



1. **C If You Can Solve This (28 points):** The following problem assumes the following declarations:

```
int x = rand();
int y = rand();
int z = rand();
float f= foo(); //f is not NaN
unsigned ux = (unsigned) x;
unsigned uy = (unsigned) y;
```

For the following C expressions, circle either Y or N (but not both).

- |   | Always True?                          |                                       |
|---|---------------------------------------|---------------------------------------|
| a. $((\text{float}) x + f) - f == (\text{float}) x$   | Y                                     | <input checked="" type="checkbox"/> N |
| <i>Let <math>x=1</math> and <math>f=1e20</math>, rounding will cause the lhs to evaluate to zero.</i> |                                       |                                       |
| b. $(ux*uy) == (x*y)$   | <input checked="" type="checkbox"/> Y | N                                     |
| <i>Mixed integers will be implicitly casted to unsigned</i>   |                                       |                                       |
| c. $((x\&8)   y) == y \implies (x\ll 28) > 0$   | Y                                     | <input checked="" type="checkbox"/> N |
| <i>Let <math>x = 0xFFFFFFFF</math> and <math>y = 0x8</math></i>                                       |                                       |                                       |
| d. $(x^y)^x + z == y+z$   | <input checked="" type="checkbox"/> Y | N                                     |
| <i>Property of xor, an example is in the inplace swap function from lecture and text</i>              |                                       |                                       |

Note that “ $\implies$ ” represents an *implication*.  $A \implies B$  means that you assume A is true, and your answer should indicate whether B should be implied by A – i.e. given that A is true, is B always true?

2. **This Problem Bytes (21 points):** Consider the following 8 bits:

10111100

We will interpret these bits in three different ways (assume the above is in big endian form):

a. An 8-bit unsigned integer

$$2^7+0^6+2^5+2^4+2^3+2^2+0^1+0^0=128+32+16+8+4=188$$

b. An 8-bit two's complement integer

$$-2^7+0^6+2^5+2^4+2^3+2^2+0^1+0^0=-128+32+16+8+4=-68$$

c. The 8-bit floating point format we covered in class

$\boxed{10111100}$   
s e f

$$\begin{aligned} &(-1)^s M \times 2^E, \quad M = 1.f, \quad E = e - \text{bias} \\ &\text{bias} = 7 \\ &M = 1.100 = 2^0 + (1/2^1) = 1.5 \\ &E = 0111 - 7 = 0 \end{aligned}$$

$$10111100 = -1.5 \times 2^0 = -3/2$$

2. **This Problem Bytes (21 points):** Consider the following 8 bits:

10101100

We will interpret these bits in three different ways (assume the above is in big endian form):

a. An 8-bit unsigned integer

$$2^7+0^6+2^5+0^4+2^3+2^2+0^1+0^0=128+32+8+4=172$$

b. An 8-bit two's complement integer

$$-2^7+0^6+2^5+0^4+2^3+2^2+0^1+0^0=-128+32+8+4=-84$$

c. The 8-bit floating point format we covered in class

$$\boxed{1} \boxed{0101} \boxed{100} \quad (-1)^s M \times 2^E, \quad M = 1.f, \quad E = e - \text{bias}$$

s   e   f   bias=7

$$M = 1.100 = 2^0 + (1/2^1) = 1.5$$
$$E = 0101 - 7 = -2$$

$$10101100 = -1.5 \times 2^{-2} = -3/8$$

2. **This Problem Bytes (21 points):** Consider the following 8 bits:

10011100

We will interpret these bits in three different ways (assume the above is in big endian form):

a. An 8-bit unsigned integer

$$2^7+0^6+0^5+2^4+2^3+2^2+0^1+0^0=128+16+8+4=156$$

b. An 8-bit two's complement integer

$$-2^7+0^6+0^5+2^4+2^3+2^2+0^1+0^0=-128+16+8+4=-100$$

c. The 8-bit floating point format we covered in class

$$\boxed{1} \boxed{0011} \boxed{100} \quad (-1)^s M \times 2^E, \quad M = 1.f, \quad E = e - \text{bias}$$

s   e   f   bias=7

$$M = 1.100 = 2^0 + (1/2^1) = 1.5$$
$$E = 0011 - 7 = -4$$

$$10011100 = -1.5 \times 2^{-4} = -3/32$$

2. **This Problem Bytes (21 points):** Consider the following 8 bits:

