

[CS M51A FALL 14] SOLUTIONS FOR MIDTERM EXAM

Date: 11/04/14

Problem 1 (10 points)

$X = (x, y, z)$ is a 3-digit weighted mixed-radix number system: x is a radix-16 digit, y is a radix-3 digit, and z is a radix-12 digit.

1. (5 points) Convert $X=(6, 1, 11)$ to a decimal number.

Solution $X = 6 \times 3 \times 12 + 1 \times 12 + 11 = 216 + 12 + 11 = 239$

2. (5 points) What is the largest number of X in decimal?

Solution $Largest X = (15, 2, 11) = 15 \times 3 \times 12 + 2 \times 12 + 11 = 575$

Problem 2 (10 points)

Simplify the following boolean expression by using postulates of Boolean Algebra.

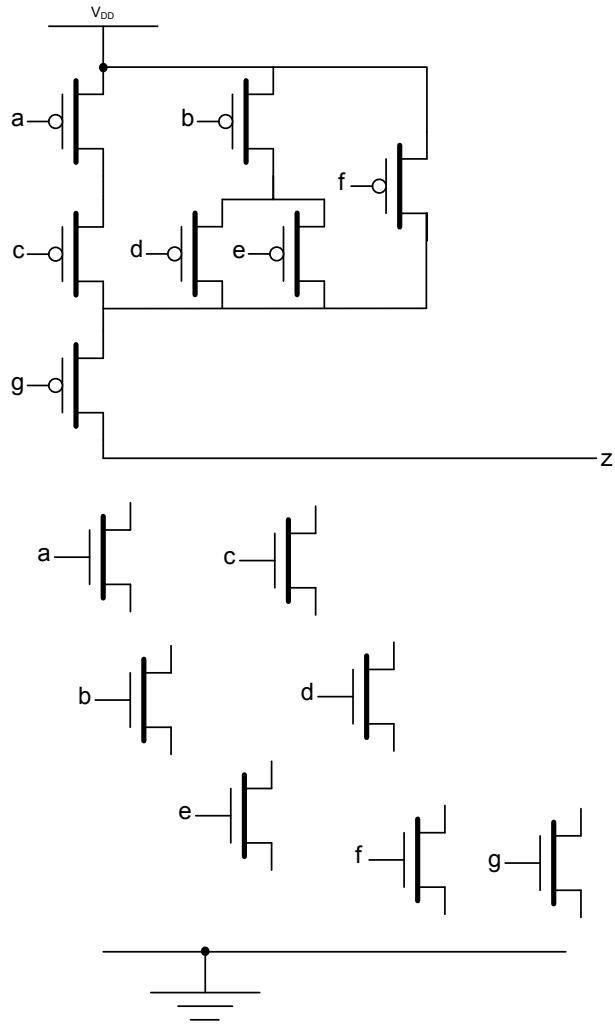
$$(a'b' + c)(a + b)(b' + a'c)'$$

Solution

$$\begin{aligned} & (a'b' + c)(a + b)(b' + a'c)' \\ = & (a'ab' + ac + a'b'b + bc)(b' + a'c)' \\ = & (ac + bc)(b' + a'c)' \\ = & (a + b)c \cdot b(a + c) \\ = & (a + b)(a + c)bc \\ = & (a + ab + ac + bc)bc \\ = & (a(1 + b + c) + bc)bc \\ = & (a + bc)bc \\ = & abc + bc \\ = & bc(a + 1) \\ = & bc \end{aligned}$$

Problem 3 (20 points)

We are given the following partial CMOS network.



1. **(10 points)** Write the expression for the pull-up network. From this, derive the expression for the pull-down network using switching algebra.

