

[CS M51A FALL 15] MIDTERM EXAM

Date: 11/3/15

- The midterm is closed books and notes. Tablets and smartphone are not allowed.
- You can use calculators and have up to 2 sheets (= 4 pages) of summary notes.
- Please show all your work and write legibly, otherwise no partial credit will be given.
- This should strictly be your own work; any form of collaboration will be penalized.

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2C

Problem	Points	Score
1	15	15
2	15	15
3	10	9
4	15	12
5	10	10
6	15	15
7	20	20
Total	100	94

Problem 1 (15 points)

1. (6 points) Given the following simplification of a boolean expression, identify all right and wrong steps and briefly explain what is wrong for each error.

(For example, (10) \rightarrow (11) wrong application of the Identity rule, (11) \rightarrow (12) correct)

$$\begin{aligned}
 E_1(w, x, y, z) &= (((w + x + x'y')y + z)' + wx' + y')' & (1) \\
 &= ((w + x + x'y')'y'z' + wx' + y')' & (2) \\
 &= ((w + x + x'y')'y'z' + (w + y')(x' + y'))' & (3) \\
 &= ((w + x' + y')'y'z' + (w + y')(x' + y'))' & (4) \\
 &= (w'xy'z' + (w + y')(x' + y'))' & (5) \\
 &= (0 + (w + y')(x' + y'))' & (6) \\
 &= wy + xy & (7)
 \end{aligned}$$

$$\begin{aligned}
 &\stackrel{1 \rightarrow 2}{\overbrace{\left(\left((w+x+x'y')y+z \right)' + wx'+y' \right)'}} \\
 &\stackrel{2 \rightarrow 3}{\overbrace{\left((w+x+x'y')'y' \right)'}} \quad \overbrace{w+y' \overline{w'y}+} \\
 &\stackrel{6 \rightarrow 7}{\overbrace{\left((wy') (x'+y') \right)'}}
 \end{aligned}$$

$\boxed{(1) \rightarrow (2)}$ should be $(w+y')' + (x'+y')'$
 $\boxed{(2) \rightarrow (3)}$, should be $w'y + xy$
 $\boxed{(3) \rightarrow (4)}$, should be $(w+y')y'z' + (w+y')(x'+y')$, $\boxed{(6) \rightarrow (7)}$ should be $w'y + xy$

- ✓ $\boxed{(1) \rightarrow (2)}$ Wrong application of De Morgan's law
- ✓ $\boxed{(2) \rightarrow (3)}$ Correct
- ✓ $\boxed{(3) \rightarrow (4)}$ Wrong application of $x+x'y = x+y$
- ✓ $\boxed{(4) \rightarrow (5)}$ correct
- ✓ $\boxed{(5) \rightarrow (6)}$ Correct
- ✓ $\boxed{(6) \rightarrow (7)}$ Wrong application of De Morgan's law

2. (5 points) Obtain the minimal sum of products form for $E_2(w, x, y, z)$ using the identities of Boolean algebra. Show all the steps in your derivation.

$$\begin{aligned}
 E_2 &= xy' + xzw + yw \\
 &= x(y' + zw) + (x+1)yw \\
 &= x(y' + yw + zw) + yw \\
 &= x(y' + w + zw) + yw \\
 &= x(y' + w) + yw \\
 &= xy' + xw + yw \\
 &= xy' + xw(y + y') + yw \\
 &= xy' + xwy + xwy' + yw \\
 &= \boxed{xy' + yw}
 \end{aligned}$$

