

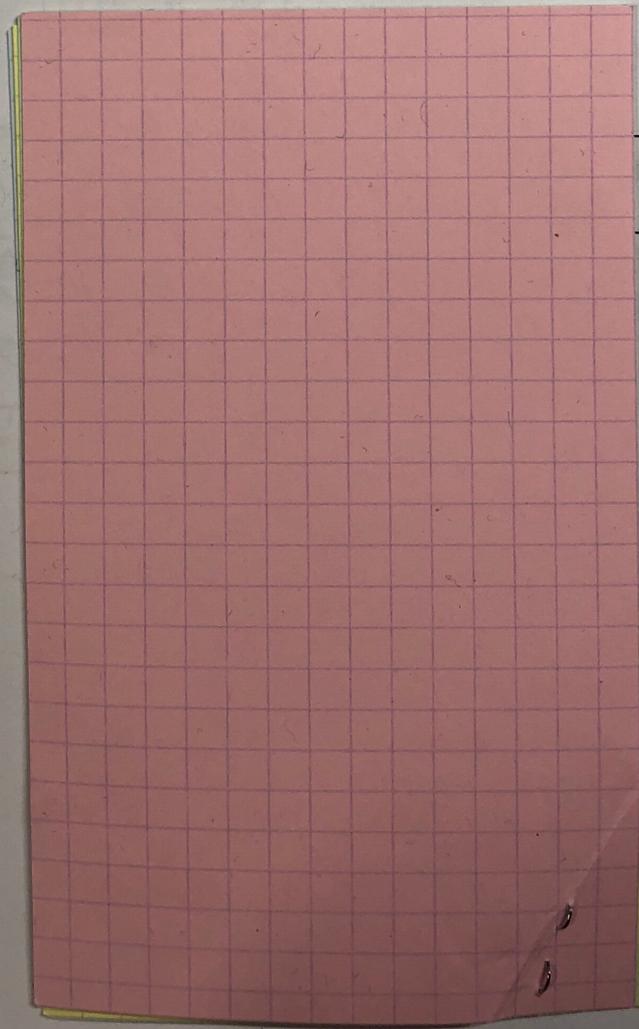
# [CS M51A FALL 18] MIDTERM EXAM

Date: 10/30/18

- The midterm is closed book, and up to 4 sheets (= 8 pages) of summary notes are allowed. You can use a calculator but not smart phones.
- For multiple choice questions, wrong answers may have a negative score value so choose carefully. It is possible that questions can have multiple answers.
- Please show all your work and write legibly, otherwise no partial credit will be given.
- **This should strictly be your own work; any form of collaboration will be penalized.**

Name :

Student ID :



Problem 1 (10 points)

1. (4 points) How many bits are required to encode a color spectrum capable of supporting 16 million using:

(a. Decimal digits in BCD

$\rightarrow (4 \times 6) \text{ million bits}$   
 $64 \text{ million bits}$   
~~X~~

BCD encode each digit (0-9) as 4 bits

b. Hexadecimal representation

~~7 digits~~  $16^6 = 16.7 \text{ million}$   
 $\therefore$  we need 7 digits in hex  
 $\Rightarrow$  Each hex = 4 bits  
 $= 7 \times 4 = 28 \text{ bits}$

Which representation is more efficient? Why?

✓ Hex is a more compact representation w/ 28 bits rather than BCD's 64m bits. Also it has 7 hex digits which is enough to encode key color properties like R, G, B, Y, alpha, gamma and within its 7 hex digits.

✓

2. (6 points) Fill in the missing entries in the table.

