

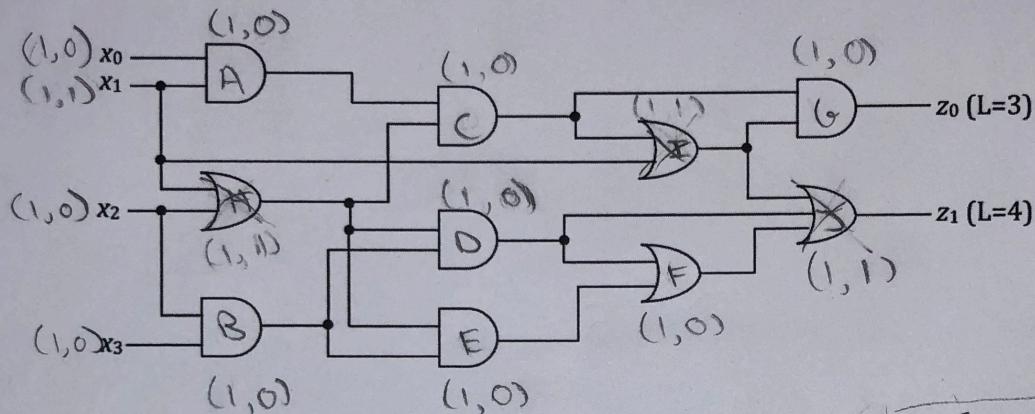
Problem 1 (20 points):

20

Given the network below, find the critical path and calculate critical path delay. Assume the values of  $(x_3, x_2, x_1, x_0)$  are initially  $(1, 1, 1, 1)$  and they change to  $(0, 0, 1, 0)$  in the next clock cycle. Now, choose a gate on the critical path which maximally decreases overall delay when the gate decreases its delay by 20%. Finally, find the critical path in the new network and its length.

Gate	Fan-in	$t_{pLH}$	$t_{pHL}$
AND	2	$0.15 + 0.037L$	$0.16 + 0.017L$
AND	3	$0.20 + 0.038L$	$0.18 + 0.018L$
OR	2	$0.12 + 0.037L$	$0.20 + 0.019L$
OR	3	$0.12 + 0.038L$	$0.34 + 0.022L$

(first, second)



#s in parens are multiplied by 0.037

06: 
$$\begin{aligned} & 0.15 + (1) \\ & 0.18 + (2) \\ & 0.12 + (3) \end{aligned}$$

06 
$$\begin{aligned} & 0.15 + (1) \\ & \text{switched} \\ & 0.15 + (2) \\ & 0.12 + (2.4) \end{aligned}$$

vs. competitor: 
$$\begin{aligned} & 0.15 + (2) \\ & 0.15 + (2) \\ & 0.12 + (1) \end{aligned}$$

	Initial delay	New delay
A	$0.15 + 0.037$	$0.15 + 0.037$
B	$0.15 + (2)0.037$	$0.15 + (2)0.037$
C	$0.15 + (2)0.037$	$0.15 + (2)0.037$
D	$0.15 + (2)0.037$	$0.15 + (2)0.037$
E	$0.15 + 0.037$	$0.15 + 0.037$
F	$0.12 + 0.037$	$0.12 + 0.037$
G	$0.15 + (3)0.037$	$\frac{4}{5}(0.15 + (3)0.037)$

initial critical path:  $(x_0, A, C, G, z_0)$ , length = 3  
 new critical path:  $(x_0, A, C, G, z_0)$ , length = 3

The critical path does NOT change (but its delay does)

**Problem 2 (20 points):**

You are given the following Boolean function.

$$F(x_6, x_5, x_4, x_3, x_2, x_1, x_0) = x_6x_5x_4x_3 + x_6x_5'x_2 + x_6x_5'x_3'x_2' + x_6'x_1x_0 + x_6'x_5x_0'$$

(1)      (2)      (3)      (4)      (5)

Given the universal operation E as specified in the table, implement F using only the gates specified by E.