Solution From the given network we can directly write for the pull-up network

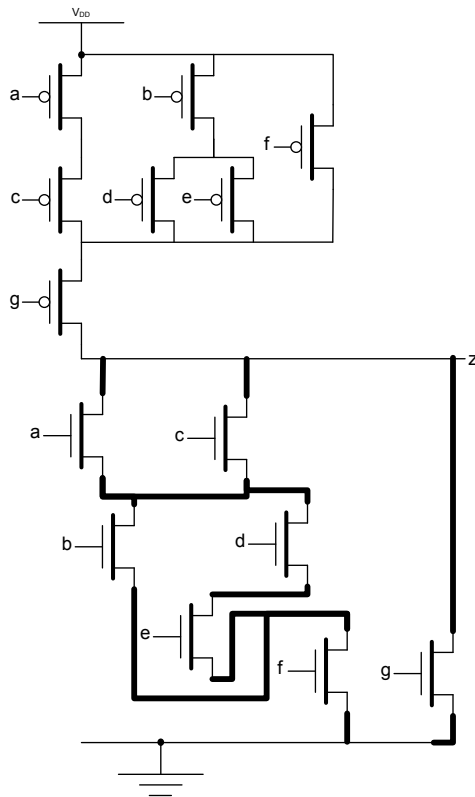
$$z = (a'c' + b'(d' + e') + f')g'$$

Therefore, the expression for the pull-down network is

$$\begin{aligned} z' &= [(a'c' + b'(d' + e') + f')g']' \\ &= (a'c' + b'(d' + e') + f')' + g \\ &= (a'c')'(b'(d' + e'))'f + g \\ &= (a + c)(b + (d' + e')')f + g \\ &= (a + c)(b + de)f + g \end{aligned}$$

2. (10 points) Connect the NMOS transistors to complete the pull-down network so that it corresponds to the expression obtained from part 1 and drives the output z to a valid output - i. e. either V_{DD} or ground - for any combination of inputs. Please only add additional wiring.

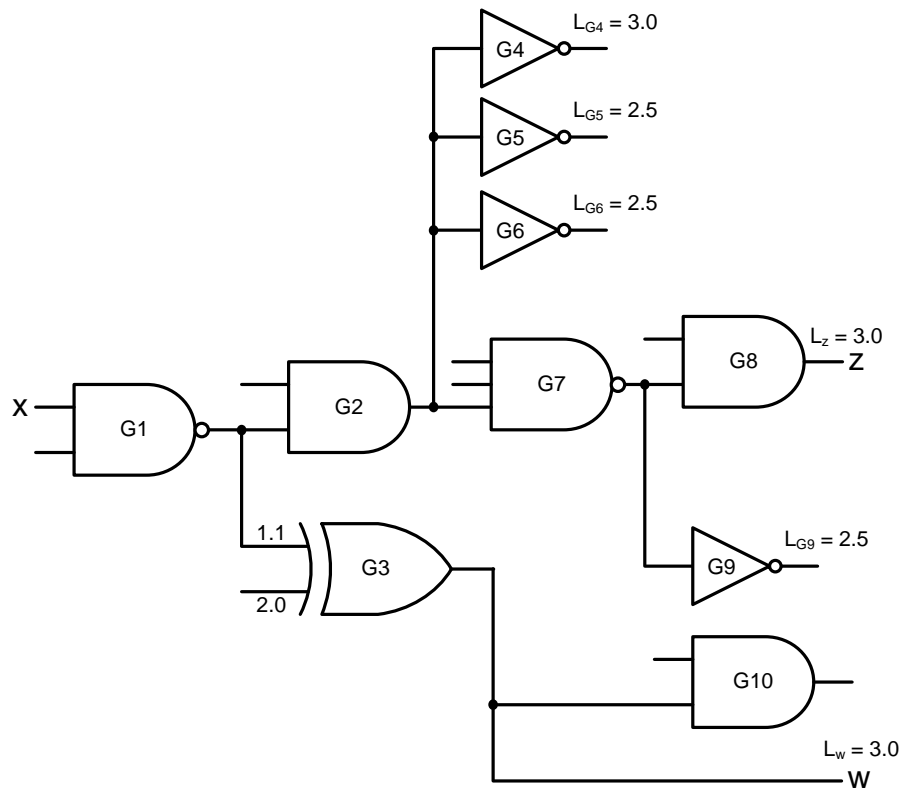
Solution Using the expression for z' , we can connect the circuit as shown. The additional lines are marked in bold.



Problem 4 (20 points)

We would like to determine the propagation delay of the gate network shown here. The output is z with a input x . The necessary gate characteristics are given in the table below.

Gate Type	Fan-in	Propagation Delays (ns)		Load Factor I
		t_{pLH}	t_{pHL}	
NOT	1	$0.02 + 0.038L$	$0.05 + 0.017L$	1.0
AND	2	$0.15 + 0.037L$	$0.16 + 0.017L$	1.0
NAND	2	$0.05 + 0.038L$	$0.08 + 0.027L$	1.0
NAND	3	$0.07 + 0.038L$	$0.09 + 0.039L$	1.0
XOR	2	$0.30 + 0.036L$	$0.30 + 0.021L$	1.1
		$0.16 + 0.036L$	$0.15 + 0.020L$	2.0



1. (10 points) (a) Determine the output load of gate G3.

Solution

The output load of gate G3 is 4, since this output is connectd to G10 (load 1), and W (load 3).

- b) If the fanout factor of gate G3 is 8, then how many additional gate inputs with load factor of 2 can be connected to the output of gate G3?