10110100

We will interpret these bits in three different ways (assume the above is in big endian form):

a. An 8-bit unsigned integer

$$2^7+0^6+2^5+2^4+0^3+2^2+0^1+0^0=128+32+16+4=180$$

b. An 8-bit two's complement integer

$$-2^7+0^6+2^5+2^4+0^3+2^2+0^1+0^0=-128+32+16+4=-76$$

c. The 8-bit floating point format we covered in class

$\boxed{1}$	$\boxed{0110}$	$\boxed{100}$	$(-1)^s M \times 2^E$ , $M = 1.f$ , $E = e - \text{bias}$
s	e	f	bias=7
			$M = 1.100 = 2^0 + (1/2^1) = 1.5$
			$E = 0110 - 7 = -1$

$$10110100 = -1.5 \times 2^{-1} = -3/4$$

3. ***And This One is a Pain in My Big Endian (51 points):*** Here's your chance to show your bomb lab skills. Below we show the disassembled function `func0()` – compiled on an ia32 machine. The function reads one integer, using `scanf()`. Your job is to ***provide the input to `scanf()` that will result in this function returning the value 42, 23, 0, 83.***

To help you with this task – we provide part of the C code for the function below:

```
int func0 (int j, int k)
{
    int a;
    int i;

    a=0; // a will be modified by the switch statement

    scanf("%d", &i); // THE INPUT YOU ARE FIGURING OUT

    switch(i){
        ...
    }

    return a;
}
```

where the ...'s above represent missing code you need to figure out.

In addition to the input that would result in this function returning a `42, 23, 0, 83`, please fill in the blanks we have on the answer key with the requested intermediate values that would help you answer this question. The next two pages contain everything you need for this problem.

Here's the function from objdump:

```

08048470 <func0>:
8048470:    55                push   %ebp
8048471:    89 e5            mov    %esp,%ebp
8048473:    83 ec 28        sub    $0x28,%esp
8048476:    8d 45 f4        lea   -0xc(%ebp),%eax
8048479:    89 5d f8        mov    %ebx,-0x8(%ebp)
804847c:    8b 5d 08        mov    0x8(%ebp),%ebx    first arg j
804847f:    89 75 fc        mov    %esi,-0x4(%ebp)
8048482:    8b 75 0c        mov    0xc(%ebp),%esi    second arg k
8048485:    89 44 24 04    mov    %eax,0x4(%esp)
8048489:    c7 04 24 64 86 04 08  movl  $0x8048664,(%esp)
8048490:    e8 07 ff ff ff  call  804839c <__isoc99_scanf@plt>
8048495:    83 7d f4 07    cmpl  $0x7,-0xc(%ebp)
8048499:    76 15          jbe   80484b0 <func0+0x40>
default 804849b:    89 d8          mov    %ebx,%eax
0,5 804849d:    89 f1          mov    %esi,%ecx
804849f:    8b 5d f8        mov    -0x8(%ebp),%ebx
80484a2:    8b 75 fc        mov    -0x4(%ebp),%esi
80484a5:    d3 e0          shl   %cl,%eax
80484a7:    89 ec          mov    %ebp,%esp        a=j<<k
80484a9:    5d            pop   %ebp
80484aa:    c3            ret
80484ab:    90            nop
80484ac:    8d 74 26 00    lea   0x0(%esi,%eiz,1),%esi
80484b0:    8b 45 f4        mov    -0xc(%ebp),%eax
80484b3:    ff 24 85 74 86 04 08  jmp   0x8048674(,%eax,4)  jump table addr
80484ba:    8d b6 00 00 00 00  lea   0x0(%esi),%esi
7 80484c0:    89 f0          mov    %esi,%eax
80484c2:    8b 75 fc        mov    -0x4(%ebp),%esi
80484c5:    21 d8          and   %ebx,%eax
80484c7:    8b 5d f8        mov    -0x8(%ebp),%ebx    a=j&k
80484ca:    89 ec          mov    %ebp,%esp
80484cc:    5d            pop   %ebp
80484cd:    c3            ret
80484ce:    66 90          xchg  %ax,%ax
1 80484d0:    8d 04 1e        lea   (%esi,%ebx,1),%eax  a=j+k
80484d3:    8b 5d f8        mov    -0x8(%ebp),%ebx
80484d6:    8b 75 fc        mov    -0x4(%ebp),%esi
80484d9:    89 ec          mov    %ebp,%esp
80484db:    5d            pop   %ebp
80484dc:    c3            ret
2 80484dd:    8d 76 00        lea   0x0(%esi),%esi
80484e0:    89 d8          mov    %ebx,%eax
80484e2:    8b 5d f8        mov    -0x8(%ebp),%ebx
80484e5:    29 f0          sub   %esi,%eax
80484e7:    8b 75 fc        mov    -0x4(%ebp),%esi    a=j-k
80484ea:    89 ec          mov    %ebp,%esp
80484ec:    5d            pop   %ebp
80484ed:    c3            ret
80484ee:    66 90          xchg  %ax,%ax
3 80484f0:    89 d8          mov    %ebx,%eax
80484f2:    89 f1          mov    %esi,%ecx
80484f4:    d3 f8          sar   %cl,%eax
80484f6:    09 f3          or    %esi,%ebx        a=j>>k
80484f8:    8b 75 fc        mov    -0x4(%ebp),%esi    a+=j|k

