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Computer Networks

UCLA CS CS 118

### Midterm Exam

**Directions:** Write your name on the exam and on every page you submit. Write something for every question. Students who do not write something for everything lose out over students who write down wild guesses. You will get some points if you attempt a solution but nothing for a blank sheet of paper. Write something down, even wild guesses. Problems take long to read but can be answered concisely.

Question	Maximum	Score
1	40	30
2	15	7
3	15	6.5
4	15	8
5	15	14
Total		65.5

1, Overview, 40 points: Give the most important reason or answer you can think of for each of the following. Each answer should be 1 or 2 lines. Only one reason or answer please! All questions are 4 points apiece. Choose best choice

• 1) Layering and DiffServ: In the DiffServ proposal, in order to provide different qualities of service to different packets, routers at the edge of a customer premise look at the traffic type of the packet in the TCP headers and set some bits in the routing header (called the TOS bits). Core routers then use the TOS bits in the routing header to give priority to some packets over others. (a) The edge routers are committing a layer violation but the core routers are not b) The edge routers are not committing a layer violation but the core routers are c) They both committing a layer violation d) Neither is committing a layer violation

• 2) Fourier Analysis: The basis for Fourier Analysis is that wires are "nice" to sine waves. This means that when a sine wave passes through a wire one basic property of the sine wave stays the same. The property is: (a) frequency (b) phase (c) amplitude d) time it takes to reach a peak. from line

• 3) Media Impairments: A fiber link suffers from both chromatic and modal dispersion. For modal dispersion, for the fastest wavelength of light in the laser, the signal that goes down the middle takes 200 nsec while the signal that bounces around takes 250 nsec. For chromatic dispersion, the slowest wavelength of light is 10 % slower than the fastest wavelength. What is the fastest spacing between bits to avoid ISI? (a) One every 25 nsec (b) Once every 50 nsec (c) Once every 75 nsec d) Once every 100 nsec

• 4) Latency and throughput: Does reducing latency always increase throughput? (a) No Yes because reducing the time for a job always increases the number of jobs per second b) No because latency and throughput are completely unrelated c) Yes because throughput is the reciprocal of latency d) No because one always use a shorter wire of the same transmission speed to reduce propagation delay

• 5) Media: Many wireless keyboards use infrared to connect from the keyboard to the computer. What would go wrong (or right) if the keyboard used radio instead of infrared to communicate with the computer? a) The radio is more expensive than infrared b) The radio would be absorbed by obstacles c) Keyboard transmissions would interfere between nearby offices (d) The radio would allow the keyboard to be very far from the computer

• 6) Preambles and Transitions: Is 11111110 a good preamble for AMI a) No because there are not enough transitions to synch clock recovery b) No because it is hard to tell where the data starts c) Yes because it guarantees at least one transition (d) Yes, because it guarantees several transitions and you can tell where the data star

30/40

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4

4

4

2

4

- **7) End-to-end argument:** To make sure your data does not get corrupted today we have two mechanisms: each hop does a checksum like CRC 32 and TCP also does an end-to-end checksum. If TCP is doing an end-to-end checksum why is each router doing a CRC as well? a) Because we want to ensure the destination address is correct before forwarding packets to the next hop b) Because corruption on links is so common, it is faster to detect it hop by hop  c) Because each router can retransmit packets hop by hop d) Because routers can corrupt packets and the CRC can catch this corruption.

- **8) Data Link Sublayering:** Why framing should be performed before error detection at a receiver, rather than vice-versa. a) So that errors in the framing characters can be caught by the CRC b) So that we can know where the destination address and other headers are in the bit stream c) So we can know which protocol we are using before doing error detection  d) So the receiver can know where the checksum is

- **9) Data Link restarting:** Consider the following scheme where a sender restarts after a crash and sending 10 frames. When the sender comes up and starts sending new frames with sequence number 0, the receiver sends an Ack back for frame 10 (which is what it expects). Suppose the sender "jumps" to 10 and starts sending. Will this jumping method guarantee that all frames sent after the sender crash are received for other cases as well?  No because this means the sender has lost frames 0 through 9 b) No, because if the sender crashes after sending one frame, it could be fooled when it comes up and sends a new frame after a crash  c) Yes because we are only guaranteeing frames sent after the last crash d) Yes, because the receiver will keep sending acknowledgements unconditionally

- **10) Ethernet collisions:** Ethernet guarantees that a) All stations detect collisions at the same time b) Only senders detect collisions c)  If any frame does collide at some point in the wire, then all stations will eventually learn that the collision occurred d) that the CRC will reliably detect collisions