Solution

The output load of gate G3 is 4. The remaining fanout of gate G3 is 4 loads. Thus, 2 gate inputs with load factor 2 can be connected.

2. (10 points) Find the worst case value of $t_{pLH}(x \rightarrow z)$. Fill in the blanks below with the appropriate values.

Solution

Gate type and fan-in G1: NAND2 \rightarrow G2: AND2 \rightarrow G7: NAND3 \rightarrow G8: AND2

LH / HL G1: HL \rightarrow G2: HL \rightarrow G7: LH \rightarrow G8: LH

Output load L G1: 2.1 \rightarrow G2: 4.0 \rightarrow G7: 2.0 \rightarrow G8: 3.0

For the propagational delay values:

$$\text{G1: } 0.08 + 0.027L = 0.08 + 0.027 \cdot 2.1 = 0.1367$$

$$\text{G2: } 0.16 + 0.017L = 0.16 + 0.017 \cdot 4.0 = 0.228$$

$$\text{G7: } 0.07 + 0.038L = 0.07 + 0.038 \cdot 2.0 = 0.146$$

$$\text{G8: } 0.15 + 0.037L = 0.15 + 0.037 \cdot 3.0 = 0.261$$

$$t_{pLH}(x \rightarrow z) = 0.1367 + 0.228 + 0.146 + 0.261 = 0.7417 \text{ (ns)}$$

Problem 5 (10 points)

Decide whether the function E defined below is a universal function. If Yes, show your proof. You can use constant 0 or 1 as your input.

(a) **(5 points)**

x	y	$E(x, y)$
0	0	1
0	1	1
1	0	0
1	1	1

Solution Yes, E is a universal function.

$$E(x, y) = x' + y$$

$$\text{NOT gate: } E(x, 0) = x'$$

$$\text{OR gate: } E(E(x, 0), y) = E(x', y) = x + y$$

(b) (5 points)

x	y	z	E(x,y,z)
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	1

Solution Yes, E is a universal function.

$$E(x, y, z) = x'yz + xy'z' + xyz' + xyz = zy + z'x$$

$$\text{NOT gate: } E(1, 0, z) = z'$$

$$\text{AND gate: } E(0, y, z) = yz$$

Problem 6 (30 points)

F is a function that accepts inputs $x \in \{0, 1, 2\}$, $y \in \{0, 1, 2, 3\}$, and outputs $z = x^2 + y$. Suppose you use binary code to encode x , y , and z . x is encoded as x_1x_0 , y is encoded as y_1y_0 , z is encoded as $z_2z_1z_0$.

(a) **(6 points)** Fill in the following truth table.

x_1	x_0	y_1	y_0	z_2	z_1	z_0
0	0	0	0			
0	0	0	1			
0	0	1	0			
0	0	1	1			
0	1	0	0			
0	1	0	1			
0	1	1	0			
0	1	1	1			
1	0	0	0			
1	0	0	1			
1	0	1	0			
1	0	1	1			
1	1	0	0			
1	1	0	1			
1	1	1	0			
1	1	1	1			

- (b) **(6 points)** Based on the truth table in (a), draw the K-map for z_2 , z_1 , and z_0 .
- (c) **(6 points)** Use the K-maps in (b) to find all the prime implicants for z_2 , z_1 , and z_0 respectively.
- (d) **(6 points)** Use the K-maps in (b) to find all the essential prime implicants for z_2 , z_1 , and z_0 respectively.
- (e) **(6 points)** Implement z_2 , z_1 , and z_0 using minimal NAND-NAND networks. Note that each output has a separate gate network. You can directly use x_1, x_0, y_1, y_0 as inputs.

Solution

(a).

x_1	x_0	y_1	y_0	z_2	z_1	z_0
0	0	0	0	0	0	0
0	0	0	1	0	0	1
0	0	1	0	0	1	0
0	0	1	1	0	1	1
0	1	0	0	0	0	1
0	1	0	1	0	1	0
0	1	1	0	0	1	1
0	1	1	1	1	0	0
1	0	0	0	1	0	0
1	0	0	1	1	0	1
1	0	1	0	1	1	0
1	0	1	1	1	1	1
1	1	0	0	-	-	-
1	1	0	1	-	-	-
1	1	1	0	-	-	-
1	1	1	1	-	-	-

