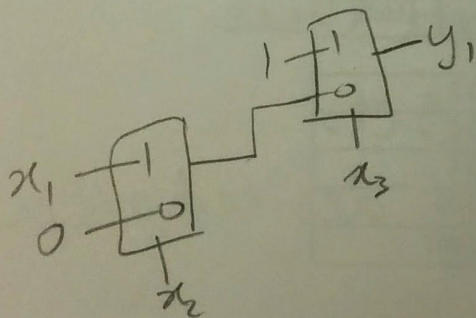
$$z_0 x_1' = x_2 x_0' + x_2' x_0$$

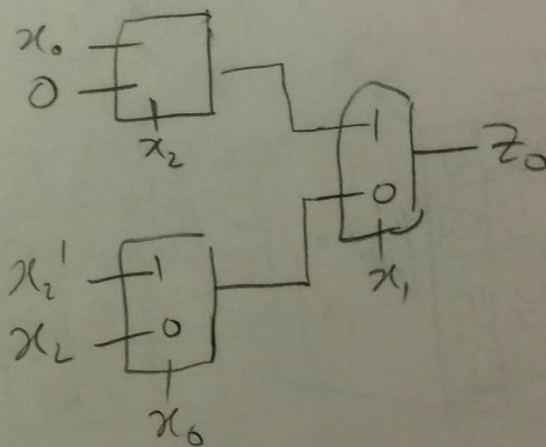
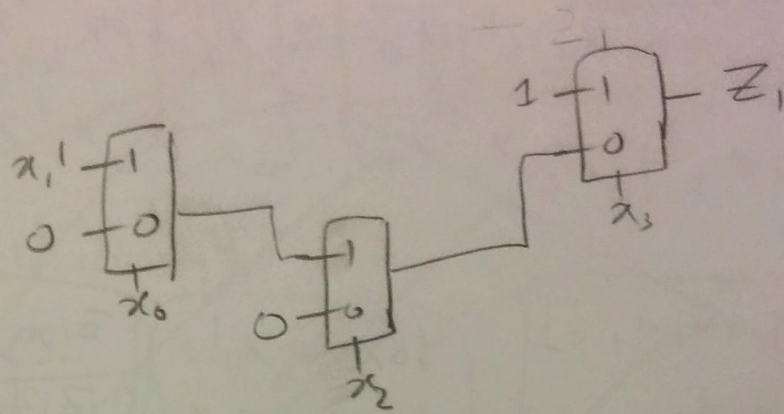
$$z_0 x_1' x_0 = x_2'$$

$$z_0 x_1' x_0' = x_2$$

$$z_0 x_1 x_2 = x_0$$

$$z_0 x_1 x_2' = 0$$





We can box all these  
multiplexer sets  
to form desired output