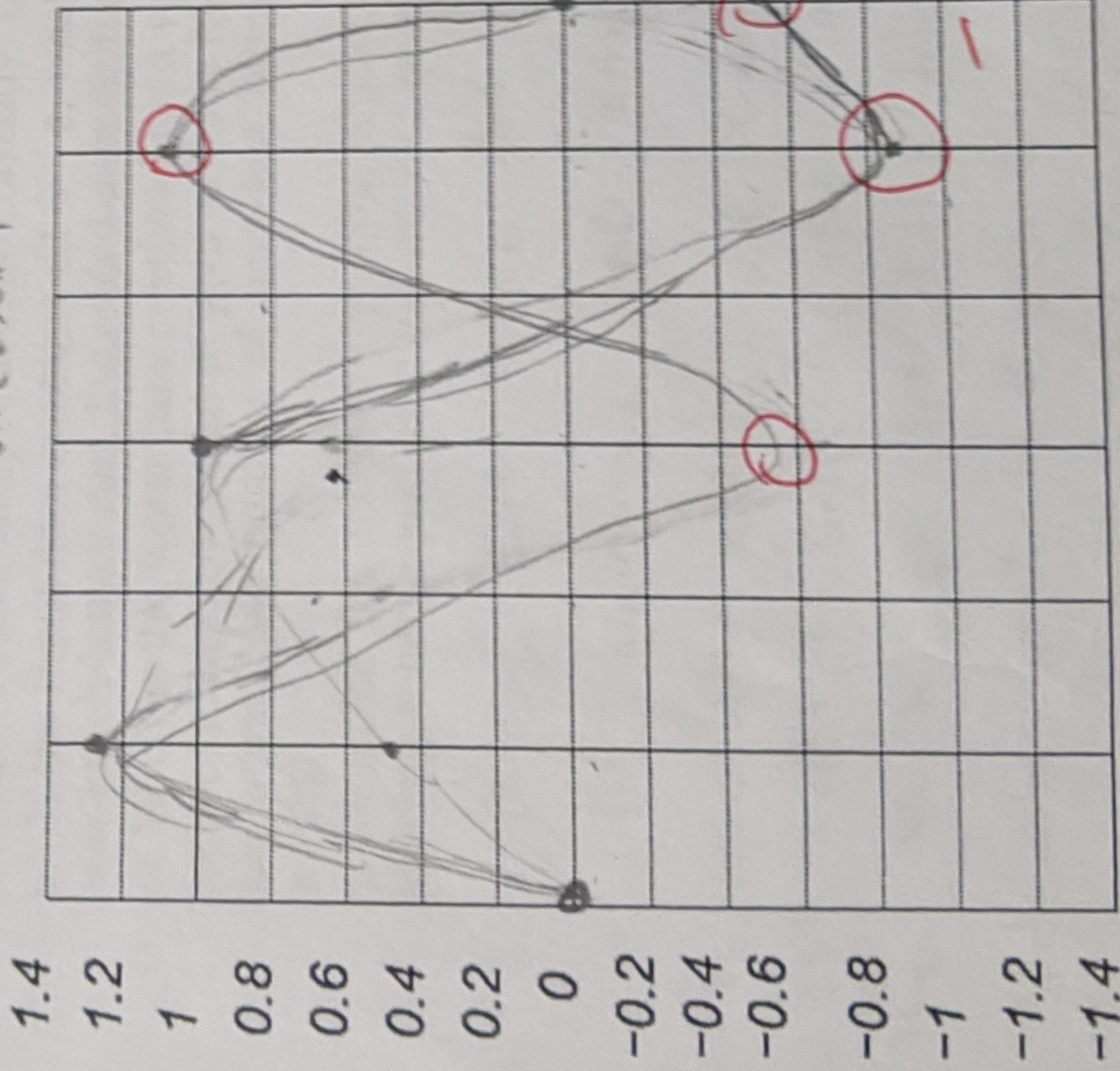
x	sin(x)	x	sin(x)	x	sin(x)
$\pi * 0/6$	0.000	$\pi * 3/6$	1.000	$\pi * 7/6$	-1.000
$\pi * 1/6$	0.500	$\pi * 4/6$	0.866	$\pi * 8/6$	-0.866
$\pi * 2/6$	0.866	$\pi * 5/6$	0.500	$\pi * 9/6$	-0.500

2. **Fourier Analysis:** An input signal's Fourier analysis has already been done for you. The signal is  $I = \sin(2\pi t) + 1/2 \sin(4\pi t) + 1/3 \sin(6\pi t)$ . We want to show that just three sine waves comes close to approximating a "nice" shape. To avoid the use of calculators, we have a table above which has the value of sine waves at various angles (in fractions of  $\pi$ ).