Radix	Digit vector $\underline{x}$	Value $x$ in decimal
16	(5, 1, 7)	1303
8	(5, 1, 7)	335
7	<u>4</u> <u>5</u> <u>3</u> <u>2</u>	1640

6

$a + b = b + a$	$ab = ba$	Commutativity
$a + (bc) = (a + b)(a + c)$	$a(b + c) = (ab) + (ac)$	Distributivity
$a + (b + c) = (a + b) + c = a + b + c$	$a(bc) = (ab)c = abc$	Associativity
$a + a = a$	$aa = a$	Idempotency
$a + a' = 1$	$aa' = 0$	Complement
$1 + a = 1$	$0a = 0$	Identity
$0 + a = a$	$1a = a$	Involution
$(a')' = a$		Absorption
$a + ab = a$	$a(a + b) = a$	Simplification
$a + a'b = a + b$	$a(a' + b) = ab$	DeMorgan's law
$(a + b)' = a'b'$	$(ab)' = a' + b'$	

Given  $E(a, b, c, d) = (ab + c)'(ac + (b' + c' + a'cd)') + a((b + c)(b + d) + c)'$ , which of the following represent the same function as  $E(a, b, c, d)$ ? Show all your work.

1.  $a + b + c + d'$
2.  $a' + b + c$
3.  $b + c' + d$

4.  $a'b'c'd$

5.  $ab'c'$

6.  $b'cd'$

$$\begin{aligned}
 E &= (ab + c)'(ac + (b' + c' + a'cd)') \\
 &= (ab + c)'(ac + b'c' + a'cd) \quad \text{--- DeMorgan's} \\
 &= (ab + c)'(ac + b'c' + a'cd) \quad \text{--- DeMorgan's} \\
 &= (ab + c)'(ac) \quad \text{--- DeMorgan's} \\
 &= (ab + c)'(ac) \quad \text{--- DeMorgan's} \\
 &= (a + b)'(a'c) \quad \text{--- DeMorgan's} \\
 &= (a' + b')(a'c) \quad \text{--- DeMorgan's} \\
 &= a'c(a' + b') \quad \text{--- DeMorgan's} \\
 &= a'c \cdot 1 \quad \text{--- DeMorgan's} \\
 &= a'c \quad \text{--- DeMorgan's}
 \end{aligned}$$

$$\begin{aligned}
 &= a((b + c)(b + d) + c)' \\
 &= a(b + c)(b + d + c)' \quad \text{--- Distributivity} \\
 &= a(b + c)' \quad \text{--- Absorption} \\
 &= ab'c' \quad \text{--- DeMorgan's} \\
 &= ab'c' \quad \text{--- DeMorgan's}
 \end{aligned}$$

Final = ① + ② =  $a + ab'c' = ab'c'$

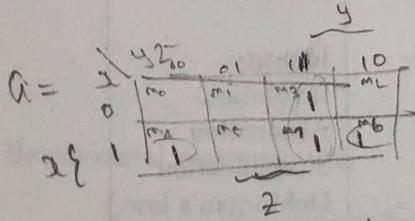
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**Problem 3 (15 points)**

Show if the gate G, described by  $G(x, y, z) = \text{one-set}\{3, 4, 6, 7\}$ , can implement NOT and AND gates. Assume that 0 and 1 are available. If it can, then use G gates to implement the following expression and show the corresponding network of G gates

$$E(a, b, c) = (a + b')(b + c')$$

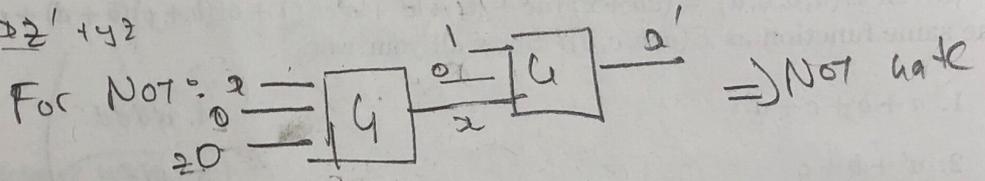
G	0	0	0
	0	0	1
	0	1	0
	1	0	0
	1	0	1
	1	1	0
	1	1	1