✓

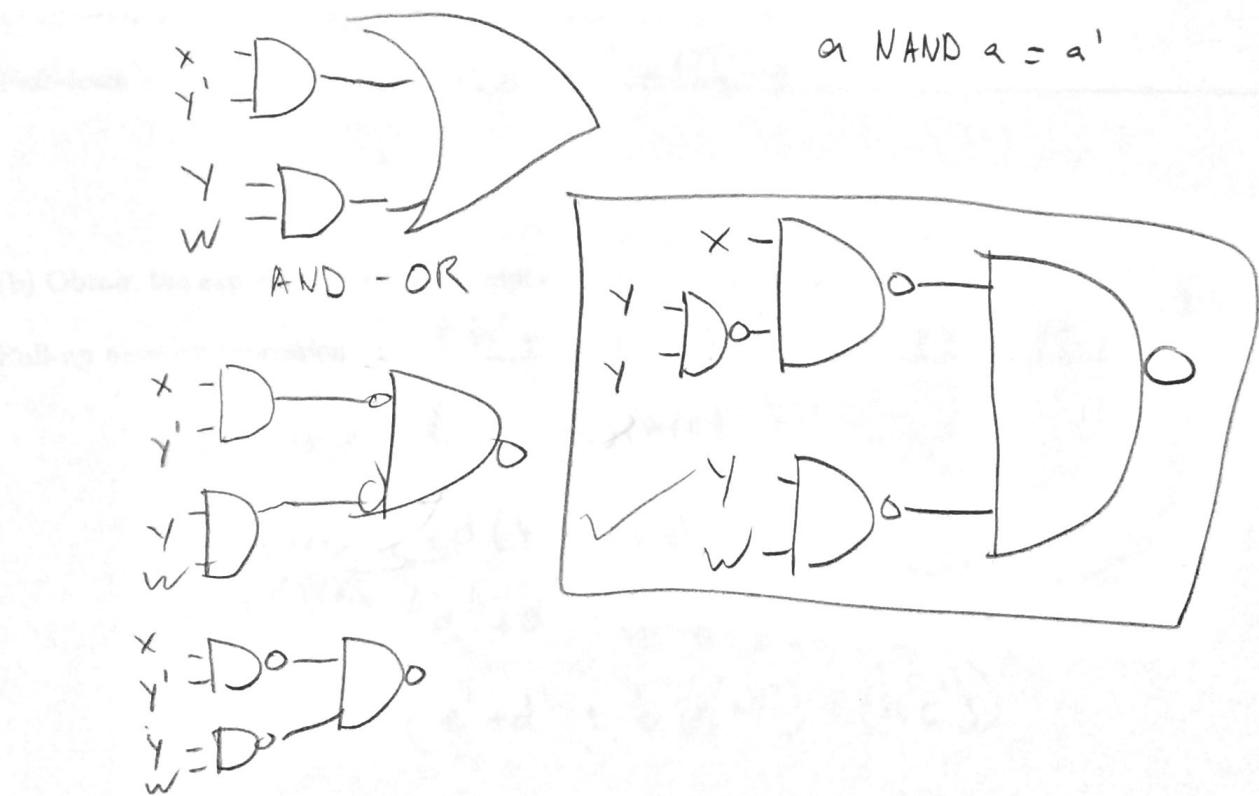
w	x	y	z	E_2
0	0	0	0	0
0	0	0	1	0
0	0	1	0	0
0	1	0	0	1
0	1	0	1	0
0	1	1	0	0
1	0	0	0	0
1	0	0	1	0
1	0	1	0	1
1	0	1	1	1
1	1	0	0	1
1	1	0	1	1
1	1	1	0	1
1	1	1	1	1
1	1	1	1	1
1	1	1	1	1
1	1	1	1	1
1	1	1	1	1
1	1	1	1	1
1	1	1	1	1

0	0	0	0	0
0	1	1	1	1
1	0	1	0	0
1	1	0	1	1
1	1	1	0	0
1	1	1	1	1

\times

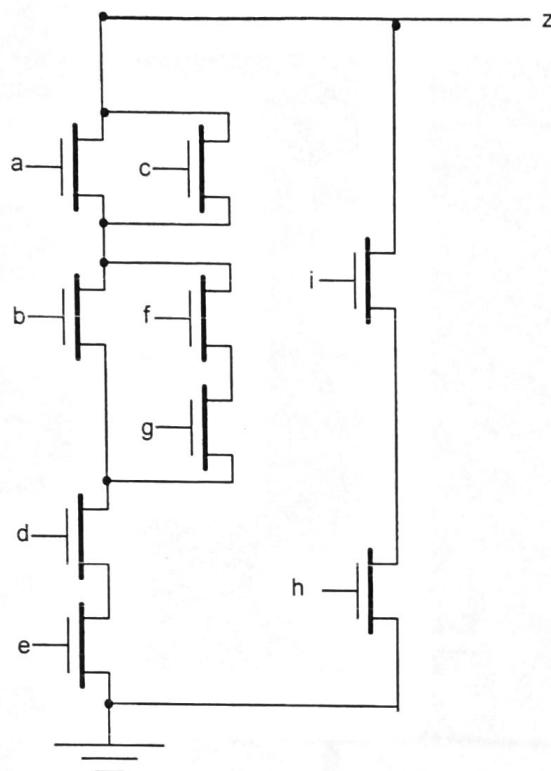
$$xy' + yw$$

3. (4 points) Using the expression obtained for E_2 from the previous step, obtain the NAND network that uses ONLY NAND gates. Inverted inputs are not available, and no constant inputs are allowed.