$$\text{NOT}(x) = E(x, 0)$$

$$\begin{aligned} \text{OR}(x, y) &= E(\text{NOT}(x), y) \\ &= E[E(x, 0), y] \end{aligned}$$

X	Y	E(X, Y)
0	0	1
0	1	1
1	0	0
1	1	1

$$\text{NAND}(x, y) = \text{OR}(\text{NOT}(x), \text{NOT}(y))$$

<u>NAND(x, y)</u>	<u>OR(x, y)</u>
1	0
1	1
1	1
0	1

$$xy = (x' + y')'$$

$$\text{AND} = \text{NOT}[\text{OR}[\text{NOT}(x), \text{NOT}(y)]]$$

$$\begin{aligned} &= \text{NOT}[\text{OR}[\text{E}(x, 0), \text{E}(y, 0)]] \\ &= \text{NOT}[\text{E}[\text{E}[\text{E}(x, 0), 0], \text{E}(y, 0)]] \\ &= \text{NOT}[\text{E}[x, \text{E}(y, 0)]] \\ &= \text{E}[\text{E}[x, \text{E}(y, 0)], 0] \end{aligned}$$

$$*: x = \text{E}[\text{E}(x, 0), 0]$$

$$x = \text{NOT}[\text{NOT}(x)]$$

$$(1) x_6x_5x_4x_3 = (x_6x_5)(x_4x_3) = \underbrace{\text{E}[\text{E}[x_6, \text{E}(x_5, 0)], 0]}_{L1} \underbrace{\text{E}[\text{E}[x_4, \text{E}(x_3, 0)], 0]}_{R1}$$

$$= \underbrace{\text{E}[\text{E}[L1, \text{E}(R1, 0)], 0]}_{G1}$$

$$(2) x_6x_5'x_2 = (x_6x_5')(x_2) = \underbrace{\text{E}[\text{E}[x_6, x_5], 0]}_{L2} \underbrace{x_2}_{R2}$$

$$= \underbrace{\text{E}[\text{E}[L2, \text{E}(R2, 0)], 0]}_{G2}$$

$$(3) x_6x_5'x_3'x_2' = (x_6x_5')(x_3'x_2') = \underbrace{\text{E}[\text{E}[x_6, x_5], 0]}_{L3} \underbrace{\text{E}[\text{E}[\text{E}(x_3, 0), x_2], 0]}_{R3}$$

$$= \underbrace{\text{E}[\text{E}[L3, \text{E}(R3, 0)], 0]}_{G3}$$

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**Problem 2 (Extra Page):**

$$(4) x_6' x_1 x_6 = (x_6' x_1)(x_6) = \underbrace{E\{E\{E(x_6, 0), E(x_1, 0)\}, 0\}}_{L4} x_6 \underbrace{R4}_{R4}$$

$$= \underbrace{E\{E\{L4, E(R4, 0)\}, 0\}}_{G4}$$

$$(5) x_6' x_5 x_6' = (x_6' x_5)(x_6') = \underbrace{E\{E\{E(x_6, 0), E(x_5, x_6)\}, 0\}}_{L5} \underbrace{E(x_6, 0)}_{R5}$$

$$= \underbrace{E\{E\{L5, E(R5, 0)\}, 0\}}_{G5}$$

$$F = G1 + G2 + G3 + G4 + G5$$

$$G1 + G2 : E\{E(G1, 0), G2\}$$

$$G3 + (G1 + G2) : E\{E(G3, 0), E\{E(G1, 0), G2\}\}$$

$$G4 + G5 : E\{E(G4, 0), G5\}$$

$$Z1 + Z2 = E\{E(Z1, 0), Z2\}$$

$$F = E\{E(Z1, 0), Z2\}$$

$$(G1 + G2 + G3)$$

Call tree diagram.

Z2 and Z1 are defined by  $G1 + G5$ , which are in turn defined by  $L1 \rightarrow L5$  and  $R1 \rightarrow R5$ . These abstractions were performed to increase readability.