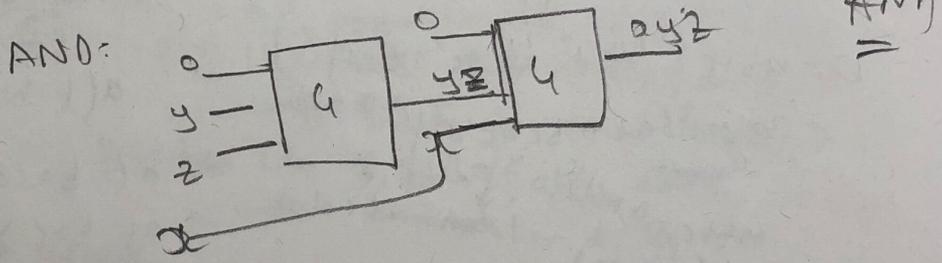
... Suitable Expression for G is  $xz' + yz$

$$G(x, y, z) = xz' + yz$$

Not:  $\text{Not}(x) = x'$   
 $G(x, 0, 0) = x$   
 $G(x, 0, 1) = 0 + x = x$   
 $G(x, 1, 0) = x$   
 $G(x, 1, 1) = x'$  ✓



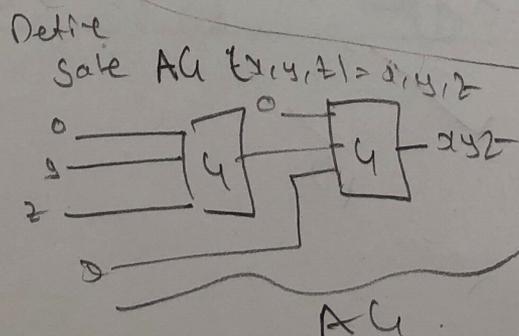
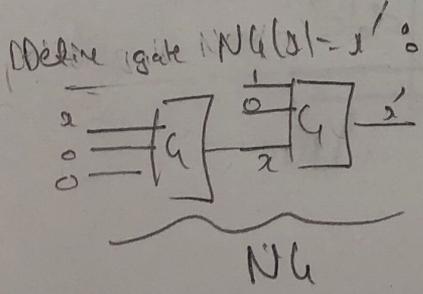
AND:  $G(0, y, z) = yz$  ✓  
 $G(1, y, z, x) = yz$



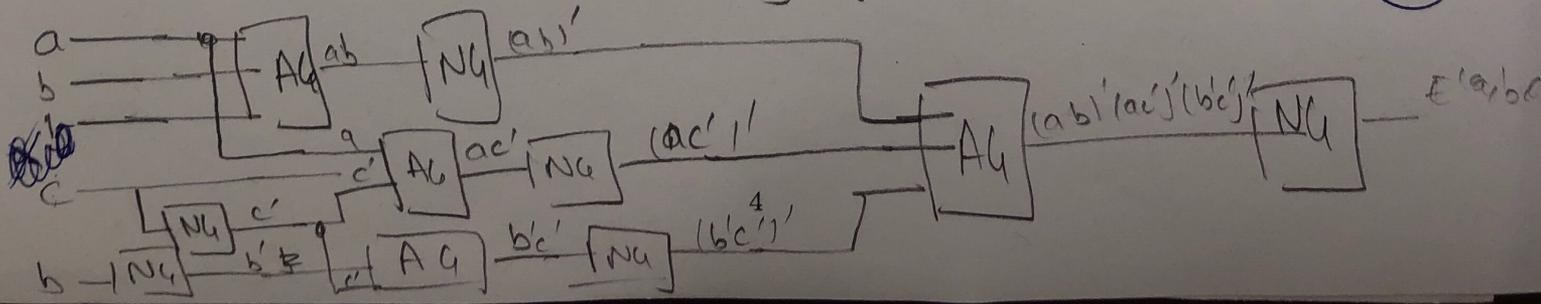
$$E(a, b, c) = (a + b')(b + c') = ab + ac' + b'c'$$

$$= ((ab)'(ac')'(b'c')')'$$

⇒ can be implemented using AND and NOT gates.  
 ⇒ can be implemented with G.