Problem 2 (15 points)

The following pull-down network is part of a complex CMOS gate that we want to implement.



1. (8 points) (a) Write the expression for the pull-down network.

Pull-down network expression : $z' = \underline{hi + ed(b+gf)(a+c)}$

✓

- (b) Obtain the expression for the corresponding pull-up network.

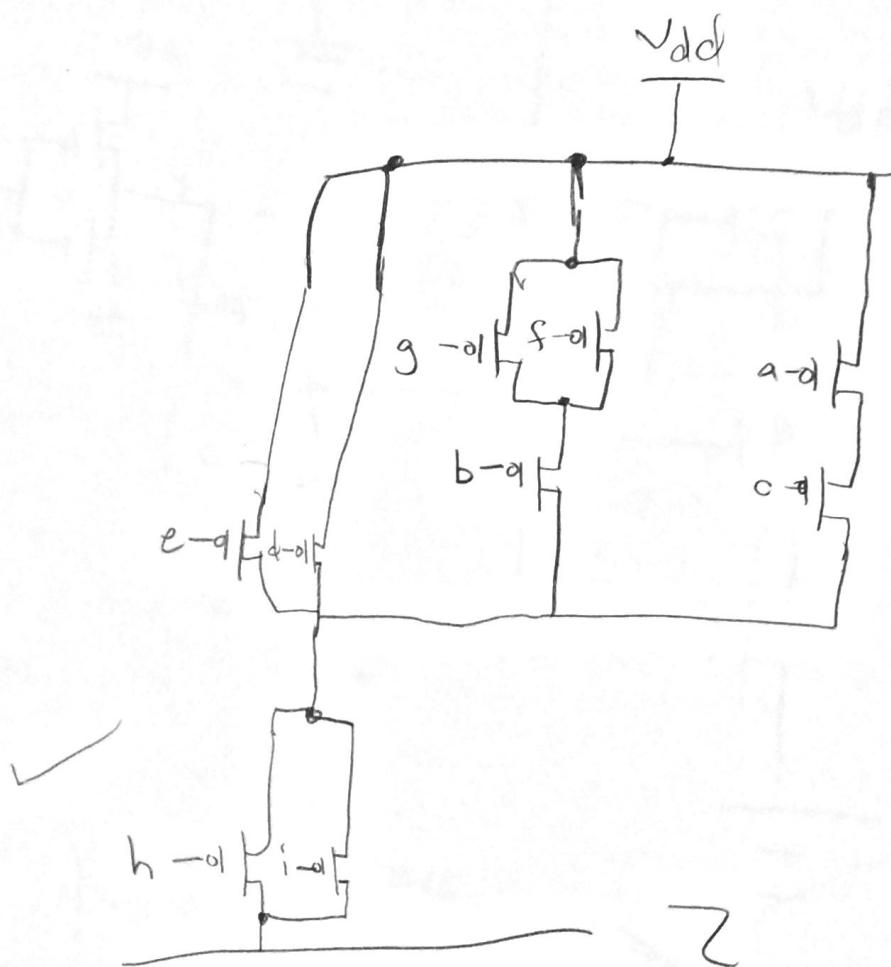
Pull-up network expression : $z = \underline{(h'i')'(e'd' + b'(g'f') + f'c')}$
 $\underline{(hi + ed(b+gf)(a+c))'}$

