

[CS M51A FALL 14] MIDTERM EXAM

Date: 11/04/14

- Please show all your work and write legibly, otherwise no partial credit will be given.
- This should strictly be your own work; any form of collaboration will be penalized.

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Problem	Points	Score
1	10	10
2	10	10
3	20	20
4	20	20 10
5	10	10
6	30	28
Total	100	100 98

Problem 1 (10 points)

10

$X = (x, y, z)$ is a 3-digit weighted mixed-radix number system: x is a radix-16 digit, y is a radix-3 digit, and z is a radix-12 digit.

1. (5 points) Convert $X = (6, 1, 11)$ to a decimal number.

$$= 6(3)(12) + 1(12) + 11$$

$$= 6(36) + 12 + 11$$

$$= 216 + 23$$

$$= \boxed{239}_{10} \checkmark$$

$$\begin{array}{r} 3 \\ 36 \\ \hline 6 \\ 216 \end{array}$$

$$\begin{array}{r} 12 \\ 11 \\ \hline 23 \end{array}$$

$$\begin{array}{r} 216 \\ 23 \\ \hline 239 \end{array}$$

2. (5 points) What is the largest number of X in decimal?

$$\text{largest} = \underline{(15, 2, 11)}$$

$$= 15(3)(12) + 2(12) + 11$$

$$= 15(36) + 24 + 11$$

$$= \boxed{15(36) + 35} \checkmark$$

$$111$$

$$1(2)(2) + 1(2) + 1$$

Problem 2 (10 points)

Simplify the following boolean expression by using postulates of Boolean Algebra.