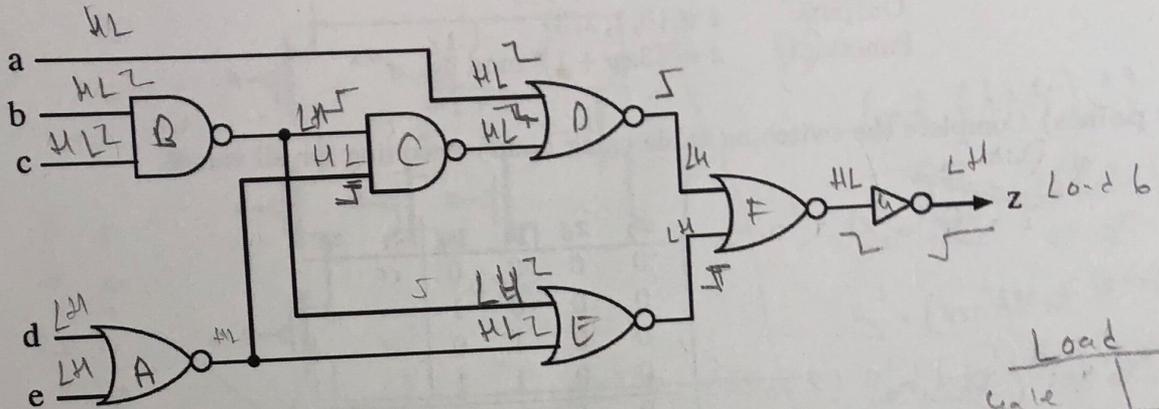
Then we can implement  $E(a, b, c)$  as follows:



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Problem 4 (15 points)

With the help of the table below, determine the low to high propagation delay  $t_{pLH}(d, z)$  of the output  $z$  of the network shown below. Assume the network output has a load of 6.



Gate Type	Fan-in	Propagation Delays (ns)		Load Factor
		$t_{pLH}$	$t_{pHL}$	
NOT	1	$0.02 + 0.038L$	$0.05 + 0.017L$	1.0
NAND	2	$0.05 + 0.038L$	$0.08 + 0.027L$	1.0
NOR	2	$0.06 + 0.075L$	$0.07 + 0.016L$	1.0

Gate	Load
A	2
B	2
C	1
D	1
E	1
F	1
z	6

$$\begin{aligned}
 &= \cancel{t_A(LH)} + t_B(LH) + t_C(HL) + t_D(LH) + t_E(LH) + t_F(HL) + t_z(LH) \\
 &= (0.07 + 0.016 * 2) + (0.05 + 0.038 * 2) + (0.08 + 0.027 * 1) \\
 &\quad + (0.06 + 0.075) + (0.06 + 0.075) + (0.07 + 0.016) \\
 &\quad + (0.02 + 0.038 * 6) \\
 &= 0.939 \text{ ns}
 \end{aligned}$$

# Problem 5 (20 points)

Obtain a two-level gate network of the following system.

Inputs:  $x, y \in \{0, 1, 2, 3\}$   
 Outputs:  $z \in \{0, 1, 2, 3\}$   
 Function:  $z = \{3xy + 1\} \pmod 4$

1. (2 points) Complete the switching table using binary encoding for all values.

*For Part 2:*

$$z_1 = x_1 x_0' y_0 + x_1 y_1' y_0 + x_1' x_0 y_1 + x_0 y_1 y_0'$$

$$= (x_1 x_0' y_0) + (x_1 y_1' y_0) + (x_1' x_0 y_1) + (x_0 y_1 y_0')$$

Always complemented input. prod.