$(hi)'(ed(b+gf)(a+c))'$

$(h'i')'(e'd' + (b+gf)'(a+c)')$

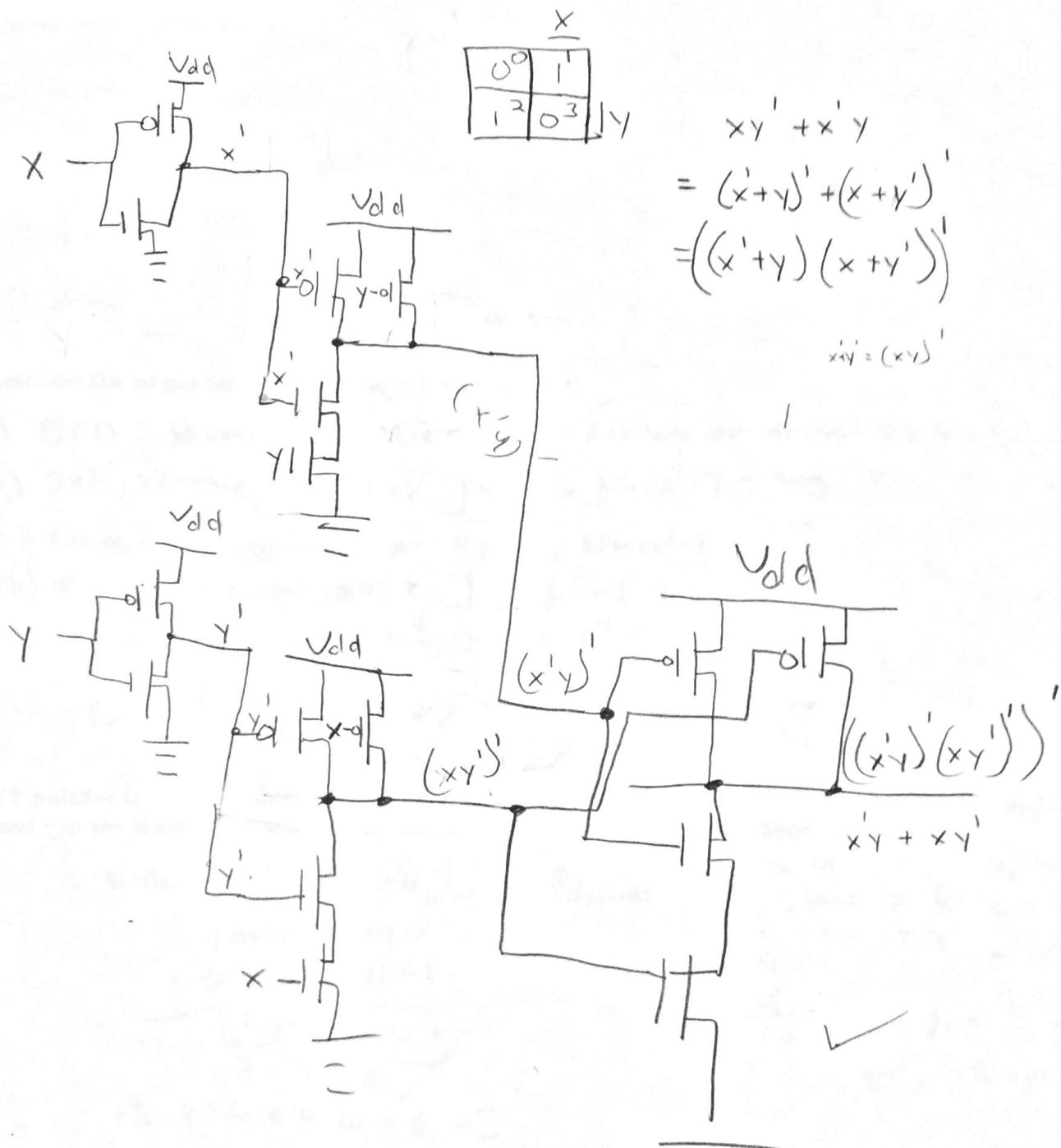
$(h'i')'(e'd' + b'(g'f') + (a'c'))$

(c) Draw the pull-up network.



2. (7 points) Draw a CMOS network that implements $f(x,y) = xy' + x'y$ (2-input XOR). x' and y' are not available as inputs. Do not use transmission gates. How many transistors does your solution have?

OPTIONAL: If you find another solution with fewer transistors, you get 3 extra points. Show your reduced design and explain why it works.



16 transistors



problem 3 (10 points)

1. (5 points) A 12-bit vector represents a set of positive integers $\{0, \dots, N\}$. Which of the following coding alternatives
- BCD
 - 2421 code (a decimal code)
 - Excess-3 code (a decimal code)
 - Octal
 - Binary

provides the largest range? Why? (Give N for each case).

