



CS 33

Midterm #2

All answers must be written on the answer sheet (last page of the exam).

All work should be written directly on the exam, use the backs of pages if needed.

This is an open book, open notes quiz – but you cannot share books or notes. An ASCII table is on the second to last page if you need it.

I will follow the guidelines of the university in reporting academic misconduct – please do not cheat.

NAME: _____

ID: _____

Problem 1: 30

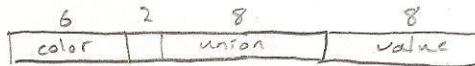
Problem 2: (15) → e.c.

Problem 3: 36

Total: 89 66/85 + 15

1. **Structured Play (30 points):** Consider the following declaration:

```
struct node_t {
    char color[6]; 6
    union {
        int numeric;
        char label[8]; } 8
    } identifier;
    long value; 8
} a[4];
```



Answer the following questions on how this data structure would be laid out on a 64-bit Linux machine:

- a) Considering alignment – how much total space (in bytes) would this data structure require?

96

- b) If the base address of array *a* is 0x600a60, what would be the string stored in *a*[1].*color*? Make use of the following gdb output:

(gdb) x/24x 0x600a60					
0x600a60:	0x68636f6d	0x00000a61	0x00000400	0x00000000	0x00000000
0x600a70:	0x46e87cc0	0x00000000	<u>0x7675616d</u>	0x00000a65	0x00000000
0x600a80:	0x00000800	0x00000000	0x3d1b58ba	0x00000000	0x00000000
0x600a90:	0x72757a61	0x00000a65	0x00001000	0x00000000	0x00000000
0x600aa0:	0x507ed7ab	0x00000000	0x7268636f	0x00000a65	0x00000000
0x600ab0:	0x00000200	0x00000000	0x2eb141f2	0x00000000	0x00000000

6d, 61, 75, 76, 65, 0a
m a u v e

Hint – don't forget that gdb reverses byte ordering within each 4-byte chunk. So in the following dump:

```
(gdb) x/4x 0x00111110
0x111110: 0x33221100 0x77665544 0xBBAA9988 0xFFEEDDCC
```

This prints out 16 bytes of memory starting at address 0x111110. In this example, the 16 bytes of memory starting at 0x111110 would contain, in order from lowest address (0x111110) to highest address (0x11111F):

00112233445566778899AABBCCDDEEFF

So address 0x111110 contains the byte 0x00, address 0x111111 contains the byte 0x11, address 0x111112 contains the byte 0x22, and so on. So in terms of just the least significant hex place of the address, gdb is actually printing out addresses in the following order:

3 2 1 0 7 6 5 4 B A 9 8 F E D C

This is useful when reading words, but can be confusing for other values.

2. **Complete Dis-Array (30 points):** Consider the following C fragment:

```
#define SIZE 10

int orange[SIZE][SIZE];
int *green;
int *purple[SIZE];
int n;

int main( int argc, const char* argv[] ) {
    ...
}
```

The code in ...'s will create three two dimensional arrays: orange, purple, and green. Arrays orange and purple are statically sized (i.e. size is known at compile time). Array orange is a nested array and array purple is a multi-level array. Array green is dynamically sized (i.e. a dynamically nested array).

Each array has a function that sets the value of one element of the array. These functions have three parameters – the row position in the array (i), the column position in the array (j), and the value to set (val). They have the following prototypes:

```
row col value
void setorange(int i, int j, int val);
void setgreen(int i, int j, int val);
void setpurple(int i, int j, int val);
```

For example, *setorange(i,j,val)* will set *orange[i][j]* to the value *val*.

Based on the information above, answer the following questions:

- a) Which of these three functions (*setorange*, *setgreen*, or *setpurple*) is shown disassembled below?

8048450:	55	push	%ebp
8048451:	8b 15 d4 9a 04 08	mov	0x8049ad4, %edx
8048457:	89 e5	mov	%esp, %ebp
8048459:	8b 45 08	mov	0x8(%ebp), %eax
804845c:	0f af 05 d0 9a 04 08	imul	0x8049ad0, %eax
8048463:	8b 4d 10	mov	0x10(%ebp), %ecx val
8048466:	03 45 0c	add	0xc(%ebp), %eax
8048469:	89 0c 82	mov	%ecx, (%edx,%eax,4)
804846c:	5d	pop	%ebp
804846d:	c3	ret	

- b) In the disassembled code above, what variable is located at address 0x8049ad0?

1011
0101
0010
0001

3. We Want the Func (40 points): Consider the following C code:

```
int foo(int i, int count) {  
  
    if (count==32)  
        return 0;  
    else  
        return (foo(i>>1, count+1)+(i&1));  
}  
  
int main( int argc, const char* argv[] ) {  
    fprintf(stderr, "%d\n", foo(atoi(argv[1]), 0));  
}
```

foo(11, 0) → foo(5, 1) + 1
↓
foo(2, 2) + 1
↓
foo(1, 3) + 0
↓
foo(0, 4) + 1

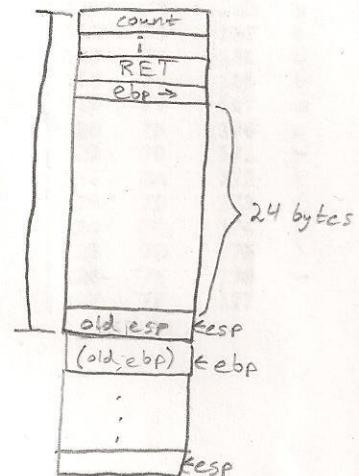
a) If this code were compiled, and then executed with the value 11 as a parameter, what would it return?

3

b) Consider the following disassembled code for function foo:

08048414 <foo>:

```
8048414:    push    %ebp  
8048415:    mov     %esp,%ebp  
8048417:    sub     $0x18,%esp  
804841a:    cmpl    $0x20, A(%ebp)  
804841e:    jne     8048427 <foo+0x13>  
8048420:    mov     $0x0,%eax  
8048425:    jmp     8048446 <foo+0x32>  
8048427:    mov     B(%ebp),%eax  
804842a:    lea     0x1(%eax),%edx  
804842d:    mov     C(%ebp),%eax  
8048430:    sar     %eax  
8048432:    mov     count,%edx, D(%esp)  
8048436:    mov     %eax, E(%esp)  
8048439:    call    8048414 <foo>  
804843e:    mov     F(%ebp),%edx  
8048441:    and     $0x1,%edx  
8048444:    add     %edx,%eax  
8048446:    leave  
8048447:    ret
```



There are six blanks in the code above – labeled A-F – and all are related to displacements relative to either %esp or %ebp. Fill in the blanks with the appropriate values to make this code work correctly. You should fill in a **hexadecimal number** for each blank – it should **not** be a register specifier. These are displacements relative to the stack or frame pointers.

- A 0xC
- B 0xC
- C 0x8
- D 0x8
- E 0x4
- F 0x8

Answer Sheet

Name: _____

1. a. 96 bytes (24 bytes per element
of the array) 15
15

b. "mauve\n"

2. a. setpurple

b. n

15

(ec)

3. a. 3

16

b. Fill in all blanks below

0xc

A

4

0xc

B

4

0x8

C

4

0x8

D

2

0x4

E

2

0x8

F

4