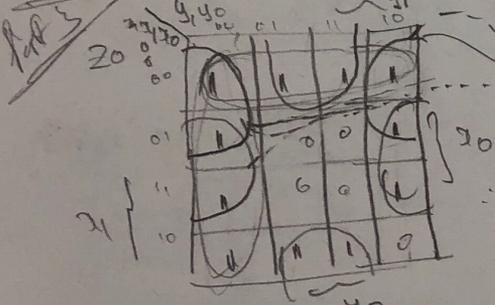
$x_1$	$x_0$	$y_1$	$y_0$	$z_1$	$z_0$
0	0	0	0	0	1
0	0	0	1	0	1
0	0	1	0	0	1
0	0	1	1	0	1
0	1	0	0	0	1
0	1	0	1	0	0
0	1	1	0	1	1
0	1	1	1	1	0
1	0	0	0	0	1
1	0	0	1	0	1
1	0	1	0	0	0
1	0	1	1	1	1
1	1	0	0	0	1
1	1	0	1	1	0
1	1	1	0	1	1
1	1	1	1	0	0

$x$	$y$	$(3xy+1)$
0	0	0+1
0	1	0+1
0	2	0+1
0	3	0+1
1	0	0+1
1	1	3+1
1	2	6+1
1	3	9+1
2	0	6+1
2	1	6+1
2	2	19
2	3	19
3	0	9+1
3	1	19
3	2	19
3	3	28

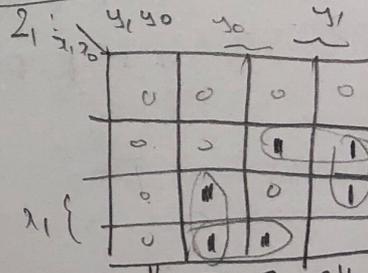
2. (5 points) Show the switching expressions of  $z_1$  and  $z_0$  in sum of minterms form.

$$z_0 = \sum m(0, 1, 2, 3, 4, 6, 8, 9, 11, 12, 14)$$

$$z_1 = \sum m(6, 7, 9, 11, 13, 14)$$



Not essential:  $\{m_0, m_2, m_4, m_6\}$   
 Essential:  $\{m_1, m_3, m_5, m_7\}, \{m_9, m_{11}, m_{13}, m_{15}\}, \{m_{10}, m_{12}, m_{14}, m_{16}\}$



3. (8 points) Show the minimal sum of products expressions of  $z_1$  and  $z_0$ . In each case, show a K-map, indicate all prime implicants, and all essential prime implicants. Show NAND-NAND networks.

Minimal SOP:

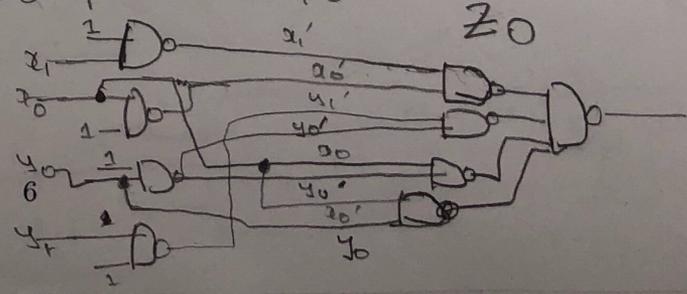
$$z_0 = x_1' x_0' + y_1' y_0' + x_0 y_0' + x_0' y_0$$

