CS33: Intro Computer Organization Fall 2021 Midterm

Name: \_\_\_\_\_

UID: \_\_\_\_\_

(Please write as legibly as possible!)

This is an open book, open notes exam, but you cannot share books/notes. Please follow the university guidelines in reporting academic misconduct.

Note: If the architecture of the machine is not specified, assume that the question is being asked in the context of a 64-bit little endian x86 machine, running Linux.

Please wait until everyone has their exam to begin. We will let you know when to start.

Good luck!

		Points Possible
1.	Multiple Choice	14
2.	Baba is Struct	8
3.	Notable Floats	5
4.	Array Interpretation	8
5.	Imaginary Stack Allocator	10
6.	Tricky Switch	7
7.	Floating Point Mystery	4
8.	Two Birds, One Bomb	6
Total		50 +12 bonus points

## **Question 1. Multiple Choice (14 points)**

For the following multiple choice questions, **<u>select all answers that apply.</u>** If none are correct, leave the question blank. Put your answers in the table on the next page. (2pts each, no partial)

- 1. What is a difference between unsigned and signed integer representations?
  - a. Unsigned integers can store a wider range for the same number of bits.
  - b. Right-shifting an unsigned integer uses "logical" shift, while right-shifting a signed integer uses "arithmetic" shift.
  - c. In a C expression that operates on two different datatypes, an unsigned datatype will take precedence over a floating-point datatype, but a signed datatype will not.
  - d. It is meaningful to ask if a signed number is greater-than or equal to zero, while the same is meaningless for unsigned numbers.
- 2. Suppose you use the objdump command to disassemble a function, and you see this:

000000000400595 <func>: 400595: 53 push %rbx

What does the "53" in the above line represent?

- a. The function has a length (in terms of instructions) of 53 bytes.
- b. The push instruction is 53-bytes offset from the beginning of the function.
- c. "53" contains an encoding of the register %rbx.
- d. "53" contains an encoding of the push operation.
- e. Address 53 is the location of where a callee-saved register is pushed.
- 3. X86\_64 contains an instruction which performs a conditional move -- **cmov** --, which moves one register to another based on the condition codes (aka. flags). This instruction can sometimes be used to perform if-statement control flow. When is performing if statements using conditional moves a better option than using ordinary branch instructions?
  - a. When the body of the if statement contains function calls.
  - b. When the body of the if statement contains side effects.
  - c. When the body of the if statement contains many instructions.
- 4. X86\_64 contains an instruction which performs an indirect jump -- **jmp\*** --, which jumps to a location specified in a memory table. This instruction can sometimes be used to perform switch-case-statement control flow. When is performing switch statements using indirect jumps a better option than using ordinary branch instructions?
  - a. When there are more fall-through cases than non-fall-through cases.
  - b. When there are few instructions in the case statements.
  - c. When there are many instructions in the case statements.
  - d. When the first case value is zero.
  - e. When the range of case values is contiguous.

- 5. Suppose that the next C language standard contains an 11-bit unsigned integer datatype. What would be a valid reason (or reasons) to reject this proposal?
  - a. There are already larger and smaller datatypes (e.g. 8-bit and 16-bit), so it's not useful.
  - b. Modern ISA use byte-addressable memory, so accessing arrays of contiguous 11-bit integers would require extra instructions to extract the number.
  - c. For a two's complement number to be well-defined, its size must be a multiple of 2.
  - d. Having an 11-bit number would make it impossible to satisfy the datatype casting rules in the C standard.
- 6. For what C datatypes is the concept of "endianness" irrelevant?
  - a. char
  - b. unsigned char
  - c. string (i.e. array of char)
  - d. int
  - e. float

# 7. What kind of control flow is contained in this assembly function? (see figure below)

- a. Loop (e.g. while/for)
- b. Conditional Branch (if/else)
- c. Indirect Branch (switch/case)

000000000000000000000000000000000000000			900	<func>:</func>	
0:	89	fØ		mov	%esi,%eax
2:	0f	bf	d0	movswl	<u>%ax,%edx</u>
5:	8d	0c	3a	lea	(%rdx,%rdi,1),%ecx
8:	83	f9	02	cmp	<pre>\$0x2,%ecx</pre>
b:	7e	0b		jle	18 < <u>func</u> +0x18>
d :	39	fa		cmp	%edi,%edx
f:	7f	04		jg	15 < <u>func</u> +0x15>
11:	01	ff		add	%edi,%edi
13:	<u>eb</u>	ed		jmp	2 < <u>func</u> +0x2>
15:	89	f8		mov	%edi,%eax
17:	c3			retq	

Answer Table (list any correct answers)

1	
2	
3	
4	
5	
6	
7	

## Question 2. Baba is Struct (8 Pts)

Test is

Consider the following structure, union, and array definitions:

<pre>typedef struct {     int baba;     short flag[5];     float* keke;     char key; } noun;</pre>	<pre>// Note: Typedef here just means that were are defining a struct type that we can use later in the array definition.</pre>
<pre>typedef union {     int hot;     short shut[5];     float* stop;     char open; } property;</pre>	
<pre>noun you[5]; property me[5];</pre>	<pre>// Create an "arrays of structs" and an "array of unions"</pre>

- 1. If &you == 0 (i.e. if the address of you[0] is zero), at what address is you[1]?
- 2. If &me == 0 (i.e. if the address of me[0] is zero), at what address is me[1]?
- 3. If we access "property.shut[1]", that will also access half of the bits in "property.hot". Will that be accessing the least significant bits of "hot" or the most significant?
- 4. Consider all the primitive data types within the arrays "you" and "me". Which of these are guaranteed to have their addresses aligned to a multiple of their size?