(7 points) Using the sine table above, first compute and plot the input sine wave at the 5 moments  $t = 0, 1/12, 1/4, 5/12, 1/2$ . Please draw the wave approximately on the ruled figure (which has amplitude spacing in units 0.25, and time spaced in units of  $1/12$  of a second).

(8 points) Assume the input is passed through a channel whose frequency response is such that all waves of frequency 1 Hz or less are passed through unscathed (and with no phase shift) while all waves of frequency greater than 1 Hz are scaled down to 0. Compute the amplitude of the output at the same 5 times and draw it on the same figure as the input. What is the final shape? (if you have no time, say what the output wave looks like).

The final signal is just a sine wave  
 $I = \sin(2\pi t)$



Input(0) = 0  
 Input(1/12) = 1.366  
 Input(3/12) = 2.15  
 Input(5/12) = 1.07  
 Input(6/12) = 0

Output(0) = 0  
 Output(1/12) = 0.5  
 Output(3/12) = 1.0  
 Output(5/12) = -0.866  
 Output(6/12) = -0.5

0 1/12 2/12 3/12 4/12 5/12 6/12

6/12 1/12 1/6 10/12 20/12 30/12  
 1/12 2/12 3/12 4/12 5/12 6/12  
 0 -1/12 -1/6 -1/4 -1/3 -1/2

**3. Framing, 20 points:** Bit stuffing can be expensive to implement if framing is done in software because it can cause the stuffed data not to be aligned on 8 bit boundaries. Thus many software framing techniques used on dialup lines do byte stuffing. Byte stuffing will guarantee that the stuffed data length is a multiple of 8. Alyssa P. Hacker has come up with what she calls the most efficient byte stuffing technique ever, much more efficient than even HDLC bit stuffing or PPP. Suppose the framing character is the zero byte (which we denote by 0). We want to code the data to remove all occurrences of 0. Alyssa's idea is as follows. We divide the data into blocks of up to 254 bytes such that: EITHER each block is less than 254 bytes and ends with a 0

OR the block is 254 bytes in length and does not contain a 0 byte. We then encode each block by placing the length of the block at the start, followed by the bytes in the block, but without the trailing 0 byte, if any. Notice that we can get rid of the zero byte because the length of each block tells us where to place the zero byte during decoding. The exception is if the length is 254 in which case the zero DOES NOT contain a trailing 0 and we use a length field of 255 to encode this case. (We cannot use 254 because that is reserved for a block of 253 bytes followed by a trailing 0). This last exception allows us to encode long sequences that do not contain a 0 byte. As an example if the input is XY00L0, where X, Y, L are arbitrary bytes, the first block is XY0, the second block is 0, the third block is L0. The first block encoded becomes 3 X Y, the second block encoded becomes 1, the third block encoded becomes 2 L.

- a) The problem does not specify how one encodes the data if the last block is less than 254 bytes in length and does not end in a 0. Obviously, one could try to pad the last block. What is the minimal amount of padding necessary to make sure we can use Alyssa's rules and yet decode correctly. (2 points)
- b.) If the input data contains all 0's, what is the length of the coded data? In general, how much longer can the coded data be than the original input data? (3 points)
- c.) What is the worst case overhead for a very large packet of length  $L \gg 254$ ? What is the worst case overhead of your scheme for a very small packet  $L \ll 254$ ? Overhead is the worst case number of non-data bytes divided by the number of data bytes. (5 points)
- d. Describe the receiver decoding algorithm briefly in words incorporating whatever padding scheme you used in Part a. (5 points)

a) 0 padding. We encode the last block with a 255 header. At the beginning of the block. This way receiver will not be incorrectly added. we can tell the receiver to read in 254 bytes or until the end of input (whichever comes first) when encountering 255.

b) The length of the coded data is equal to the length of the input data since each 0 will be encoded as 1. For small packets, overhead can be  $\frac{1}{3}$  (ex.  $X0 \rightarrow 2X0$ ). If every block while there are blocks: in the large packet case of the form length = first bytes of block  $X0$ , then the overhead would also be  $\frac{1}{3}$  if length  $\leq 254$ : (ex. length - 1 bytes from input to receiver buffer starting after length bytes appear 0 to receive better else: read in 254 bytes from input to receiver buffer starting after 254 bytes OR until end of input stream (if there are less than 254 bytes remaining)