$$z_1 = x_1 x_0' y_0 + x_1 y_1' y_0 + x_1' x_0 y_1 + x_0 y_1 y_0'$$

*z1 above next part 1*

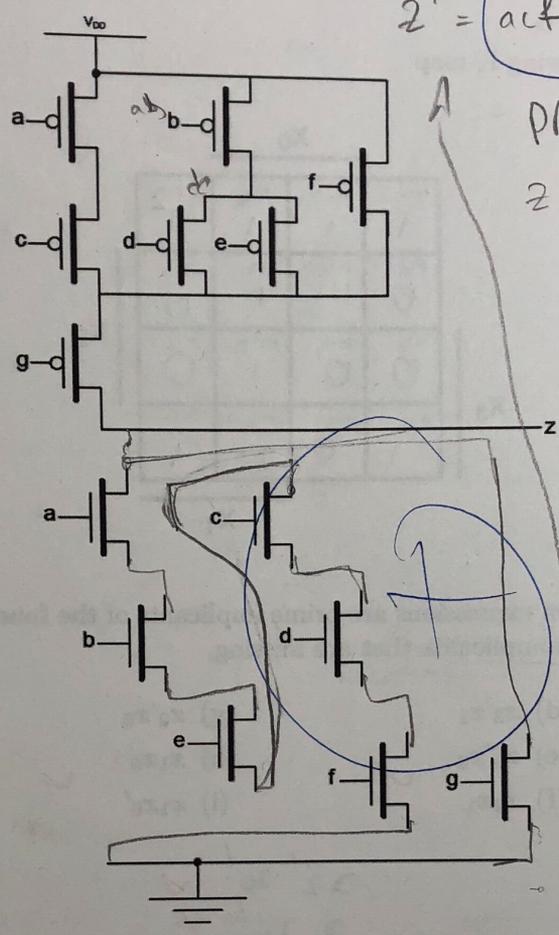
*For z1 all prime imp are ess*

$$z_0 = (x_1' x_0') (y_1' y_0') (x_0 y_0') (x_0' y_0)$$



Problem 6 (10 points)

are given the following partial CMOS network.



$$z' = acf(b+de) + g$$

PD:

$$z = (acf(b+de) + g)'$$

$$g' \cdot (acf(b+de))'$$

$$g' \cdot (act' + (b+de)')$$

$$g' \cdot ((a'+c'+t') + b'(de)')$$

$$g' \cdot ((a'+c'+t') + b'(d'te'))$$

$$g' \cdot (a'b'+c'+d'te'+t')$$

1. (5 points)

Write the expression for the pull-up network. From this, derive the expression for the pull-down network using switching algebra.

2. (5 points)

Connect NMOS transistors to complete the circuit according to the pull-down expression. Please only add missing wires.

**Problem 7 (15 points)**

For  $f(x_3, x_2, x_1, x_0) = (x_3 + x_2' + x_1 + x_0)(x_3 + x_2' + x_1' + x_0)(x_3' + x_2 + x_1 + x_0')(x_3' + x_2' + x_1 + x_0)(x_3' + x_2 + x_1 + x_0')(x_3' + x_2' + x_1' + x_0)$

1. (2 points) Fill out the following K-map.

		$x_0$				
		0	1	1	0	
	$x_3$	0	1	1	0	$x_2$
	1	0	1	1	0	
$x_3$	0	0	0	1	0	
1	0	1	1	1	0	
		$x_1$				

Prime implicants  
 $x_2'x_3' + x_0x_1'$   
 $+ x_0x_3'$   
 $+ x_2'x_0'$   
 $+ x_1x_2'$   
 $+ x_3x_2'$

2. (4 points) Which of the given expressions are prime implicants of the function given above? Circle all that apply. Write down any prime implicants that are missing.

- (a)  $x_1$
- (b)  $x_3x_1$
- (c)  $x_3'x_2'$  ✓
- (d)  $x_3'x_1$
- (e)  $x_3'x_0$  ✓
- (f)  $x_2x_1$
- (g)  $x_2'x_0$
- (h)  $x_1x_0$  ✓
- (i)  $x_1x_0'$  ✓
- (j)  $x_3'x_2'x_1$
- (k)  $x_2x_1x_0$
- (l)  $x_3x_2x_1x_0$

$x_2'x_0'$  ✓  
 $x_1x_2'$  ✗  
 $x_3x_2'$  ✗

3. (2 points) Write down the complete set of essential prime implicants.

$x_2'x_3'$  ✓  
 $x_0x_1'$  ✓  
 $x_0x_3'$  ✓  
 $x_2'x_0'$  ✓

4. (1 point) Write the minimal sum of products expression for  $f$ . Is it unique?

Minimum SOP  $f = x_2'x_3' + x_0x_1' + x_0x_3' + x_2'x_0'$   
 → Unique ✓ ✓ ✓