```

```

80484fb:    01 d8      add    %ebx,%eax
80484fd:    8b 5d f8   mov    -0x8(%ebp),%ebx
8048500:    89 ec     mov    %ebp,%esp
8048502:    5d       pop    %ebp
8048503:    c3       ret
8048504:    8d 74 26 00  lea   0x0(%esi,%eiz,1),%esi
4 8048508:    31 c0     xor    %eax,%eax
804850a:    eb ea     jmp    80484f6 <func0+0x86>
804850c:    8d 74 26 00  lea   0x0(%esi,%eiz,1),%esi
6 8048510:    89 f0     mov    %esi,%eax
8048512:    8b 75 fc   mov    -0x4(%ebp),%esi
8048515:    31 d8     xor    %ebx,%eax
8048517:    8b 5d f8   mov    -0x8(%ebp),%ebx
804851a:    89 ec     mov    %ebp,%esp
804851c:    5d       pop    %ebp
804851d:    c3       ret

```

And here is some gdb interaction which should prove useful:

```

(gdb) break *0x8048470
Breakpoint 1 at 0x8048470
(gdb) run
Starting program

```

Breakpoint 1, 0x08048470 in func0 () **breakpoints stop before executing the current line/instr, stack setup for func0 has not been executed.**

```

(gdb) p $esp
$1 = (void *) 0xffffd3dc
(gdb) p $ebp
$2 = (void *) 0xffffd408
(gdb) x/32x 0xffffd3dc
ret addr      arg1:j      arg2:k      These values may be different
0xffffd3dc:  0x08048572  0x00000030  0x0000001a  0x0000000a
0xffffd3ec:  0x006fdfff  0x080485b0  0x080483b0  0x006fdfff
0xffffd3fc:  0x00000000  0x080485b0  0x00000000  0xffffd488
0xffffd40c:  0x00581ce6  0x00000003  0xffffd4b4  0xffffd4c4
0xffffd41c:  0xf7ffd428  0x080483b0  0xffffffff  0x00567fc4
0xffffd42c:  0x0804827e  0x00000001  0xffffd470  0x00557a05
0xffffd43c:  0x00568ab0  0xf7ffd708  0x006fdfff  0x00000000
0xffffd44c:  0x00000000  0xffffd488  0xf67726a3  0xb9e6515c
(gdb) x/16x 0x8048660
addr 0x8048674 jump table
0x8048660 <__dso_handle>:  0x00000000  0x74006425  0x69747365  0x2020676e
0x8048670 <__dso_handle+16>:  0x000a6425  0x0804849b  0x080484d0  0x080484e0
0x8048680 <__dso_handle+32>:  0x080484f0  0x08048508  0x0804849b  0x08048510
0x8048690 <__dso_handle+48>:  0x080484c0  0x3b031b01  0x00000020  0x00000003

```

Hint – don't forget that gdb reverses byte ordering within each 4-byte chunk. So in the following dump:

```

(gdb) x/4x 0x00111110
0x111110:  0x33221100  0x77665544  0xBBAA9988  0xFFEEDDCC

```

This prints out 16 bytes of memory starting at address 0x111110. In this example, the 16 bytes of memory starting at 0x111110 would contain, in order from lowest address (0x111110) to highest address (0x11111F):

00112233445566778899AABBCCDDEEFF

So address 0x111110 contains the byte 0x00, address 0x111111 contains the byte 0x11, address 0x111112 contains the byte 0x22, and so on. So in terms of just the least significant hex place of the address, gdb is actually printing out addresses in the following order:

3 2 1 0 7 6 5 4 BA 9 8 FED C

This is useful when reading words, but can be confusing for other values.

Return Val	42	23	0	83
j	0x30 = 48	0x41 = 65	0x22 = 34	0x36 = 54
k	0x1a = 26	0x2a = 42	0x10 = 16	0x1d = 29
default addr	0x0804849b			
case	6	2	7	1