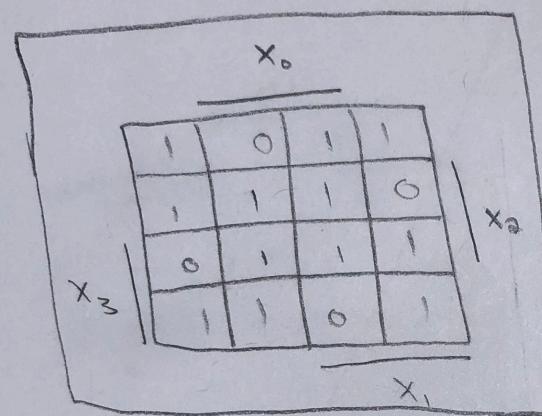
$$\text{where } F = F(x_6, x_5, x_4, x_3, x_2, x_1, x_0)$$

**Problem 3 (20 points):**

For a K-map, M denotes the number of prime implicants of the K-map, and N denotes the number of essential prime implicants of the K-map. Draw a  $4 \times 4$  K-map that has the largest value of  $P = M - N$  among all the  $4 \times 4$  K-maps.

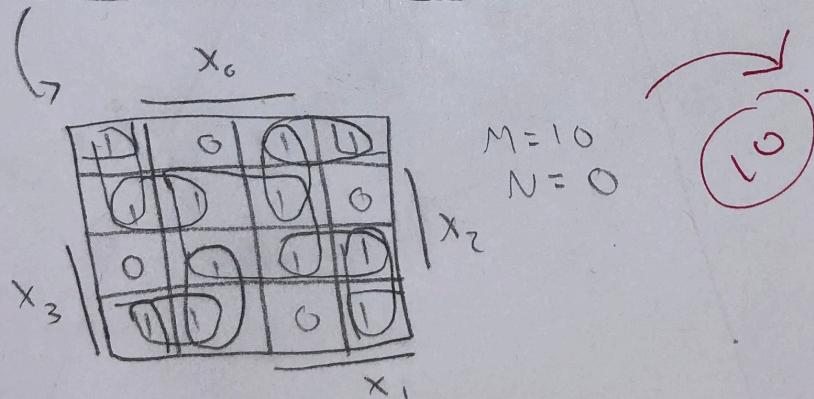
For example, in the following  $4 \times 4$  K-map,  $M=3$ ,  $N=2$ ,  $P=M-N=1$ .

$x_0$			
0	0	0	0
1	1	0	0
1	1	1	0
0	0	1	0



$$P = M - N = 10 - 0 = 10$$

with  
prime implicants  
shown



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**Problem 4 (20 points):**

Given an input stream X, we want to recognize interchangably patterns A and B. We recognize A first, then B, followed by A again, then B again and so on.

For example,

1. Assume  $X = 01011010010101$ ,  $A=101$  and  $B=001$ .

We will first recognize A, then look for B. Please note that we ignore the second '101' (A) in X and we only search for B once we have found A. After finding B, we again search for A.