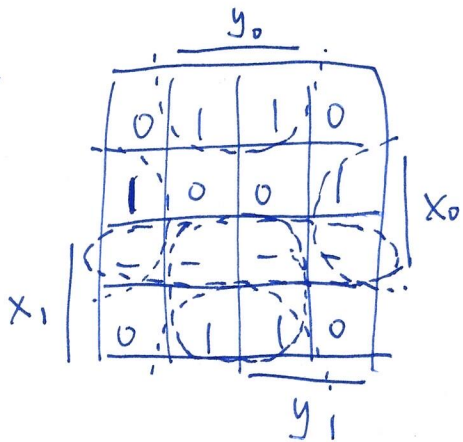
(b)

z_0	y_0				
	0	1	1	0	
	1	0	0	1	x_0
	-	-	-	-	
x_1	0	1	1	0	
	y_1				

z_1	y_0				
	0	0	1	1	
	0	1	0	1	x_0
	-	-	-	-	
x_1	0	0	1	1	
	y_1				

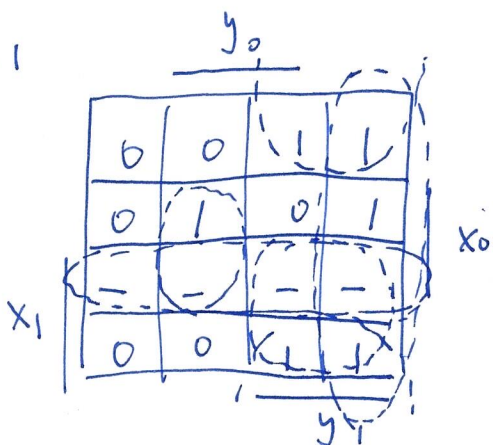
z_2	y_0				
	0	0	0	0	
	0	0	1	0	x_0
	-	-	-	-	
x_1	1	1	1	1	
	y_1				

(c). z_0 .



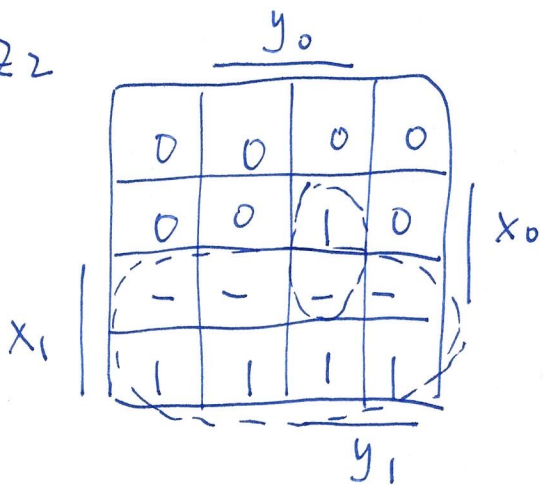
$$\begin{cases} y_0 x_0' \\ y_0' x_0 \\ x_1 y_0 \\ x_1 x_0 \end{cases}$$

z_1



$$\begin{cases} y_0 x_0 y_1' \\ x_1 x_0 \\ x_1 y_1 \\ x_0' y_1 \\ y_0' y_1 \end{cases}$$

z_2



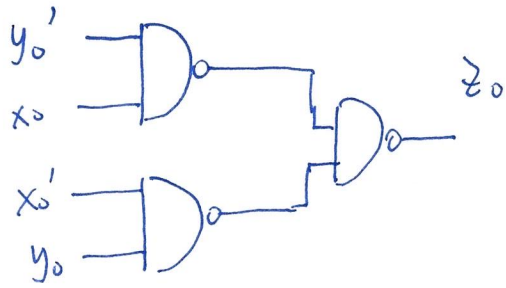
$$\begin{cases} x_1 \\ \cancel{x_1' x_0} \\ x_0 y_0 y_1 \end{cases}$$

(d). $z_0: y_0' x_0, x_0' y_0$

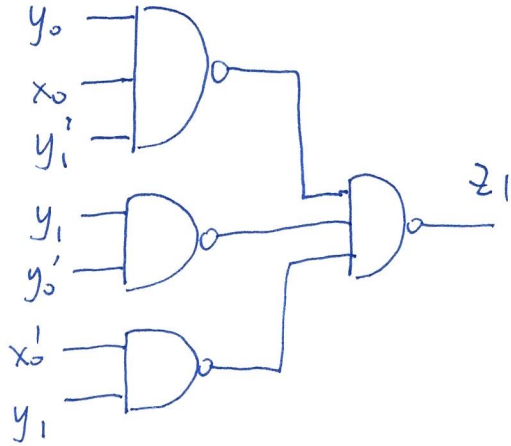
$z_1: y_0 x_0 y_1', y_1 y_0', x_0' y_1$

$z_2: x_1, x_0 y_0 y_1$

(e). z_0



z_1



z_2