$$(a'b' + c)(a + b)(b' + a'c)'$$

$$= (a'b'a + ac + bc)(b' + a'c)'$$

distributive

$$= (0 + ac + bc)((b)(a + c))$$

De Morgan's, complement

$$= (ac + bc)(ab + bc)$$

distributive

$$= abc + abc + abc + bc$$

distributive

$$= abc + bc$$

idempotent

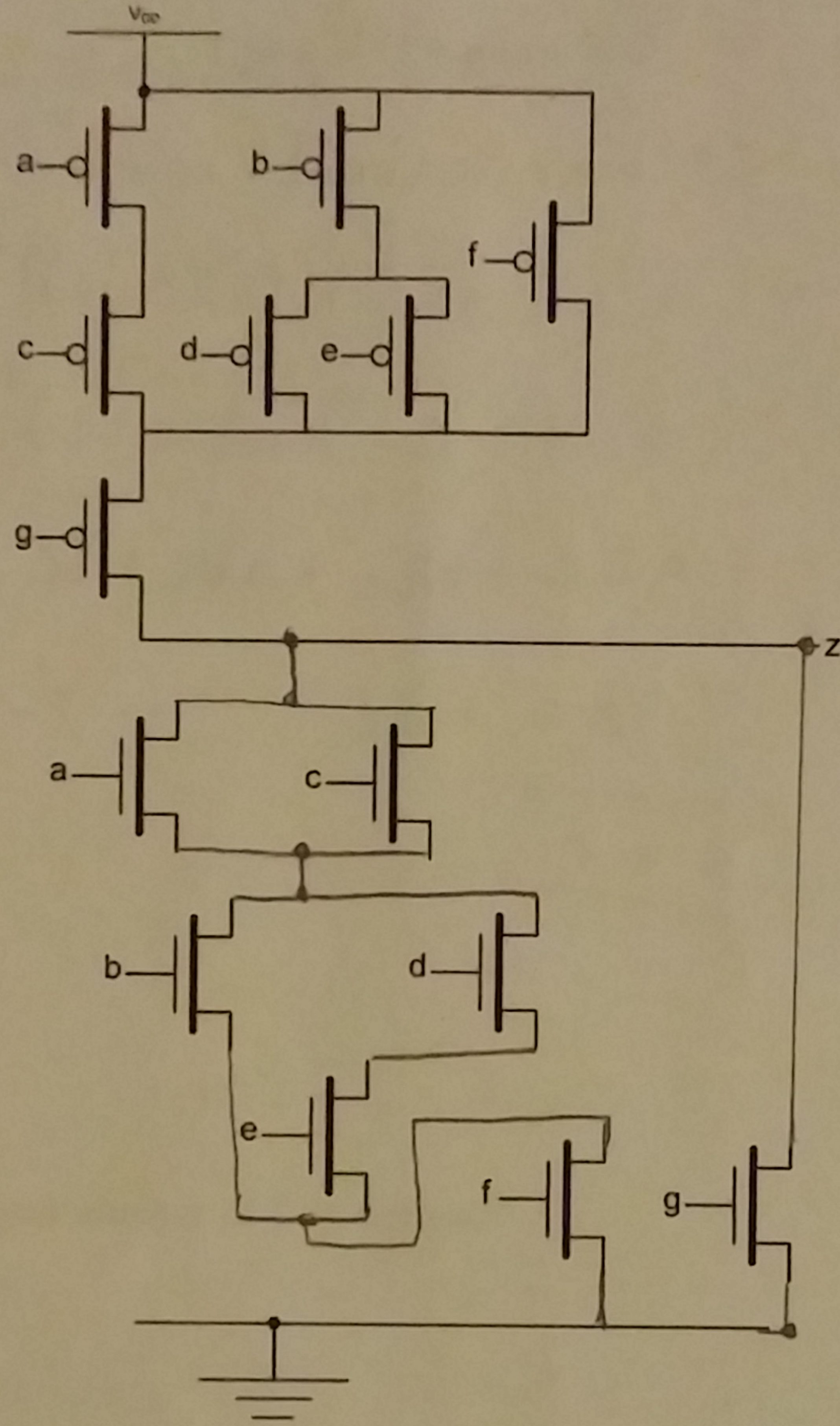
$$\boxed{= bc}$$

absorption

/D

Problem 3 (20 points)

We are given the following partial CMOS network.



1. (10 points) Write the expression for the pull-up network. From this, derive the expression for the pull-down network using switching algebra.

$$Z = (g') ((a'c') + (b'(d'+e')) + (f'))$$

$$[b'(d'+e')]'$$

$$\text{Pull Down} = Z'$$

$$b + de$$

$$Z = g' (a'c' + b'(d'+e') + f')$$

$$Z' = (g) + (a'c' + b'(d'+e') + f')' \quad \text{DeMorgan's}$$

$$= g + (a + e) (b + de) (f)$$

Pull-Down Network Expression

b

2. (10 points) Connect the NMOS transistors in page 4 to complete the pull-down network so that it corresponds to the expression obtained from part 1 and drives the output z to a valid output - i. e. either V_{DD} or ground - for any combination of inputs.

You can directly draw on the CMOS implementation on page 4.