4. CRCs, 20 points: Consider a CRC generator  $x^2 + 1$  and a message 110.
- (3 points) a) First compute the message together with the CRC. Show your working so we can give you partial credit.
  - (1 point) b) Suppose we wish to consider all error patterns in which the 5th bit is corrupted. What term must all such error polynomials start with?
  - (5 points) c) We wish to find all errors such that Bit 5 is corrupted but the error is not detected by the CRC. As in the HWs, the key is to use multiplication. Write down all polynomials the generator must be multiplied with to ensure that the 5th bit is corrupted and yet the result is undetectable. Note that unlike the case of a burst in your HW, we do not need the last term in the error polynomial to be 1 since we only want Bit 5 to be corrupted.
  - (5 points) d) Multiply each of these polynomials by the generator to find all possible undetectable error patterns with the 5th bit corrupted. Remember that when you multiply you use mod 2 addition unlike ordinary algebra. For example,  $x^2 + 1$  times mod 2 = 0.
  - (1 point) e) How many of these undetectable error patterns that corrupt Bit 5 have a total of 3 bit errors?

bits (CRC = 101)  
 $r = 3$

a) Shifted  $m = 11000$

$$\begin{array}{r} 101 \ 1100 \\ 0110 \\ \hline 101 \\ 011 \end{array}$$

Shifted  $m + remainder = 11011$

x1  
 b)  $x^4$

x1  
 e) 011

c)  $x^4$

$$\begin{array}{r} 10000 \\ 10001 \\ 10010 \\ 10011 \\ 10100 \\ 10101 \\ 10110 \\ 10111 \\ 11000 \\ 11001 \\ 11010 \\ 11011 \\ 11100 \\ 11101 \\ 11110 \\ 11111 \end{array}$$

d)  $x^2 + 1$

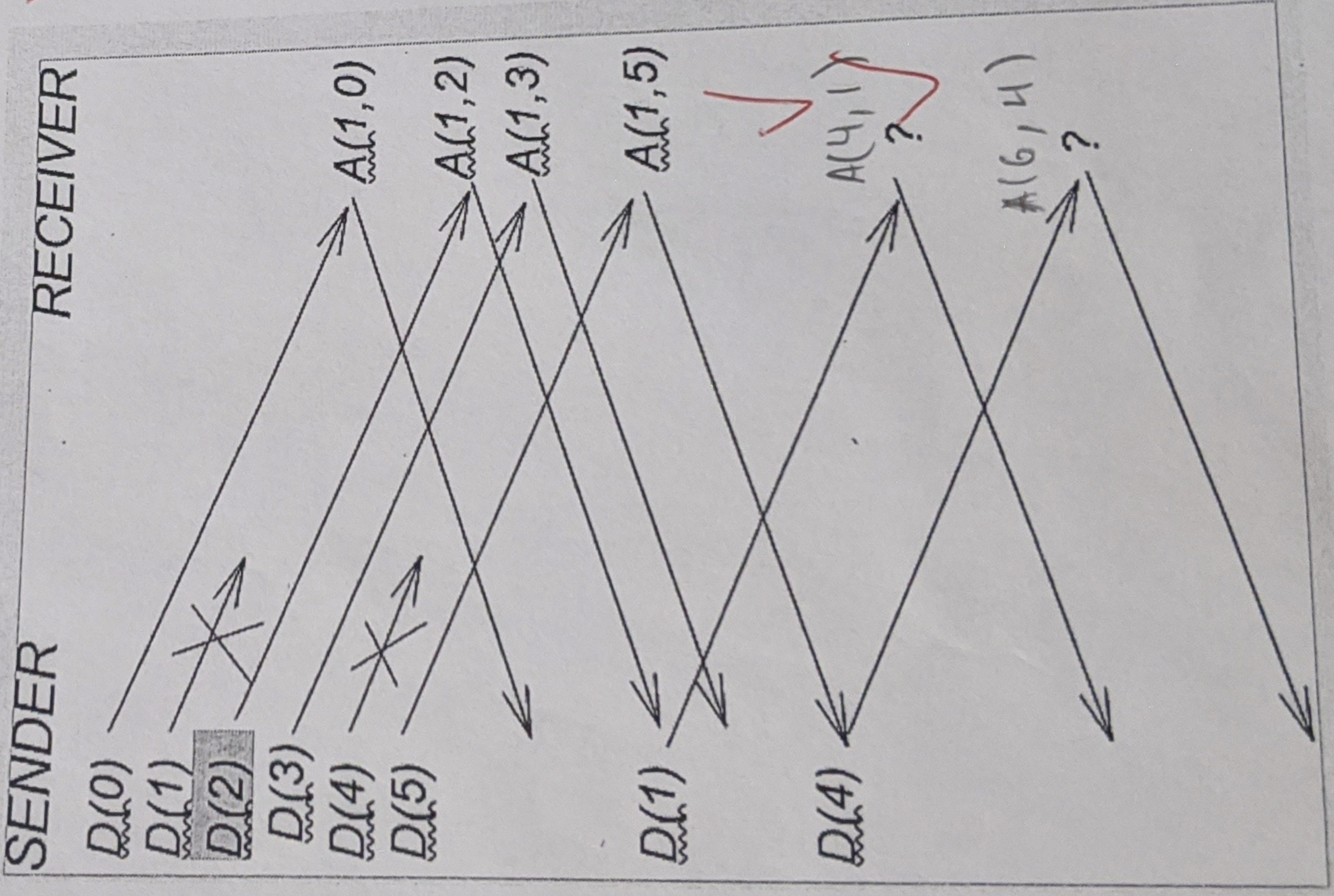
$$\begin{array}{r} x^6 + x^2 + x^0 + 1 \\ x^6 + x^3 + x^0 + x^1 + x \\ x^6 + x^3 + x^2 + x^0 + x^1 + x + 1 \\ x^6 + x^4 + x^1 + x^2 + x^0 + x^1 + x + 1 \\ x^6 + x^4 + x^3 + x^0 + x^1 + x^2 + x + 1 \\ x^6 + x^4 + x^3 + x^2 + x^0 + x^1 + x^2 + x + 1 \\ x^6 + x^5 + x^0 + x^1 + x^2 + x + 1 \end{array}$$