# **Question 3. Notable Floats (5 Pts)**

The following table shows a number of "interesting" values for floating-point numbers, along with their encoding into sign, exp, and frac fields.

Pattern	sign	ехр	frac
А	0	0000	0000
В	0	0000	0001
С	0	0001	1111
D	0	0001	0000
E	0	0111	0000
F	0	1110	1110
G	0	1111	0000
Н	0	1111	1111

Match the following definitions to the interesting numbers above (write "A", "B", etc. in the box)

0. Zero	
1. One	
2. Smallest Non-zero Denorm	
3. Smallest Possible Normalized Number	
4. Not a Number (NaN)	

Note: These patterns don't specify the number of bits, but that won't matter for answering the question. Assume the same strategy for representation of denormalized numbers, NaNs and infinities as other IEEE 754 standard numbers.

#### Question 4. Array Interpretation (8 pts)

- 1. For the following datatype definitions, answer the following questions: (6pts)
- **sizeof(array):** What is the size of "array" in terms of number of bytes? For this problem, we are only talking about the memory allocated for the variable "array", and not any other supporting data structures.
- **How many dereferences?:** List the number of memory dereferences that it would take to access an integer for that datatype. In other words, how many times do you have to access memory *total* (how many loads), to eventually access a single integer. Include the load of the integer itself.

	Array Declaration	sizeof(array)	How many dereferences?
	int* array	8	1
1	int array[3][2]		
2	int (*array)[5]		
3	int *array[5]		

2. What is the array declaration for the following array, represented visually below? (2pts)

array		•			
Allocated pointe	Unalloca er to unalloc	ited int ated int		 •	→
Datatype for array:					

# **Question 5. Imaginary Stack Allocator (10pts)**

**Part 1:** Suppose we have two functions, FuncP and FuncQ. FuncP calls FuncQ, and the stack frames of both functions are depicted below.

All functions require a "prologue" and "epilogue" to manage the stack. The prologue allocates stack space, and usually appears at the beginning of the function. The epilogue deallocates stack space, and usually appears at the end of the function.

Based on the stack frame, write the prologue and epilogue for FuncQ. Don't use more instructions than you need to.

		Stack Frames	
Prologue for FuncQ:		Arg. 8	F
		Arg. 7	u n
		Return Address	c P
		Old r12	_
		Old rbp	r u
Epilogue for FuncQ:		Unused (16 bytes)	n C Q
		Space for C array int array[4]	
	rsp->		

**Part 2:** FuncR is another function, and its assembly is shown below. Fill in the prologue and epilogue for this function too!

%rbx
3(%rsp)
%esi
), %rdi
%rax
<pre>icR_epilogue</pre>

Question 6. Tricky Switch (7 pts)

Source Code	Compiled Assembly			
<pre>int func(int x, int y, int r) {</pre>	func(in	t, int,	int):	
<pre>switch (x) {</pre>		cmpl	<b>\$5, %edi</b>	
case 0: Blank 0		ja	.L9	
case 1: Blank 1		movl	%edi, %edi	
case 2: Blank 2		jmp	*.L4(,%rdi,8)	
	.L4:			(hint, this is the
		.quad	.L8	<pre>jump table!)</pre>
Case 4: Blank 4		.quad	.L7	
case 5: <u>Blank 5</u>		.quad	.L6	
}		.quad	.L5	
return <u>Blank 6</u> ;		.quad	.L5	
}		.quad	.L3	
	.L7:			
		addl	<b>\$6, %edx</b>	
	.L6:			
		leal	(%rdx,%rdx,4),	%eax
		ret		
	.L5:			
		leal	(%rdx,%rsi,2),	%eax
		ret		
	.L3:			
		movl	%edx, %eax	
		xorl	%esi, %eax	
		ret		
	.L8:			
		movl	<b>\$-4,</b> %eax	
		Ret		
	.L9:			
		movl	%edx, %eax	
		ret		

Reverse engineer the assembly code on the previous page to figure out what each case of the switch-case statement is doing. Don't forget about "break" statements!