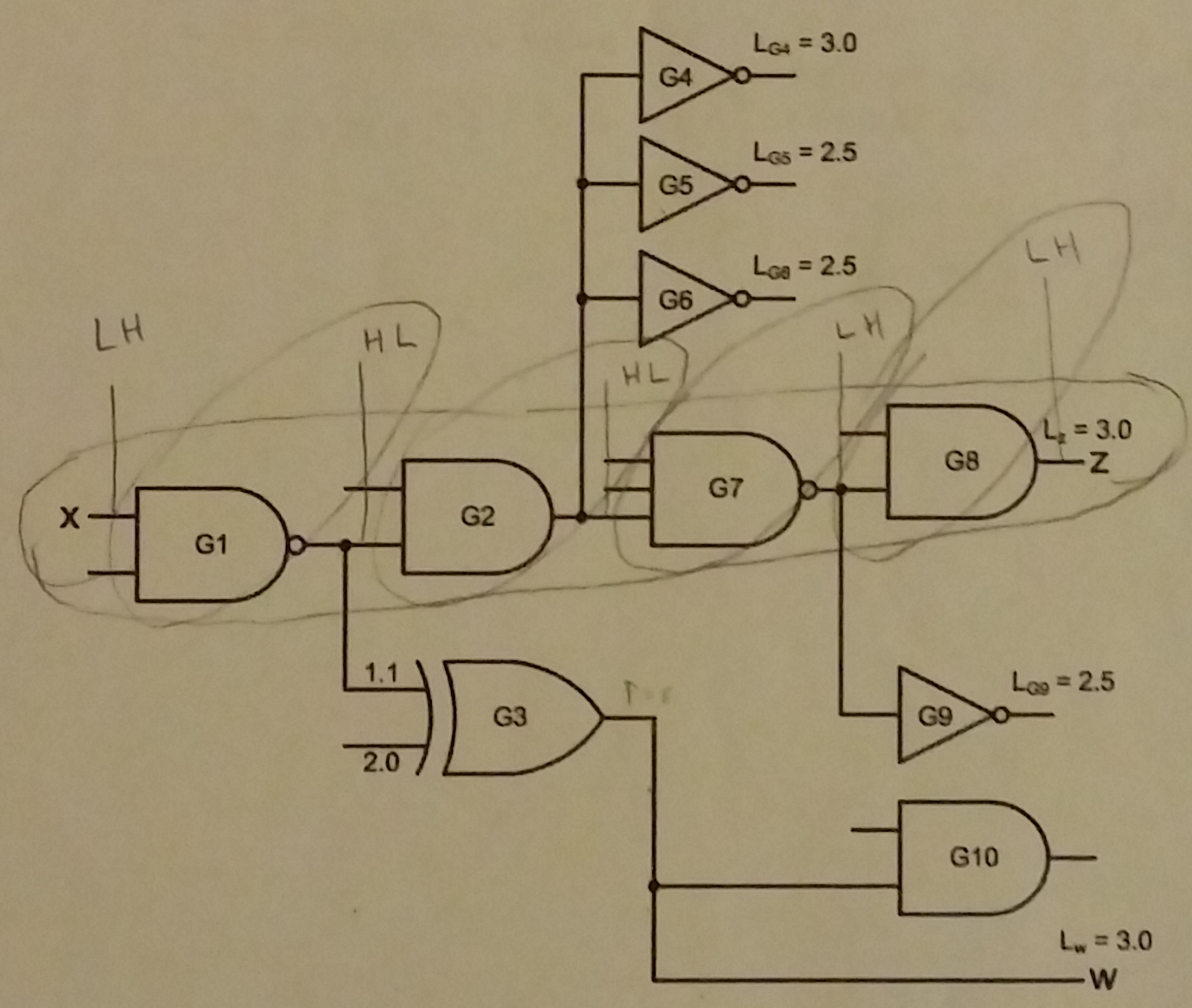
Drawn on pg. 4

Problem 4 (20 points)

#5 15

We would like to determine the propagation delay of the gate network shown here. The output is z with a input x . The necessary gate characteristics are given in the table below.

Gate Type	Fan-in	Propagation Delays (ns)		Load Factor I
		t_{pLH}	t_{pHL}	
NOT	1	$0.02 + 0.038L$	$0.05 + 0.017L$	1.0
AND	2	$0.15 + 0.037L$	$0.16 + 0.017L$	1.0
NAND	2	$0.05 + 0.038L$	$0.08 + 0.027L$	1.0
NAND	3	$0.07 + 0.038L$	$0.09 + 0.039L$	1.0
XOR	2	$0.30 + 0.036L$	$0.30 + 0.021L$	1.1
		$0.16 + 0.036L$	$0.15 + 0.020L$	2.0



2.1 4 2 3

70
1. (10 points) (a) Determine the output load of gate G3.

(b) If the fanout factor of gate G3 is 8, then how many additional gate inputs with load factor of 2 can be connected to the output of gate G3?

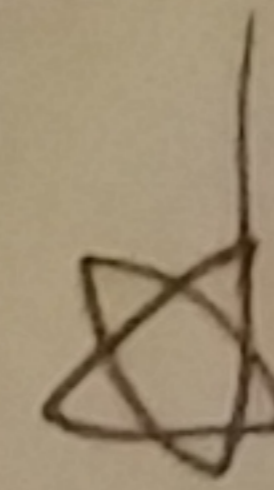
$$= 8.0 - L(G_{10}) - L_w$$

$$= 8.0 - 1 - 3$$

$$= 8.0 - 4$$

$$= \underline{4.0} + 2$$

(a) ?



additional inputs of $I=2$

$$\frac{4.0}{2.0} = \boxed{2}$$

10

2. (10 points) Find the worst case value of $t_{pLH}(x \rightarrow z)$. Fill in the blanks below with the appropriate values.

Gate type and fan-in G1: NAND 2 → G2: AND 2 → G7: NAND 3 → G8: AND 2

LH / HL G1: HL → G2: HL → G7: LH → G8: LH

Output load L G1: 2.1 → G2: 4 → G7: 2 → G8: 3
0.08 + 0.16 + 0.07 + 0.15 +

Propagational delay G1: 0.027(2.1) → G2: 0.017(4) → G7: 0.038(2) → G8: 0.037(3)

$t_{pLH}(x \rightarrow z)$ ↓

$$= 0.08 + 0.16 + 0.07 + 0.15 +$$

$$0.027(2.1) + 0.017(4) + 0.038(2) + 0.037(3)$$

10

Problem 5 (10 points)

Decide whether the function E defined in (a) and (b) are universal functions.. If Yes, show your proof. You can use constant 0 or 1 as your input.

(a) (5 points)

x	y	$E(x, y)$
0	0	1
0	1	1
1	0	0
1	1	1

$$E(x, y) = (x' \wedge y)$$

$$E(x, y) = x' y' + x' y + x y$$

$$= x' + y$$

* NOT-CASE

$$E(x, 0) = x'(1) + x'(0) + x(0) = x' = \text{NOT}(x) \checkmark$$

* OR-CASE

$$E(E(x, 0), y) = ((x' + 0)' + y) = ((x')' + y)$$

$$= (x + y)$$

$$= \text{OR}(x, y)$$

OR-NOT is a known universal set, and I have shown both can be constructed through E

Therefore E is also universal.

(b) (5 points)

$$x'yz + xy'z' + xyz' + xyz$$

$$x'yz + x(y'z' + yz' + yz)$$

$$x'yz + xz' + xy$$

x	y	z	E(x,y,z)
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	1

$$E(x,y,z) =$$

$$x'yz +$$

$$xy'z' +$$

$$xyz' +$$

$$xyz$$

Not Universal

because there is no input that can generate an inversion with this setup.