e)  $x^5$

$$\begin{array}{r} x^6 + x^5 + x^0 + x^1 + x^2 + x + 1 \\ x^6 + x^5 + x^1 + x^2 + x + 1 \\ x^6 + x^5 + x^2 + x^0 + x^1 + x^2 + x + 1 \\ x^6 + x^5 + x^3 + x^0 + x^1 + x^2 + x + 1 \\ x^6 + x^5 + x^3 + x^2 + x^0 + x^1 + x^2 + x + 1 \\ x^6 + x^5 + x^4 + x^0 + x^1 + x^2 + x + 1 \\ x^6 + x^5 + x^4 + x^2 + x^0 + x^1 + x^2 + x + 1 \\ x^6 + x^5 + x^4 + x^3 + x^0 + x^1 + x^2 + x + 1 \\ x^6 + x^5 + x^4 + x^3 + x^2 + x^0 + x^1 + x^2 + x + 1 \\ x^6 + x^5 + x^4 + x^3 + x^2 + x^0 + x^1 + x^2 + x + 1 \end{array}$$

- 5. **SMART Retransmission Strategy, 20 points:** The SMART retransmission strategy is a retransmission strategy that tries to do as well as selective reject without receiver sends the same cumulative ack as before but also sends one extra number the number of the received frame that caused the ack to be sent. To see how this works, imagine the sender sends frames (see figure below) with sequence numbers 0, 1, 2, 3, 4, 5 but frame 1 and 4 get lost. On receiving 0, the receiver will send back Ack (1,0); on receiving 2, the receiver will send back Ack (1, 2); on receiving 3, the receiver sends back (1, 3); and on receiving 5, the receiver sends back (1, 5).
- When the sender receives Ack(1,2) why does the sender know enough to retransmit D(1)? (2 points)
- When sender receives Ack (1,3) why does the sender know enough to retransmit D(2)? (2 points)
- What information does the sender need to maintain to not take any action at this point? (6 points)
- Write down on the picture (in place of the question number) the acks that are sent after receiver gets D(1) and D(4). (4 points)
- When the ack to the retransmitted D(1) arrives, how does the sender know not to retransmit D(4)? (4 points)
- Why is this scheme better than the fast retransmit scheme (done by TCP) where 3 duplicate acks trigger a retransmission. (4 points)

a)  
b)  
c)  
d)  
e)



a) The first number in the Ack tells the sender what the receiver has received up to but not including the number. In this case, the sender knows the receiver has received all messages up to but not including D(1), so it needs to retransmit D(1).

b) The sender needs to store the second number in the Ack's sent by the receiver. This way, the sender would have stored the 2 from the A(1,2). It received earlier and now retransmits D(2).

c) It just sent out D(1) and the timer hasn't expired for retransmission.

e) The fast retransmit scheme requires a delay of at least 3x propagation delay before a duplicate frame gets sent. Since the prior frame hasn't been received yet, it requires no duplicate Acks to trigger a retransmission.

7. (1)

extra bytes added = 0-0-0-0-1