- BCD = Binary Coded Decimal \Rightarrow 4 bits per decimal $\Rightarrow N = 10^3 - 1 = 999$
- 2421 \Rightarrow range of 10 per 4 bits $\Rightarrow N = 10^3 - 1 = 999$
- Excess-3 \Rightarrow range of 10 per 4 bits $\Rightarrow N = 10^3 - 1 = 999$
- Octal \Rightarrow 4 octal bits $\Rightarrow N = 8^4 - 1 = (2^3)^4 - 1 = 2^{12} - 1 = 4095$
- Binary \Rightarrow 12 binary bits $\Rightarrow N = 2^{12} - 1 = 4095$ ✓

Octal and binary have the largest range. They don't waste any space for decimal.

2. (5 points) Let $a = (1011|1001|0110)$ and $b = (0011|0111|101)$. If a represents a number in the Excess-3 code and b in the binary code, what is the value in decimal of their sum $a + b$? Show all your work.

~~$$2 a+b = \begin{array}{r} a_{\text{binary}} + b_{\text{binary}} \\ \hline 1011|1001|0110 \\ + 0011|0111|101 \\ \hline 11101001011 \end{array} - 3_{\text{decimal}}$$

$F \cdot 4 \quad B$

$$15 \cdot 256 + 4 \cdot 16 + 11 - 3$$

$$3840 + 64 + 11 - 3$$

$$3904 + 11 - 3$$

3912~~

~~$$\begin{array}{r}
 a_2 = 1011 = 8 \quad b_2 = 0011 = 3 \\
 a_1 = 1001 = 6 \quad b_1 = 1011 = B = 11 \\
 a_0 = 0110 = 3 \quad b_0 = 0101 = 5 = 5
 \end{array}$$

~~256~~

$$\begin{array}{r}
 a = 863 \quad \checkmark + 2 \\
 b = 3 \cdot 256 + 4 \cdot 11 + 5 \quad X
 \end{array}$$

$$3840$$~~

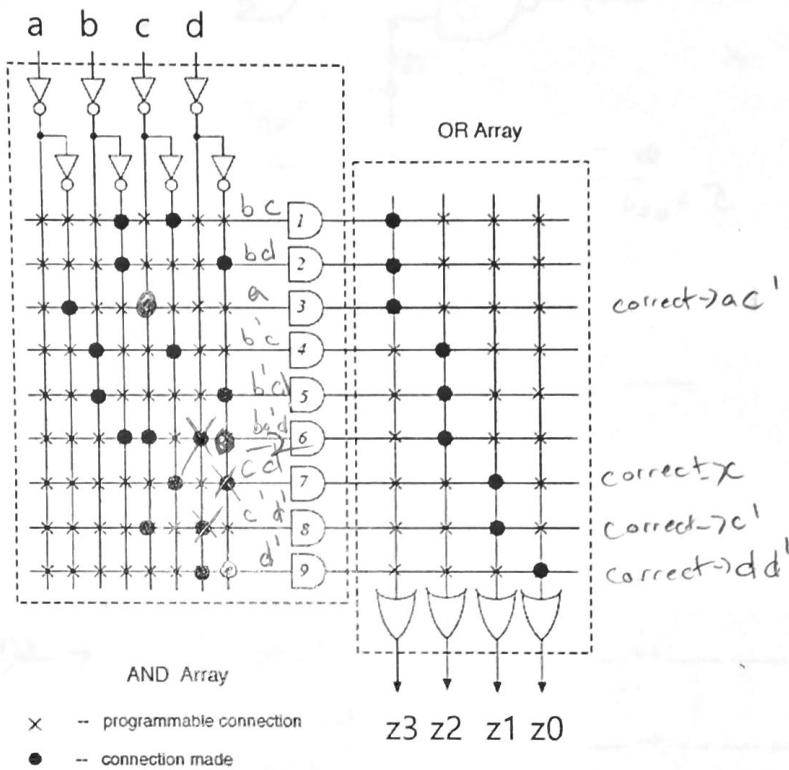
$$\begin{aligned}
 a+b &= 863 + 3 \cdot 256 + 4 \cdot 11 + 5 \\
 &= 863 + 768 + 44 + 5 \\
 &= 1631 + 44 + 5 \\
 &= 1680
 \end{aligned}$$

problem 4 (15 points)

~~12~~

We would like to verify that the PLA implementation shown here implements the following switching functions:

$$\begin{aligned}z_3 &= bc + bd + ac' \\z_2 &= b'c + b'd + bc'd \\z_1 &= 1 \\z_0 &= 0\end{aligned}$$



1. (7 points) Analyze the PLA shown above and show the output expressions.