For any two zeros the output will not invert the third ones

Therefore Not, NAND, NOR are Not Possible

$$E(x, 1, y) = x'y + xy' + xy = x + y = OR(x, y)$$

$$E(x, y, 1) = x'y + y + xy = xy = AND(x, y)$$

Therefore E is Universal

because it can form OR and AND

which are known to be universal

28 Problem 6 (30 points)

F is a function that accepts inputs $x \in \{0, 1, 2\}$, $y \in \{0, 1, 2, 3\}$, and outputs $z = x^2 + y$. Suppose you use binary code to encode x , y , and z . x is encoded as x_1x_0 , y is encoded as y_1y_0 , z is encoded as $z_2z_1z_0$.

(a) (5 points) Fill in the following truth table.

x	y	$x^2 + y$
0	0	0
0	1	1
0	2	2
0	3	3
1	0	1
1	1	2
1	2	3
1	3	4
2	0	4
2	1	5
2	2	6
2	3	7

	x_1	x_0	y_1	y_0	z_2	z_1	z_0
0	0	0	0	0	0	0	0
	0	0	0	1	0	0	1
	0	0	1	0	0	1	0
3	0	0	1	1	0	1	1
4	0	1	0	0	0	0	1
	0	1	0	1	0	1	0
	0	1	1	0	0	1	1
7	0	1	1	1	1	0	0
8	1	0	0	0	1	0	0
	1	0	0	1	1	0	1
	1	0	1	0	1	1	0
11	1	0	1	1	1	1	1
12	1	1	0	0	-	-	-
	1	1	0	1	-	-	-
	1	1	1	0	-	-	-
15	1	1	1	1	-	-	-

$x = 3 \rightarrow \{-\}$

(b) (5 points) Based on the truth table in (a), draw the K-map for z_2 , z_1 , and z_0 .

z_2 :

	y_1	y_0	00	01	11	10
x_1	x_0		y_0			
00			0	0	0	0
01			0	0	1	0
11	x_0		-	-	-	-
10		x_1	1	1	1	1
			y_1			

$P_1 = x_1 x_0'$ ✓
 $P_2 = y_1 y_0 x_0$ ✓

essential

z_1 :

0	0	1	1
0	1	0	1
-	-	-	-
0	0	1	1

$P_1 = y_0' y_1$ ✓
 $P_2 = y_1 x_0'$ ✓
 $P_3 = y_0 y_1' x_0$ ✓

z_0 :

0	1	1	0
1	0	0	1
-	-	-	-
0	1	1	0

$P_1 = x_0' y_0$ ✓
 $P_2 = y_0' x_0$ ✓

(c) (5 points) Use the K-maps in (b) to find all the prime implicants for z_2 , z_1 , and z_0 respectively

$$z_2: \quad p_1 = x_1 x_0'$$

$$p_2 = y_1 y_0 x_0$$

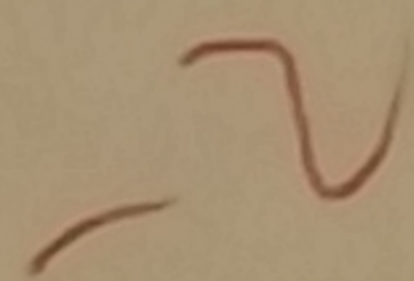
$$z_1: \quad p_1 = y_0' y_1$$

$$p_2 = y_1 x_0'$$

$$p_3 = y_0 y_1' x_0$$

$$z_0: \quad p_1 = x_0' y_0$$

$$p_2 = y_0' x_0$$



(d) (5 points) Use the K-maps in (b) to find all the essential prime implicants for z_2 , z_1 , and z_0 respectively.

$$z_2: \quad e_1 = x_1 x_0'$$

$$e_2 = y_1 y_0 x_0$$

$$z_1: \quad e_1 = y_0' y_1$$

$$e_2 = y_1 x_0'$$

$$e_3 = y_0 y_1' x_0$$

$$z_0: \quad e_1 = x_0' y_0$$

$$e_2 = y_0' x_0$$

(e) (5 points) Give minimal SOP expression for z_2 , z_1 , and z_0 . State if the expressions are unique.

$$z_2 = x_1 x_0' + y_1 y_0 x_0$$

$$z_1 = y_0' y_1 + y_1 x_0' + y_0 y_1' x_0$$

$$z_0 = x_0' y_0 + y_0' x_0$$

All 3 are unique

(f) (5 points) Implement z_2 , z_1 , and z_0 using minimal NAND-NAND networks. Note that each output has a separate gate network. You can directly use x_1, x_0, y_1, y_0 as inputs.