Blank 0	
Blank 1	
Blank 2	
Blank 3	
Blank 4	
Blank 5	
Blank 6	

### **Question 7. Floating Point Mystery (4pts)**

A long time ago, we used to put floating-point questions on the datalab. I found a solution to one of these problems lying around, but can't figure out what it's doing anymore:

```
unsigned mystery_function(unsigned uf) {
  unsigned sign = uf >> 31;
  unsigned exp = uf>>23 & 0xFF;
  unsigned frac = uf & 0x7FFFFF;
  if (exp == 0) {
   frac = 2*frac;
    if (frac > 0x7FFFFF) {
      frac = frac & 0x7FFFFF;
      exp = 1;
    }
  } else if (exp < 0xFF) {</pre>
    exp++;
    if (exp == 0xFF) {
      frac = 0;
    }
  }
  return (sign << 31) | (exp << 23) | frac;
}
```

- 1. In what cases will this function return the same thing as the input argument? (1pt)
- 2. What does the above function do? (3pts)

#### Question 8. Two Birds, One Bomb (6pts)

```
Dump of assembler code for function phase 5:
   0x0000555555551c8 <+0>:
                                       $0x8,%rsp
                                sub
   0x00005555555551cc <+4>:
                                mov
                                       $0xa,%edx
   0x0000555555551d1 <+9>:
                                mov
                                       $0x0,%esi
   0x0000555555551d6 <+14>:
                                callg 0x55555555670 <strtol@plt>
   0x0000555555551db <+19>:
                                mov
                                       %rax,%rsi
   0x0000555555551de <+22>:
                                       0x2e5b(%rip),%rdi
                                lea
                                                                 # 0x55555558040 <<u>nodes</u>>
   0x00005555555551e5 <+29>:
                                callg 0x5555555559e <recurse>
   0x0000555555551ea <+34>:
                                       $0x21,%eax
                                cmp
   0x0000555555551ed <+37>:
                                       0x5555555551fe <phase 5+54>
                                je
   0x0000555555551ef <+39>:
                                mov
                                       $0x0,%eax
   0x0000555555551f4 <+44>:
                                callq 0x5555555517e <explode bomb>
   0x0000555555551f9 <+49>:
                                add
                                       $0x8,%rsp
   0x0000555555551fd <+53>:
                                reta
   0x0000555555551fe <+54>:
                                mov
                                       $0x0,%eax
   0x000055555555203 <+59>:
                                callq
                                       0x55555555569 <phase defused>
   0x0000555555555208 <+64>:
                                       0x5555555551f9 <phase_5+49>
                                jmp
Dump of assembler code for function recurse:
   0x00005555555519e <+0>:
                                       %rdi,%rdi
                                test
   0x00005555555551a1 <+3>:
                                je
                                       0x5555555551c2 <recurse+36>
   0x0000555555551a3 <+5>:
                                push
                                       %rbx
   0x0000555555551a4 <+6>:
                                       0x10(%rdi),%ebx
                                mov
   0x00005555555551a7 <+9>:
                                cmp
                                       %esi,%ebx
   0x0000555555551a9 <+11>:
                                       0x5555555551b8 <recurse+26>
                                ige
   0x0000555555551ab <+13>:
                                mov
                                       0x8(%rdi),%rdi
   0x0000555555551af <+17>:
                                callq 0x5555555519e <recurse>
   0x0000555555551b4 <+22>:
                                add
                                       %ebx,%eax
   0x0000555555551b6 <+24>:
                                       %rbx
                                рор
   0x0000555555551b7 <+25>:
                                retq
   0x0000555555551b8 <+26>:
                                       (%rdi),%rdi
                                mov
   0x0000555555551bb <+29>:
                                callq
                                       0x55555555519e <recurse>
   0x0000555555551c0 <+34>:
                                jmp
                                       0x5555555551b4 <recurse+22>
   0x0000555555551c2 <+36>:
                                mov
                                       $0x0,%eax
   0x00005555555551c7 <+41>:
                                reta
(gdb) x/21gx &nodes
0x555555558040 <nodes>:
                            0x0000555555558058
                                                0x000055555558070
                                                                     0x0000000000000000
0x555555558058 <nodes+24>:
                            0x0000555555558088
                                                0x0000555555580a0
                                                                     0x00000000000000005
0x5555555558070 <nodes+48>:
                            0x0000555555580b8
                                                0x0000555555580d0
                                                                     0x0000000000000000
0x555555558088 <nodes+72>: 0x0000000000000000
                                                0x0000000000000000
                                                                     0x00000000000000000
0x5555555580a0 <nodes+96>: 0x0000000000000000
                                                0x00000000000000000
                                                                     0x000000000000000007
0x5555555580b8 <nodes+120>: 0x0000000000000000
                                                0x0000000000000000
                                                                     0x000000000000000b
0x5555555580d0:<nodes+144>: 0x00000000000000000
                                                0x00000000000000000
                                                                     0x0000000000000011
```

So, it turns out that I need a new phase\_5 for the bomb lab, because the old phase\_5 is waaaay too easy to cheat on ... therefore, I made this phase\_5 for next year, with a similar kind of flavor.

Please help me make sure this new phase\_5 has no errors by solving it for me. I've printed out a couple functions and some memory values. Next year's class will surely thank you.; )

Note that "strtol" just converts an ascii string (specified by a char\* input) into an integer.

- 1. What datastructure is this problem concerned with? Please be specific. (1 pt)
- 2. What string passed to "phase\_5(char\* input)" will defuse the phase? (5 pts)