$$z_3 = bc + bd + a$$

$$z_2 = b'c + b'd + \cancel{bc'd} - 1$$

$$z_1 = cd + c'd'$$

$$z_0 = d'$$

$$(cd)' = c'd'$$

$$z_1 = cd + c'd'$$

2. (8 points) Is the PLA implementation correct? If not, find errors and show the correct implementation (cross out wrong connections and insert correct ones)

c' on line 3 is connected

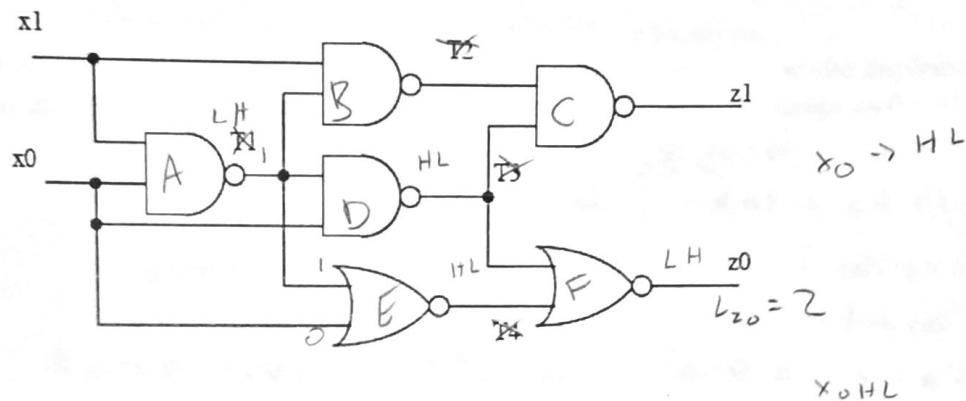
remove d on line 7

d on line 9 is connected

remove d' on line 8

Problem 5 (10 points) 10

Calculate the propagation delay $t_{PLH}(z_0)$ when x_0 changes. Assume that z_0 's load value is 2. Fill in the blanks below with the appropriate values. You don't need to fill all the blanks.



Gate Type	Fan-in	Propagation Delays (ns)		Load Factor I
		t_{PLH}	t_{PHL}	
NOT	1	$0.02 + 0.038L$	$0.05 + 0.017L$	1.0
AND	2	$0.15 + 0.037L$	$0.16 + 0.017L$	1.0
OR	2	$0.12 + 0.037L$	$0.20 + 0.019L$	1.0
NAND	2	$0.05 + 0.038L$	$0.08 + 0.027L$	1.0
NOR	2	$0.06 + 0.075L$	$0.07 + 0.016L$	1.0

A D F

Gate type & Fan-in

NAND → NAND → NOR → _____ → _____ → _____

LH / HL LH → HL → LH → _____ → _____ → _____

Output load L 3 → 2 → 2 → _____ → _____ → _____

Prop. Delay $0.05 + 0.038 \cdot 3 \rightarrow 0.08 + 0.027 \cdot 2 \rightarrow 0.06 + 0.075 \cdot 2 \rightarrow$ _____ → _____ → _____

$t_{PLH}(z_0)$ when x_0 changes is ✓

$$0.05 + 0.038 \cdot 3 +$$

$$0.08 + 0.027 \cdot 2 +$$

$$0.06 + 0.075 \cdot 2$$

Problem 6 (15 points)

A gate G is defined by the following expression

$$E(a, b, c, d) = \cancel{ad} + \cancel{a'b'c} + \cancel{a'b'd} + \cancel{bd} + \cancel{ab'd}$$

15

Show that gate G forms a universal set assuming that constants 1 and 0 are available.

Specify a pre-established universal set you are using in the proof, and explicitly show the implementation for each element in the set using gate G with 1 and/or 0 as needed. For example, you can assign $a=0$ and $b=1$ in the expression E.

$$\begin{aligned} E(a, b, 0, 1) &= a'b' + ab' \\ &= b' = \text{NOT} & E(a=0) &= cd + b'c' + b'd + bcd' \\ &&& a=1 = cd + bcd' + b'c'd \\ E(0, b, 0, d) &= b' + b'd' & b=0 &= cd + a'c' + a'd' + ac'd' \\ &= b' & b=1 &= cd + cd' = c \\ &&& c=0 = a'b' + a'b'd' + ab'd \\ &&& c=0 \} = a'b' + a'b' + \\ &&& d=0 \} = a'b' + a'b' + \end{aligned}$$

El.