2. Assume that we have  $X = 1011$ ,  $A = 101$  and  $B = 011$

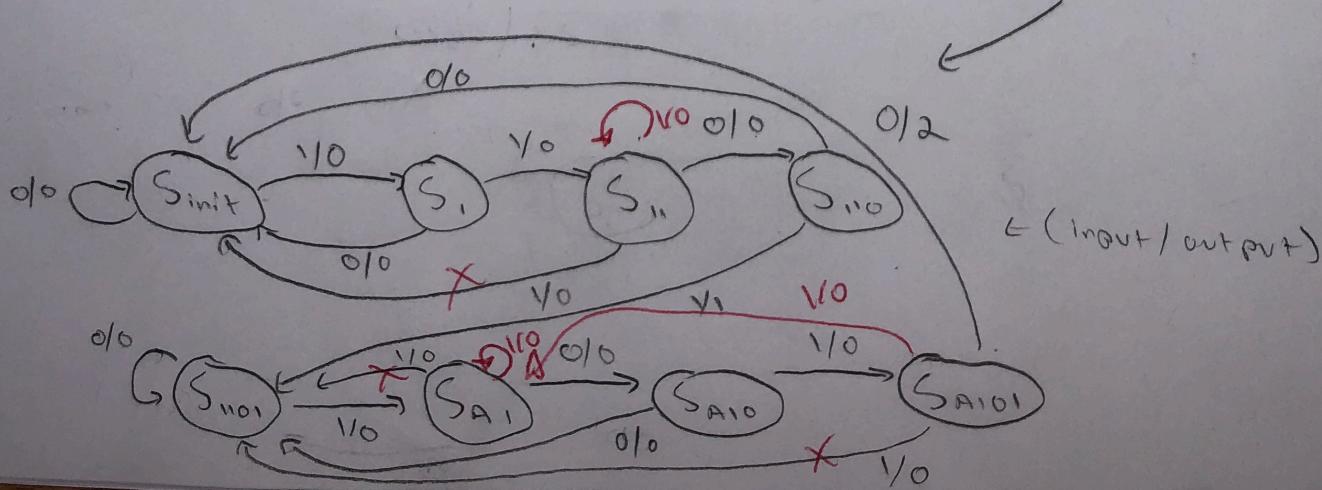
We recognize A, but we do not recognize B as we only start looking for B once we have detected A. In other words, A and B do not overlap.

Now, you are given any input stream X. Design a finite state machine such that the system outputs 1 when it recognizes pattern A='1101' and outputs 2 when it recognizes pattern B='1010' after recognizing pattern A. In all other cases the machine should output 0. Show the state transition table and state transition diagram.

<u>PS</u>	(OUTPUT/NS)		$\left\{ \begin{array}{l} 6,13 \text{ inputs} \\ 2,3 \text{ outputs} \end{array} \right.$	$A = 1101$
$S_{init}$	0	1		$B = 1010$
$S_1$	$0/S_{init}$	$0/S_1$		
$\times S_{II}$	$0/S_{II}$	$0/S_{init}$		
$S_{II}$	$0/S_{II}$	$1/S_{II}$		
$S_{III}$	$0/S_{III}$	$1/S_{III}$		
$\times S_{A1}$	$0/S_{A1}$	$0/S_{II01}$		
$S_{A10}$	$0/S_{A10}$	$0/S_{A101}$		
$\times S_{A101}$	$2/S_{init}$	$0/S_{II01}$		

STATE TRANSITION TABLE

STATE  
TRANSITION  
DIAGRAM



A=10, B=11, C=12, D=13, E=14, F=15

Problem 5 (20 points):

Perform the following conversions:

- a)  $(B2451)_{16} \rightarrow (x)_8$
- b)  $(354)_7 \rightarrow (y)_5$

(a)  $(B2451)_{16} \rightarrow (x)_8$

to  
binary  
 $\begin{array}{cccccc} 8 & 2 & 4 & 5 & 1 \\ 1011 & 0010 & 0100 & 0101 & 0001 \\ \hline 2 & 6 & 2 & 2 & 2 & 1 \end{array}$

$\hookrightarrow (x)_8 = (2622121)_8$

(b)  $(354)_7 \rightarrow (y)_5$

$$\begin{array}{r} \text{(+) } 147 \\ + 35 \\ \hline 182 \\ + 4 \\ \hline 186 \end{array}$$

$$(354)_7 \rightarrow 3(7^2) + 5(7) + 4(7^0)$$

$$= 3(49) + 5(7) + 4 = 147 + 35 + 4$$

$$= (186)_{10}$$

$$\begin{array}{r} 186 \\ - 125 \\ \hline 61 \end{array}$$

$$(186)_{10} \rightarrow (y)_5 : \quad 186 - 1(125) = 61$$

$$61 - 2(25) = 11$$

$$11 - 2(5) = 1$$

$$1 - 1(1) = 0$$

$$\begin{array}{r} 61 \\ - 50 \\ \hline 11 \\ - 10 \\ \hline \end{array}$$

$(y)_5 = (1221)_5$