$$E(a, b, 0, 0) = a'b' = \text{NOR } b$$

NOR is universal

$$a'b' = (a+b)'$$

problem 7 (20 points)

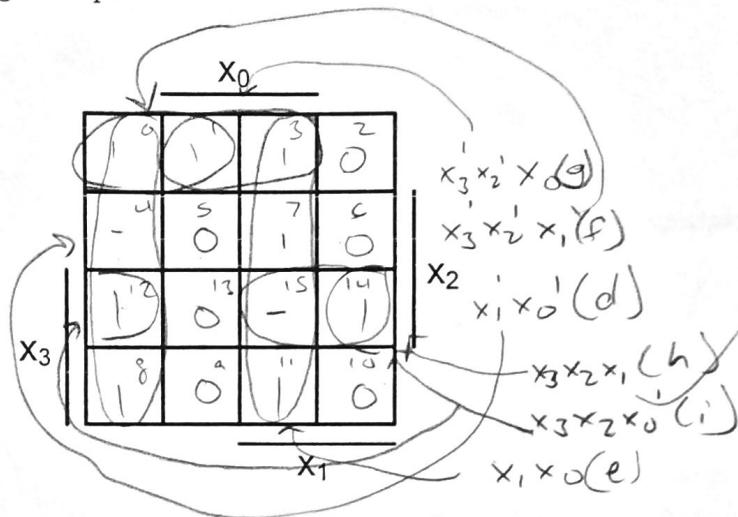
for the switching function $f(x_3, x_2, x_1, x_0)$, we are given the information below for the dc-set and zero-set.

dc-set = (4, 15)

zero-set = zero-set of function

$$(x_3 + x_2 + x_1' + x_0)(x_3 + x_2' + x_1 + x_0')(x_3 + x_2' + x_1' + x_0)(x_3' + x_2 + x_1 + x_0')(x_3' + x_2 + x_1' + x_0)(x_3' + x_2' + x_1 + x_0')$$

1. (2 points) Fill out the following K-map.



2. (4 points) Which of the given expressions are prime implicants of the function given above? Circle all that apply. Do not circle implicants that are not prime.

- (a) x_3x_1
- (b) x_3x_2'
- (c) $x_3'x_1$

- (d) $x_1'x_0'$
- (e) x_1x_0
- (f) $x_3'x_2'x_1'$

- (g) $x_3'x_2'x_0$
- (h) $x_3x_2x_1$
- (i) $x_3x_2x_0'$

- (j) $x_3'x_2x_1'x_0$
- (k) $x_3x_2'x_1x_0$
- (l) $x_3'x_2x_1x_0'$

3. (3 points) Write down the complete set of essential prime implicants.

$$x_1'x_0 \quad x_1x_0$$

4. (2 points) Write down the minimal sum of products expressions for f . If there are multiple forms of minimal sum of products expressions, you only need to write down one of them.

$$x_1'x_0 + x_1x_0 + x_3x_2x_0' + x_3'x_2x_0$$

5. (4 points) Which of the given expressions are prime implicants of the function given above? Circle all that apply. Do not circle implicants that are not prime.

(a) $(x_3 + x_2' + x_1)$

(b) $(x_3' + x_2')$

(c) $(\cancel{x_2'} + x_1 + x_0')$

(d) $(x_3' + x_1 + x_0')$

(e) $(x_3 + x_2' + x_0)$

(f) $(x_3 + x_1' + x_0)$

(g) $(x_2 + x_1' + x_0)$

(h) $(x_3 + x_1 + x_0)$

(i) $(x_3' + x_2' + x_1)$

(j) $(x_3' + x_1')$

(k) $(x_3 + x_2 + x_1 + x_0')$

(l) $(x_3 + x_2' + x_1' + x_0)$

6. (3 points) Write down the complete set of essential prime implicants.

$$\begin{array}{c} x_2 \\ \times \quad + x_1 + x_0 \\ x_3 \\ \times \quad + x_1 + x_0 \\ x_2 \\ \times \quad + x_1 + x_0 \\ x_3 \\ \times \quad + x_1 + x_0 \end{array}$$

7. (2 points) Write down the minimal product of sums expressions for f . If there are multiple forms of minimal product of sums expressions, you only need to write down one of them.

$$(x_2 + x_1 + x_0)(x_3 + x_1 + x_0)(x_2 + x_1 + x_0)(x_3 + x_1 + x_0)$$

